## CS1210 Lecture 33 <br> Nov. 8, 2021

- HW 7 due
- DS 9 available at the end of class, due Friday 8pm
- HW 8 available after class, due Thursday next week
- you MAY NOT import any modules (such as Fraction) to help with HW8 Q1, the q1() function. Do the necessary basic math operations directly.
_ for buildWordGraph in DS9 and HW8, it should *not* take several minutes to build the graph for words5.txt. If it takes several minutes, you probably are using an $\mathrm{O}\left(\mathrm{n}^{\wedge} 3\right)$ algorithm rather than $O\left(n^{\wedge} 2\right)$, often because you use a linear time operation such as g.hasNode(...) inside the inner loop of your nested loops. (See slide 17 of this lecture)


## Last time

- Intro to optimization algorithms, greedy algorithms
- Introduced graphs (the computer science/mathematical kind), not the charts you've been plotting with pylab


## Today

- Graph representations
- Basic graph algorithms


## Graphs and optimization problems based on graphs

- Many important real-world problems can be modeled as optimization problems on graphs
- A graph is:
- A set of nodes (vertices)
- A set of edges (arcs) representing connections between pairs of nodes
- There are several types of graphs:
- Directed. Edges are "one way" from source to destination)
- Undirected. Edges have no particular direction - can travel either way, "see" each node from other, etc.
- Weighted. Edges have associated numbers called weights that can be used to represent cost, time, flow capacity, etc.
- See, e.g., Ch. 2 of free online book - Think Complexity - http:// greenteapress.com/complexity/html/thinkcomplexity003.html


Directed graph

- edges are "one way" from source to destination
- Can have two (one each way) between a pair of nodes
- Node can have edge to self
- Example relationships:
- course prerequisite
- hyperlink between web pages
- street between intersections
- Twitter follower
- Infection spread from-to



## Undirected graph

- Edges have no direction. Can "travel" either direction
- Can have only one edge between a pair of nodes*
- Node cannot have edge to self

- Example relationships:
- Facebook friend
- Bordering countries/states

*another kind of graph - multigraph -
relaxes this rule


## Weighted graphs

- Variant of both directed and undirected graphs in which each edge has an associate number called a weight or cost
- Edge weight provides additional information about the relationship between the nodes.
- Example relationships:
- Airfare between two cities
- Distance between two cities
- Flow capacity of oil/water pipeline between two points
- Network bandwidth between two ISP nodes



## Classic graph problems

- Determine if a graph has a cycle, a path that loops back to start points (e.g. 2-4-5-3-2)
- Find a path (non-branching) that traverses each (undirected) edge exactly once
- Leonhard Euler and the Bridges of Königsberg
- Not possible in graph on top right
- Find the shortest path between source $s$ and destination
- Different algorithms for weighted/unweighted graphs

- Find longest path between source and destination
- Find a path that visits each vertex exactly once
- A, E, D, C, F, B, A in example on bottom right
- Path of minimum cost that visits each vertex once
- A, E, D, C, B, F, A (cost 15) in example
- Assign no more than n different colors to vertices under constraint that no pair of connected vertices has the same color

Some of these are easy (have fast algorithms), others hard (no known efficient solution)


## Representing graphs

- How can we represent general graph in Python?
- Need to keep track of nodes
- Need to keep track of edges
- Several ways to represent graphs have been developed
- List of nodes and list of edges
- Adjacency matrix
- Adjacency lists
- Dictionary of dictionaries
- Efficiency of algorithms that solve graph problems can vary greatly depending on how graph are representated
- a strong influence on choice is the fact that one of the most common things needed in graph algorithms is access to immediate neighbors of a node (nodes that are destinations of edges for which "current" node is source)


## Adjacency matrix

|  |  | 1 | 2 | 3 | 4 | 5 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | False | True | True | True | False |
| - | 2 | True | False | False | False | False |
| - | 3 | True | False | False | True | False |
|  | 4 | True | False | True | False | False |
|  | 5 | False | False | False | False | False |

- Appealingly simple to understand and implement
- Use, e.g. a list of lists containing True/False, 0/1, or similar
- NOT the most common graph representation for most problems. Can you think of a reason why?
- Consider representing Facebook friends graph where each node is a FB user and an edge exists between two nodes whenever the two are FB friends.
- One billion nodes. Adjacency matrix 1B x 1B in size! Your computer doesn't have that much storage. But FB graph can be represented in computer! How?
- The 1B x 1B would be mostly False/0 - most people don't have huge number of friends. Should be representable in closer to 1 B * median number of friends. Other representations enable this huge memory savings.


## Adjacency list

## Use a dictionary with

- Nodes as keys
- Values are lists of neighbor nodes


Compared to adjacency matrix:

+ Much less space (when, as is common, most nodes have only a small relatively small number of neighbors). Facebook graph. People have hundreds of friends, not many milliions
- Query of "does edge ( $\mathrm{i}, \mathrm{j}$ ) exist?" not $\mathrm{O}(1)$. Must search list associated with node i to see if j is there. Turns out this is not crucial in many graph algorithms. (could address this using dictionary of dictionaries but often not necessary)


## Adjacency list graph representation

 Suitable for both undirected and directed graphs (and can be use for weighted graphs as well)

An adjacency list representation for undirected graphs in Python

## Two classes: Node and Graph

## basicgraph.py

Node

- properties:
- name: string
- status: string (we'll use this to "mark" nodes during traversals)
- methods
- getName
_ __repr__ : we'll print nodes as <name>

Note: in your HW8 you'll add one or more additional properties that help with traversing/walking through graphs to solve specific problems

## Adjacency list representation for undirected graphs

 Graph- properties
- nodes: a list of Node objects
basicgraph.py
- adjacencyLists: a dictionary with all nodes as keys. The value associated with a key n 1 (where n 1 is a node) is a list of all the nodes, n 2 , for which ( $\mathrm{n} 1, \mathrm{n} 2$ ) is an edge.
- methods
- addNode(node) : nodes must be added to graphs before edges
- addEdge(node1, node2) : presumes both nodes in graph already
- neighborsOf(node) : returns list of neighboring nodes
- getNode(name)
- hasNode(node)
- hasEdge(node1, node2) : return T if edge node1-node2 in graph
- __repr


$\longrightarrow$| KEY | VALUE |
| :---: | :--- |
| A |  |
| B |  |
| B $, \mathrm{C}, \mathrm{E}, \mathrm{H}]$ |  |
| $\mathrm{C}, \mathrm{C}, \mathrm{D}, \mathrm{F}]$ |  |
| $[\mathrm{A}, \mathrm{B}, \mathrm{F}]$ |  |
| D | $[\mathrm{B}]$ |
| E | $[\mathrm{A}, \mathrm{G}]$ |
| F | $[\mathrm{B}, \mathrm{C}]$ |
| G | $[\mathrm{E}, \mathrm{H}]$ |
| H | $[\mathrm{A}, \mathrm{G}]$ |

This graph is generated by genDemoGraph() in basicGraph.py

Note: for exams, you need to be able to 1) draw graph given adjacency list dictionary, and/or 2) show adjacency list dictionary given graph drawing



As I've said, many real-world problems can be represented as problems involving graphs. The algorithms to solve those problems often involve graph traversals, organized exploration or "walkthroughs" of the graph.
Two famous ones are: depth-first search and breadth-first search. I will present breadth-first search.
You will not be responsible for knowing the details of breadth-first search (for exam purposes) but you need to understand it well enough to use and extend it in HW8.

## Word ladder puzzles

CAT
???
???
DOG

CAT
COT
???
DOG

CAT
COT
DOT
DOG

Find 3-letter English words for ??? Positions. Each must differ from previous and next word in only one location

This problem is easily representable and solvable using graphs!

## DS9 buildWordGraph

Be careful in buildWordGraph - what is potentially slow about this?
for w1 in wordList:
$\mathrm{n} 1=$ g.getNode(w1)
for w2 in wordList:
$\mathrm{n} 2=$ g.getNode(w2)
If shouldHaveEdge( $\mathrm{w} 1, \mathrm{w} 2$ ):
g.addEdge(getNode(w1), getNode(w2))

Instead, recommend organizing as
for n 1 in g.nodes:
for n 2 in g.nodes:
If shouldHaveEdge(n1.getName(), n2.getName())

```
            g.addEdge(n1, n2)
```

Better yet:
for i in range (len(g.nodes)):
$\mathrm{n} 1=\operatorname{g} . \operatorname{nodes}[1]$
for j in range( i , len(g.nodes)):
n2 $=$ g.nodes[j]

This has two problems:

1) error because
2) very slow because ?
<- : You could address problem 1 here with and 'if ...'

Note: fixes both problems

## Next time

- Graph traversals
- Breadth first search
- Depth first search
- HW 8

