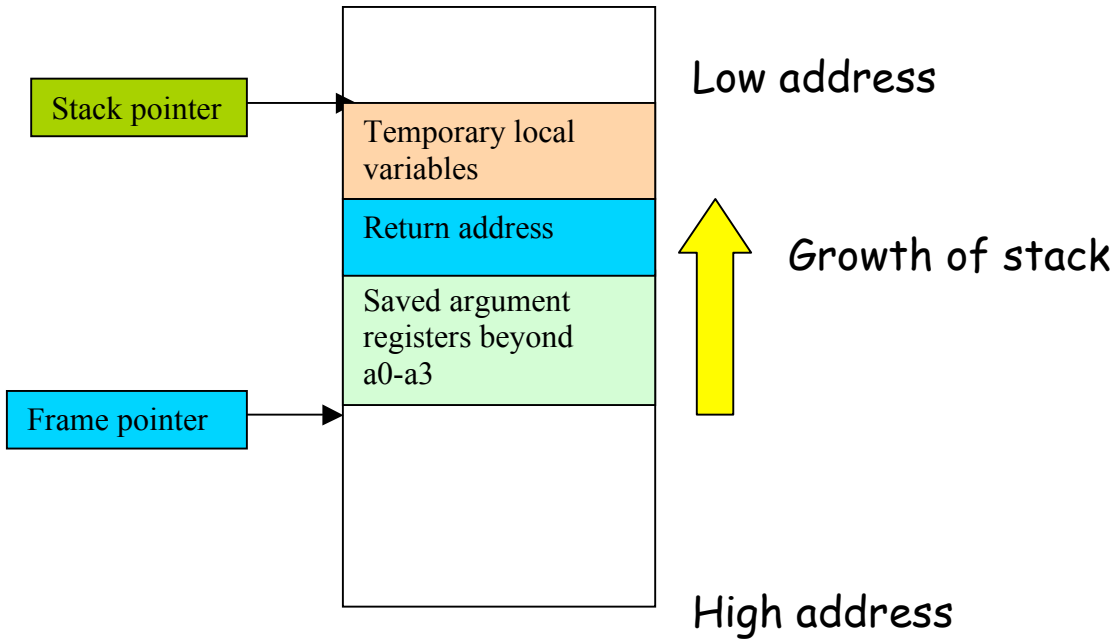
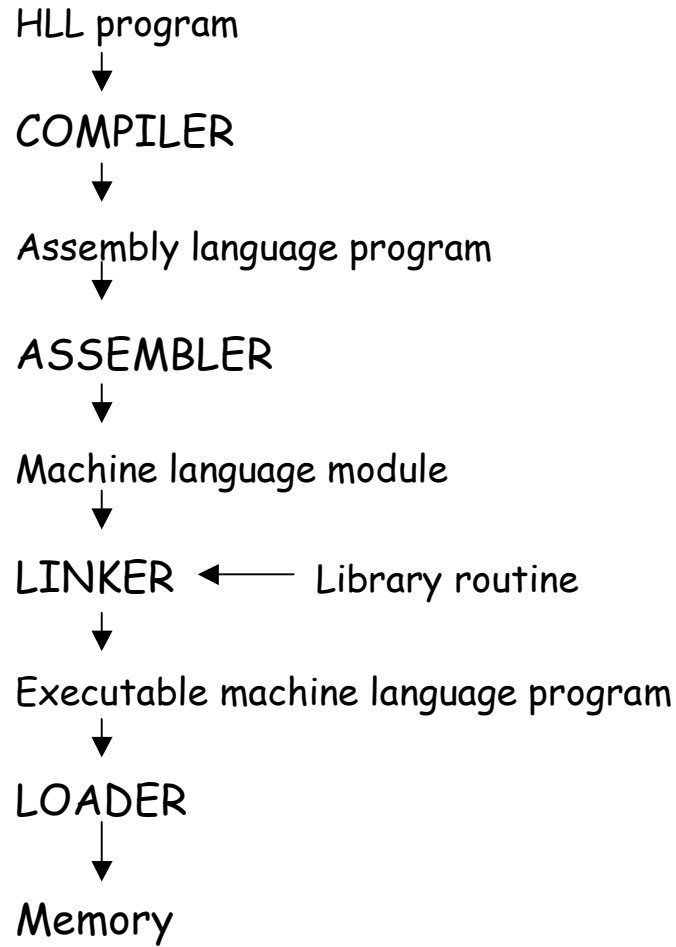


# Run time environment of a MIPS program



## A translation hierarchy



## What are Assembler directives?

Instructions that are not executed, but they tell the assembler about how to interpret something. Here are some examples:

```
. text
```

```
{Program instructions here}
```

```
. data
```

```
{Data begins here}
```

```
. byte 84, 104, 101
```

```
. asciiz "The quick brown fox"
```

```
. float f1, . . . , fn
```

```
. word w1, . . . . wn
```

# How does an assembler work?

In a **two-pass assembler**

PASS 1: Symbol table generation

PASS 2: Code generation

The operation of a two-pass assembler has been briefly explained in the class on 2/9/10.

## Other architectures

Not all processors are like MIPS.

### Example. Accumulator-based machines

A single register, called the **accumulator**, stores the operand before the operation, and stores the result after the operation.

Load	x	# into acc from memory
Add	y	# add y from memory to the acc
Store	z	# store acc to memory as z

Can we have an instruction

`add z, x, y # z := x + y, (x, y, z in memory) ?`

For some machines, YES, not in MIPS

## Load-store machines

MIPS is a **load-store architecture**. Only **load** and **store** instructions access the memory, all other instructions use registers as operands. What is the motivation?

*Register access is much faster than memory access, so the program will run faster.*

## Reduced Instruction Set Computers (RISC)

- The instruction set has only a small number of **frequently used instructions**. This lowers processor cost, without much impact on performance.
- All instructions have the same length.
- Load-store architecture.

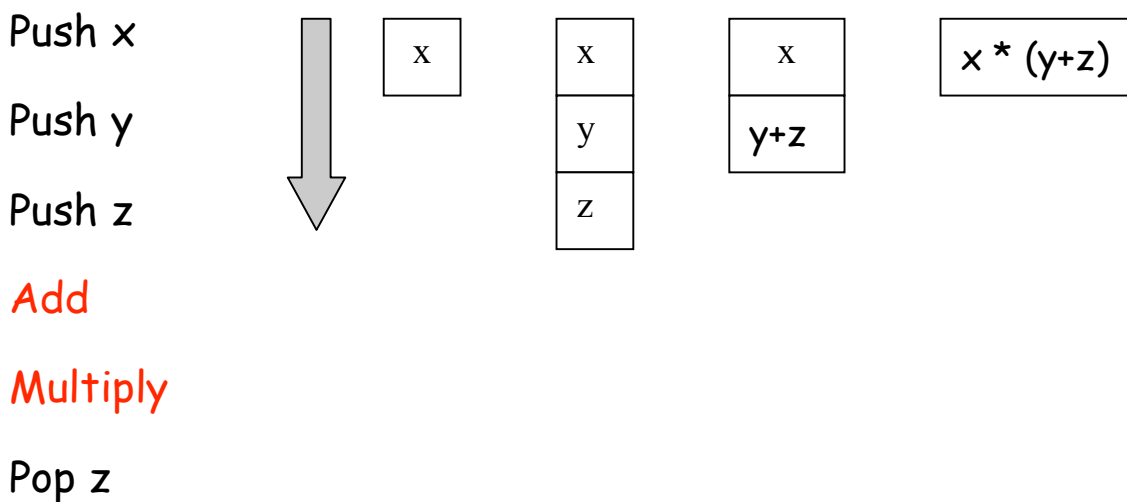
Non-RISC machines are called CISC

(Complex Instruction Set Computer). Example: Pentium

## Another classification

3-address	add r1, r2, r3	( $r1 \leftarrow r2 + r3$ )
2-address	add r1, r2	( $r1 \leftarrow r1 + r2$ )
1-address	add r1	(to the accumulator)
0-address or stack machines		(see below)

## Example of stack architecture



Computes  $z = x * (y + z)$

## Computer Arithmetic

How to represent negative integers? The most widely used convention is 2's complement representation.

$$+14 = 0,1110$$

$$-14 = 1,0010$$

Largest integer represented using n-bits is  $+ 2^{n-1} - 1$

Smallest integer represented using n-bits is  $- 2^{n-1}$

Review binary-to decimal and binary-to-hex conversions.

Review BCD (Binary Coded Decimal) and ASCII codes.

How to represent fractions?



## Overflow

$$+12 = 0,1100$$

$$+2 = 0,0010$$

add \_\_\_\_\_

$$+14 = 0,1110$$

$$+12 = 0,1100$$

$$+7 = 0,0111$$

\_\_\_\_\_ add

$$? = 1,0011$$

Addition of a positive and a negative number does not lead to overflow. **How to detect overflow?**

The following sequence of MIPS instructions can detect overflow in signed addition of \$t1 and \$t2:

```
addu $t0, $t1, $t2      # add unsigned
xor  $t3, $t1, $t2      # check if signs differ
slt  $t3, $t3, $zero     # $t3=1 if signs differ
bne  $t3 $zero, no_overflow
xor  $t3, $t0, $t1      # sum sign = operand sign?
slt  $t3, $t3, $zero     # if not, then $t3=1
bne  $t3, $zero, overflow
no_overflow:
...
...
overflow:
<Do something to handle overflow>
```

## The Basics of Exception Handling

MIPS uses two coprocessors: C0 and C1 for additional help. C0 primarily helps with **exception handling**, and C1 helps with **floating point** arithmetic.

### **Interrupts**

Initiated outside the instruction stream

Arrive asynchronously (at no specific time),

Example:

- I/O device status change
- I/O device error condition

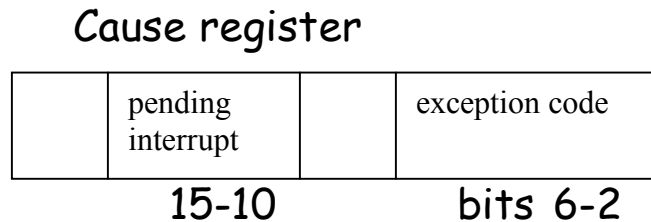
### **Traps**

Occur due to something in instruction stream.

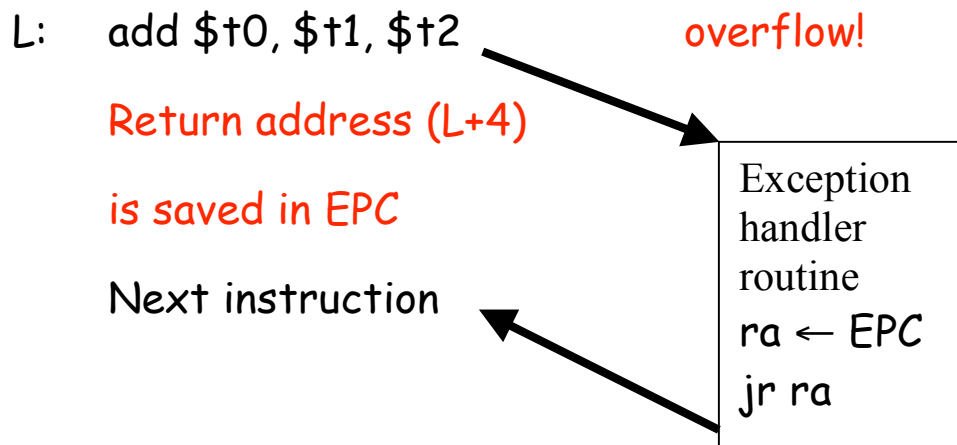
Examples:

- Unaligned address error
- Arithmetic overflow
- System call

**MIPS coprocessor C0** has a **cause register** (Register 13) that contains a 5-bit code to identify the cause of an **exception**



MIPS instructions that cause overflow (or some other violation) lead to an **exception**, which sets the **exception code**. It then switches to the **kernel mode** (designated by a bit in the **status register** of C0, register 12) and transfer control to a predefined address to invoke a routine (**exception handler**) for handling the exception.

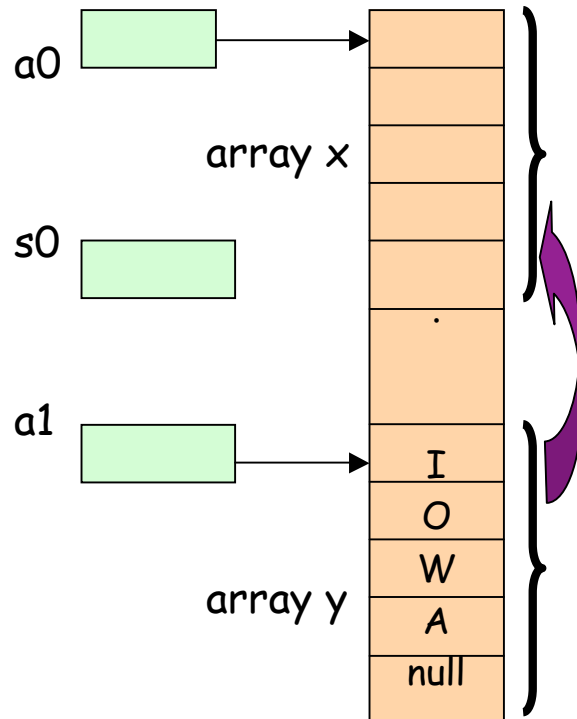


Exceptions cause **unscheduled procedure calls**.

## More Programming Examples

### Copying a string

Each char is represented by an ASCII byte. The string is terminated by a **Null** in ASCII). Reg s0 will hold the array index.



```
add $s0, $zero, $zero           # i = 0
L1: add $t1, $a1, $s0           # address of y[i] in t1
    lb $t2, 0($t1)             # t2 = y[i]
    add $t3, $a0, $s0         # address of x[i] in t3
    sb $t2, 0($t3)           # x[i] = y[i]
    addi $s0, $s0, 1         # i = i+1
    bne $t2, $zero, L1       # if y[i]≠0 then goto L
```

Load  
byte