CS:4420 Artificial Intelligence

Spring 2018

Propositional Logic

Cesare Tinelli

The University of Iowa

Readings

- Chap. 7 of [Russell and Norvig, 2012]

Logics

A logic is a triple $\langle \mathcal{L}, \mathcal{S}, \mathcal{R} \rangle$ where

- $\mathcal{L}$, the logic's language, is a class of sentences described by a formal grammar
- $\mathcal{S}$, the logic's semantics is a formal specification of how to assign *meaning* in the "real world" to the elements of $\mathcal{L}$
- $\mathcal{R}$, the logic's inference system, is a set of formal derivation *rules* over $\mathcal{L}$

There are several logics: propositional, first-order, higher-order, modal, temporal, intuitionistic, linear, equational, non-monotonic, fuzzy, ...

We will concentrate on propositional logic and first-order logic

Propositional Logic

Each sentence is made of

- propositional variables ($A, B, \ldots, P, Q, \ldots$)
- logical constants ($\mathbf{True}, \mathbf{False}$)
- logical connectives ($\wedge, \vee, \Rightarrow, \ldots$)

Every propositional variable stands for a basic fact

Examples: *I'm hungry, Apples are red, Joe and Jill are married*

Propositional Logic

Ontological Commitments

Propositional Logic is about facts in the world that are either true or false, nothing else

Semantics of Propositional Logic

Since each propositional variable stands for a fact about the world, its meaning ranges over the Boolean values $\{true, false\}$

Note: Do note confuse

- $true$, $false$, which are values (i.e., semantical entities) here with
- $\mathbf{True}$, $\mathbf{False}$, which are logical constants (i.e., symbols of the language)

Propositional Logic

The Language

- Each propositional variable ($A, B, \ldots, P, Q, \ldots$) is a sentence
- Each logical constant ($\mathbf{True}, \mathbf{False}$) is a sentence
- If φ and ψ are sentences, all of the following are also sentences

$$(\varphi) \qquad \neg\varphi \qquad \varphi \wedge \psi \qquad \varphi \vee \psi \qquad \varphi \Rightarrow \psi \qquad \varphi \Leftrightarrow \psi$$

- Nothing else is a sentence

The Language of Propositional Logic

More formally, it is the language generated by the following grammar

Symbols:

- Propositional variables: $A, B, \ldots, P, Q, \ldots$
- Logical constants:

$\mathbf{True}$	(true)	$\wedge$	(and)	$\Rightarrow$	(implies)	$\neg$	(not)
$\mathbf{False}$	(false)	$\vee$	(or)	$\Leftrightarrow$	(equivalent)		

Grammar Rules:

$$
\begin{array}{lcl}
\mathit{Sentence} & ::= & \mathit{AtomicS} \mid \mathit{ComplexS} \\
\mathit{AtomicS} & ::= & \mathbf{True} \mid \mathbf{False} \mid A \mid B \mid \ \ldots \ \mid P \mid Q \mid \ \ldots \\
\mathit{ComplexS} & ::= & (\mathit{Sentence}) \mid \mathit{Sentence}\ \mathit{Connective}\ \mathit{Sentence} \mid \neg \mathit{Sentence} \\
\mathit{Connective} & ::= & \wedge \mid \ \vee \ \mid \ \Rightarrow \ \mid \ \Leftrightarrow
\end{array}
$$

Wumpus world sentences

Let $P_{i,j}$ denote that there is a pit in position (i, j)
Let $B_{i,j}$ denote that there is a breeze in position (i, j)

"There is no pit in the initial position but there is one in $(2, 2)$"

$$\neg P_{1,1} \land P_{2,2}$$

Wumpus world sentences

Let $P_{i,j}$ denote that there is a pit in position (i, j)
Let $B_{i,j}$ denote that there is a breeze in position (i, j)

"There is no pit in the initial position but there is one in $(2, 2)$"

$$\neg P_{1,1} \land P_{2,2}$$

"A square is breezy if and only if there is an adjacent pit"

$$\begin{array}{ccc} B_{1,1} & \Leftrightarrow & (P_{2,1} \lor P_{1,2}) \\ B_{2,2} & \Leftrightarrow & (P_{2,1} \lor P_{3,2} \lor P_{2,1} \lor P_{1,2}) \\ \vdots & \vdots & \vdots \end{array}$$

Semantics of Propositional Logic

The meaning of **True** is always *true*
The meaning of **False** is always *false*

The meaning of the other sentences depends on the meaning of the propositional variables

- Base cases: truth tables

P	Q	$\neg P$	$P \wedge Q$	$P \vee Q$	$P \Rightarrow Q$	$P \Leftrightarrow Q$
false	*false*	*true*	*false*	*false*	*true*	*true*
false	*true*	*true*	*false*	*true*	*true*	*false*
true	*false*	*false*	*false*	*true*	*false*	*false*
true	*true*	*false*	*true*	*true*	*true*	*true*

- Non-base Cases: given by reduction to the base cases

  Ex: the meaning of $(P \vee Q) \wedge R$ is the same as the meaning of $A \wedge R$ where A has the same meaning as $P \vee Q$

Semantics of Propositional Logic

An assignment of Boolean values to the propositional variables of a sentence is an interpretation of the sentence

P	H	$P \lor H$	$(P \lor H) \land \neg H$	$((P \lor H) \land \neg H) \Rightarrow P$
False	*False*	*False*	*False*	*True*
False	*True*	*True*	*False*	*True*
True	*False*	*True*	*True*	*True*
True	*True*	*True*	*False*	*True*

Interpretations: $\{P \mapsto \mathit{false}, H \mapsto \mathit{false}\}, \{P \mapsto \mathit{false}, H \mapsto \mathit{true}\}, \ldots$

An interpretation is a model of a sentence φ if it makes the sentence true

Note: The semantics of Propositional Logic is compositional — the meaning of a sentence is defined recursively in terms of the meaning of the sentence's components

Semantics of Propositional Logic

The meaning of a sentence in general depends on its interpretation

Some sentences, however, have always the same meaning

P	H	$P \vee H$	$(P \vee H) \wedge \neg H$	$((P \vee H) \wedge \neg H) \Rightarrow P$
False	*False*	*False*	*False*	*True*
False	*True*	*True*	*False*	*True*
True	*False*	*True*	*True*	*True*
True	*True*	*True*	*False*	*True*

A sentence is

- satisfiable if it is true in some interpretation
- unsatisfiable if it is true in no interpretation
- valid if it is true in every possible interpretation
- invalid if it is false in some possible interpretation

A Warning

Disjunction

- $A \vee B$ is true when A or B or *or both* are true (inclusive or)
- $A \oplus B$ is sometimes used to mean "either A or B but not both" (exclusive or)

Implication

- $A \Rightarrow B$ does not require a causal connection between A and B
  Ex: *Sky-is-blue* $\Rightarrow$ *Snow-is-white*
- When A is false, $A \Rightarrow B$ is always true *regardless* of B
  Ex: *Two-equals-four* $\Rightarrow$ *Apples-are-red*
- Beware of negations in implications
  Ex: *Is-a-female-bird* $\Rightarrow$ *Lays-eggs*
  $\neg$*Is-a-female-bird* $\Rightarrow$ $\neg$*lays-eggs*

Entailment in Propositional Logic

Given

- a set Γ of sentences and
- a sentence φ,

we write

$$\Gamma \models \varphi$$

iff every interpretation that makes all sentences in Γ true makes φ also true

$\Gamma \models \varphi$ is read as "Γ entails φ" or "φ logically follows from Γ"

Entailment in Propositional Logic

Examples

$$
\begin{array}{lcl}
\{A, A \Rightarrow B\} & \models & B \\
\{A\} & \models & A \vee B \\
\{A, B\} & \models & A \wedge B \\
\{\} & \models & A \vee \neg A \\
\{A\} & \not\models & A \wedge B \\
\{A \vee \neg A\} & \not\models & A
\end{array}
$$

	A	B	$A \Rightarrow B$	$A \vee B$	$A \wedge B$	$A \vee \neg A$
1.	*false*	*false*	*true*	*false*	*false*	*true*
2.	*false*	*true*	*true*	*true*	*false*	*true*
3.	*true*	*false*	*false*	*true*	*false*	*true*
4.	*true*	*true*	*true*	*true*	*true*	*true*

Properties of Entailment

- $\Gamma \models \varphi$, for all $\varphi \in \Gamma$ (inclusion property of PL)
- if $\Gamma \models \varphi$, then $\Gamma' \models \varphi$ for all $\Gamma' \supseteq \Gamma$ (monotonicity of PL)
- φ is valid iff $\{\,\} \models \varphi$ (also written as $\models \varphi$)
- φ is unsatisfiable iff $\varphi \models \mathbf{False}$
- $\Gamma \models \varphi$ iff the set $\Gamma \cup \{\neg\varphi\}$ is unsatisfiable

Logical Equivalence

Two sentences φ_1 and φ_2 are logically equivalent, written

$$\varphi_1 \equiv \varphi_2$$

if $\varphi_1 \models \varphi_2$ and $\varphi_2 \models \varphi_1$

Note:

- $\varphi_1 \equiv \varphi_2$ if and only if every interpretation assigns the same Boolean value to φ_1 and φ_2
- Implication and equivalence ($\Rightarrow$, $\Leftrightarrow$), which are syntactical entities, are intimately related to entailment and logical equivalence ($\models$, $\equiv$), which are semantical notions:

$$\begin{array}{lll} \varphi_1 \models \varphi_2 & \text{iff} & \models \varphi_1 \Rightarrow \varphi_2 \\ \varphi_1 \equiv \varphi_2 & \text{iff} & \models \varphi_1 \Leftrightarrow \varphi_2 \end{array}$$

Properties of Logical Connectives

- $\wedge$ and $\vee$ are commutative

$$\begin{aligned} \varphi_1 \wedge \varphi_2 &\equiv \varphi_2 \wedge \varphi_1 \\ \varphi_1 \vee \varphi_2 &\equiv \varphi_2 \vee \varphi_1 \end{aligned}$$

- $\wedge$ and $\vee$ are *associative*

$$\begin{aligned} \varphi_1 \wedge (\varphi_2 \wedge \varphi_3) &\equiv (\varphi_1 \wedge \varphi_2) \wedge \varphi_3 \\ \varphi_1 \vee (\varphi_2 \vee \varphi_3) &\equiv (\varphi_1 \vee \varphi_2) \vee \varphi_3 \end{aligned}$$

- $\wedge$ and $\vee$ are mutually *distributive*

$$\begin{aligned} \varphi_1 \wedge (\varphi_2 \vee \varphi_3) &\equiv (\varphi_1 \wedge \varphi_2) \vee (\varphi_1 \wedge \varphi_3) \\ \varphi_1 \vee (\varphi_2 \wedge \varphi_3) &\equiv (\varphi_1 \vee \varphi_2) \wedge (\varphi_1 \vee \varphi_3) \end{aligned}$$

- $\wedge$ and $\vee$ are related by $\neg$ (DeMorgan's Laws)

$$\begin{aligned} \neg(\varphi_1 \wedge \varphi_2) &\equiv \neg\varphi_1 \vee \neg\varphi_2 \\ \neg(\varphi_1 \vee \varphi_2) &\equiv \neg\varphi_1 \wedge \neg\varphi_2 \end{aligned}$$

Properties of Logical Connectives

$\wedge$, $\Rightarrow$, and $\Leftrightarrow$ are actually redundant:

$$\begin{array}{lcl} \varphi_1 \wedge \varphi_2 & \equiv & \neg(\neg\varphi_1 \vee \neg\varphi_2) \\ \varphi_1 \Rightarrow \varphi_2 & \equiv & \neg\varphi_1 \vee \varphi_2 \\ \varphi_1 \Leftrightarrow \varphi_2 & \equiv & (\varphi_1 \Rightarrow \varphi_2) \wedge (\varphi_2 \Rightarrow \varphi_1) \end{array}$$

We keep them all mainly for convenience

Exercise Use the truth tables to verify all the logical equivalences seen so far

Inference Systems for Propositional Logic

An inference system $\mathcal{I}$ for PL is a procedure that

given a set $\Gamma = \{\alpha_1, \ldots, \alpha_m\}$ of sentences and a sentence φ, may reply "yes", "no", or run forever

If $\mathcal{I}$ replies positively to input (Γ, φ), we say that Γ derives φ in $\mathcal{I}$ (or, $\mathcal{I}$ derives φ from Γ, or, φ derives from Γ in $\mathcal{I}$) and write

$$\Gamma \vdash_{\mathcal{I}} \varphi$$

Intuitively, $\mathcal{I}$ should be such that it replies "yes" on input (Γ, φ) only if φ is in fact entailed by Γ

All These Fancy Symbols!

Note:

$A \wedge B \Rightarrow C$ is a sentence, a bunch of symbols manipulated by an inference system $\mathcal{I}$

$A \wedge B \models C$ is a mathematical abbreviation standing for the statement: "every interpretation that makes $A \wedge B$ true, makes C also true"

$A \wedge B \vdash_{\mathcal{I}} C$ is a mathematical abbreviation standing for the statement: "$\mathcal{I}$ derives C from $A \wedge B$"

In other words,

$\Rightarrow$ is a formal symbol of the logic, which is used by the inference system

$\models$ is a shorthand we use to talk about the meaning of formal sentences

$\vdash_{\mathcal{I}}$ is a shorthand we use to talk about the output of the inference system $\mathcal{I}$

All These Fancy Symbols!

The connective $\Rightarrow$ and the shorthand $\models$ are related as follows

The sentence $\varphi_1 \Rightarrow \varphi_2$ is valid (always true) if and only if $\varphi_1 \models \varphi_2$

Example: $A \Rightarrow (A \lor B)$ is valid and $A \models (A \lor B)$

	A	B	$A \lor B$	$A \Rightarrow (A \lor B)$
1.	*false*	*false*	*false*	*true*
2.	*false*	*true*	*true*	*true*
3.	$\underline{\textit{true}}$	*false*	*true*	*true*
4.	$\underline{\textit{true}}$	*true*	*true*	*true*

All These Fancy Symbols!

The shorthands $\models$ and $\vdash_{\mathcal{I}}$ are related as follows

- A sound inference system can derive only sentences that logically follow from a given set of sentences:

  $$\text{if } \ \Gamma \vdash_{\mathcal{I}} \varphi \ \text{ then } \ \Gamma \models \varphi$$

- A complete inference system can derive all sentences that logically follow from a given set of sentences:

  $$\text{if } \ \Gamma \models \varphi \ \text{ then } \ \Gamma \vdash_{\mathcal{I}} \varphi$$

Inference systems for PL

Divided into (roughly) two kinds:

Rule-based

- Sound generation of new sentences from old
- Proof = a sequence of inference rule applications
  Can use inference rules as operators as in a standard search procedures
- Typically require translation of sentences into some normal form

Model-based

- Truth table enumeration (always exponential in n)
- Improved backtracking, e.g., Davis–Putnam–Logemann–Loveland
- Heuristic search in model space (incomplete)
  e.g., min-conflicts-like hill-climbing algorithms

Truth table enumeration

The proof system $\mathcal{TT}$ is specified as follows:

$\{\alpha_1, \ldots, \alpha_m\} \vdash_{\mathcal{TT}} \varphi$ *iff* *all the values in the truth table of* $(\alpha_1 \wedge \cdots \wedge \alpha_m) \Rightarrow \varphi$ *are* $true$

Inference by Truth Tables

- The truth-tables-based inference system is sound:

$\alpha_1, \ldots, \alpha_m \vdash_{\mathcal{TT}} \varphi$	implies	truth table of $(\alpha_1 \wedge \cdots \wedge \alpha_m) \Rightarrow \varphi$ all true
	implies	$(\alpha_1 \wedge \cdots \wedge \alpha_m) \Rightarrow \varphi$ is valid
	implies	$\models (\alpha_1 \wedge \cdots \wedge \alpha_m) \Rightarrow \varphi$
	implies	$(\alpha_1 \wedge \cdots \wedge \alpha_m) \models \varphi$
	implies	$\alpha_1, \ldots, \alpha_m \models \varphi$

- It is also complete (exercise: prove it)
- Its time complexity is $O(2^n)$ where n is the number of propositional variables in $\alpha_1, \ldots, \alpha_m, \varphi$
- We cannot hope to do better in general because the dual problem: determining the satisfiability of a sentence, is NP-complete
- However, really hard cases of propositional inference are somewhat rare in practice

Rule-Based Inference in PL

An inference system in Propositional Logic can also be specified as a set $\mathcal{R}$ of inference (or derivation) rules

- Each rule is just a *pattern* premises/conclusion
- A rule applies to Γ and derives φ if
  - some of the sentences in Γ match with the premises of the rule and
  - φ matches with the conclusion
- A rule is sound it the set of its premises entails its conclusion

Some Inference Rules

- And-Introduction

$$\frac{\alpha \qquad \beta}{\alpha \wedge \beta}$$

- And-Elimination

$$\frac{\alpha \wedge \beta}{\alpha} \qquad\qquad \frac{\alpha \wedge \beta}{\beta}$$

- Or-Introduction

$$\frac{\alpha}{\alpha \vee \beta} \qquad\qquad \frac{\alpha}{\beta \vee \alpha}$$

Some Inference Rules (cont'd)

- Implication-Elimination (aka Modus Ponens)

$$\frac{\alpha \Rightarrow \beta \qquad \alpha}{\beta}$$

- Unit Resolution

$$\frac{\alpha \vee \beta \qquad \neg\beta}{\alpha}$$

- Resolution

$$\frac{\alpha \vee \beta \qquad \neg\beta \vee \gamma}{\alpha \vee \gamma}$$ or, equivalently,

$$\frac{\neg\alpha \Rightarrow \beta, \qquad \beta \Rightarrow \gamma}{\neg\alpha \Rightarrow \gamma}$$

Some Inference Rules (cont'd)

- Double-Negation-Elimination

$$\frac{\neg\neg\alpha}{\alpha}$$

- **False**-Introduction

$$\frac{\alpha \quad \neg\alpha}{\mathbf{False}}$$

- **False**-Elimination

$$\frac{\mathbf{False}}{\beta}$$

Inference by Proof

We say there is a proof of φ from Γ in an inference system $\mathcal{R}$ if we can derive φ by applying the rules of $\mathcal{R}$ repeatedly to Γ and its derived sentences

Example: a proof of P *from* $\{(P \vee H) \wedge \neg H\}$

1.	$(P \vee H) \wedge \neg H$	*by assumption*
2.	$P \vee H$	*by* $\wedge$*-elimination applied to (1)*
3.	$\neg H$	*by* $\wedge$*-elimination applied to (1)*
4.	P	*by unit resolution applied to (2),(3)*

We can represent a proof more visually as a proof tree:

Example:

$$\dfrac{\dfrac{(P \vee H) \wedge \neg H}{P \vee H} \qquad \dfrac{(P \vee H) \wedge \neg H}{\neg H}}{P}$$

Rule-Based Inference in Propositional Logic

More precisely, there is a proof of φ from Γ in $\mathcal{R}$ if

1. $\varphi \in \Gamma$ or,
2. there is a rule in $\mathcal{R}$ that applies to Γ and produces φ or,
3. there is a proof of each $\varphi_1, \ldots, \varphi_m$ from Γ in $\mathcal{R}$ and a rule that applies to $\{\varphi_1, \ldots, \varphi_m\}$ and produces φ

Then, the inference system $\mathcal{R}$ is specified as follows:

$\Gamma \vdash_{\mathcal{R}} \varphi$ *iff there is a proof of* φ *from* Γ *in* $\mathcal{R}$

An Inference System $\mathcal{R}$

$$\frac{\alpha \qquad \beta}{\alpha \wedge \beta} \qquad \frac{\alpha}{\alpha \vee \beta} \qquad \frac{\alpha}{\beta \vee \alpha}$$

$$\frac{\alpha \wedge \beta}{\alpha} \qquad \frac{\alpha \wedge \beta}{\beta}$$

$$\frac{\alpha \Rightarrow \beta \qquad \alpha}{\beta} \qquad \frac{\alpha \vee \beta \qquad \neg\beta}{\alpha} \qquad \frac{\alpha \vee \beta \qquad \neg\beta \vee \gamma}{\alpha \vee \gamma}$$

$$\frac{\neg\neg\alpha}{\alpha} \qquad \frac{\alpha \quad \neg\alpha}{\textbf{False}} \qquad \frac{\textbf{False}}{\beta}$$

Soundness of $\mathcal{R}$

The given system $\mathcal{R}$ is sound because all of its rules are

Example: the Resolution rule $\dfrac{\alpha \vee \beta, \quad \neg\beta \vee \gamma}{\alpha \vee \gamma}$

	α	β	γ	$\neg\beta$	$\alpha \vee \beta$	$\neg\beta \vee \gamma$	$\alpha \vee \gamma$
1.	*false*	*false*	*false*	*true*	*false*	*true*	*false*
2.	*false*	*false*	*true*	*true*	*false*	*true*	*true*
3.	*false*	*true*	*false*	*false*	*true*	*false*	*false*
4.	*false*	*true*	*true*	*false*	$\underline{true}$	$\underline{true}$	*true*
5.	*true*	*false*	*false*	*true*	$\underline{true}$	$\underline{true}$	*true*
6.	*true*	*false*	*true*	*true*	$\underline{true}$	$\underline{true}$	*true*
7.	*true*	*true*	*false*	*false*	*true*	*false*	*true*
8.	*true*	*true*	*true*	*false*	$\underline{true}$	$\underline{true}$	*true*

All the interpretations that satisfy both $\alpha \vee \beta$ and $\neg\beta \vee \gamma$ (4,5,6,8) satisfy $\alpha \vee \gamma$ as well

Soundness of $\mathcal{R}$

The given system $\mathcal{R}$ is sound because all of its rules are

Exercise: prove that the other inference rules are sound as well

Is $\mathcal{R}$ also complete?

Resolution

Literal: prop. symbol (P) or negated prop. symbol ($\neg P$)

Clause: set of literals $\{l_1, \ldots, l_k\}$ (understood as $l_1 \vee \cdots \vee l_k$)

Conjunctive Normal Form: set of clauses $\{C_1, \ldots, C_n\}$ (understood as $C_1 \wedge \cdots \wedge C_n$)

Resolution rule for CNF:

$$\frac{l_1 \vee \cdots \vee l_k \vee P \qquad \neg P \vee m_1 \vee \cdots \vee m_n}{l_1 \vee \cdots \vee l_k \vee m_1 \vee \cdots \vee m_n}$$

E.g.,

$$\frac{P \vee Q \qquad \neg Q}{P} \qquad \frac{P \vee Q \qquad R \vee \neg Q \vee \neg S}{P \vee R \vee \neg S} \qquad \frac{P \vee Q \qquad \neg Q \vee P \vee R}{P \vee R}$$

Resolution is sound and complete for CNF KBs

Conversion to CNF

Ex.: $A \Leftrightarrow (B \vee C)$

1. Eliminate $\Leftrightarrow$, replacing $\alpha \Leftrightarrow \beta$ with $(\alpha \Rightarrow \beta) \wedge (\beta \Rightarrow \alpha)$

$$(A \Rightarrow (B \vee C)) \wedge ((B \vee C) \Rightarrow A)$$

2. Eliminate $\Rightarrow$, replacing $\alpha \Rightarrow \beta$ with $\neg\alpha \vee \beta$

$$(\neg A \vee B \vee C) \wedge (\neg(B \vee C) \vee A)$$

3. Move $\neg$ inwards using de Morgan's rules and double-negation

$$(\neg A \vee B \vee C) \wedge ((\neg B \wedge \neg C) \vee A)$$

4. Apply distributivity law ($\vee$ over $\wedge$) and flatten

$$(\neg A \vee B \vee C) \wedge (\neg B \vee A) \wedge (\neg C \vee A)$$

Resolution Procedure

Proof by contradiction: show $KB \models \alpha$ by showing $KB \wedge \neg\alpha$ unsatisfiable.

Do the latter by deriving **False** from CNF of $KB \wedge \neg\alpha$

function PL-RESOLUTION(KB, α) **returns** *true* or *false*

$clauses \leftarrow$ the set of clauses in the CNF representation of $KB \wedge \neg\alpha$

$new \leftarrow \{\,\}$

loop do

  for each C_i, C_j **in** *clauses* **do**

    $resolvents \leftarrow$ PL-RESOLVE(C_i, C_j)

    if *resolvents* contains the empty clause **then return** *true*

    $new \leftarrow new \cup resolvents$

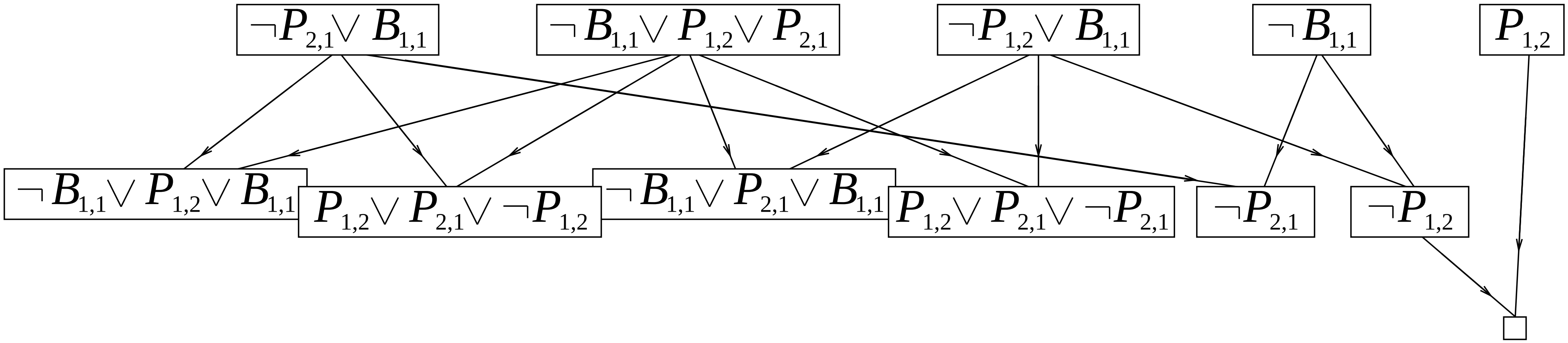
  if $new \subseteq clauses$ **then return** *false*

  $clauses \leftarrow clauses \cup new$

Resolution Example

$$KB = \{ B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1}),\ \neg B_{1,1} \}$$
$$\alpha = \neg P_{1,2}$$

$$\mathsf{CNF} = \{ \neg P_{1,2} \vee B_{1,1},\ \neg B_{1,1} \vee P_{1,2} \vee P_{2,1},\ \neg P_{1,2} \vee B_{1,1},\ \neg B_{1,1},\ P_{1,2} \}$$



Forward and backward chaining

Horn clause: prop. symbol (p) or implication $p_1 \wedge \cdots \wedge p_n \Rightarrow p$

Horn Form: set of Horn clauses $\{C_1, \ldots, C_n\}$ (understood as $C_1 \wedge \cdots \wedge C_n$)

E.g., $\{\ C,\ B \Rightarrow A,\ C \wedge D \Rightarrow B\ \}$

Modus Ponens for Horn Form

$$\frac{\alpha_1 \quad \cdots \quad \alpha_n \qquad \alpha_1 \wedge \cdots \wedge \alpha_n \Rightarrow \beta}{\beta}$$

Sound and complete for Horn Form KBs

Can be used with forward chaining or backward chaining

These algorithms are very natural and run in linear time

Forward chaining

Idea:
Fire any rule whose premises are satisfied in the KB,
add its conclusion to the KB, until query is found

Ex.: query is Q

KB is

$P \implies Q$

$L \wedge M \implies P$

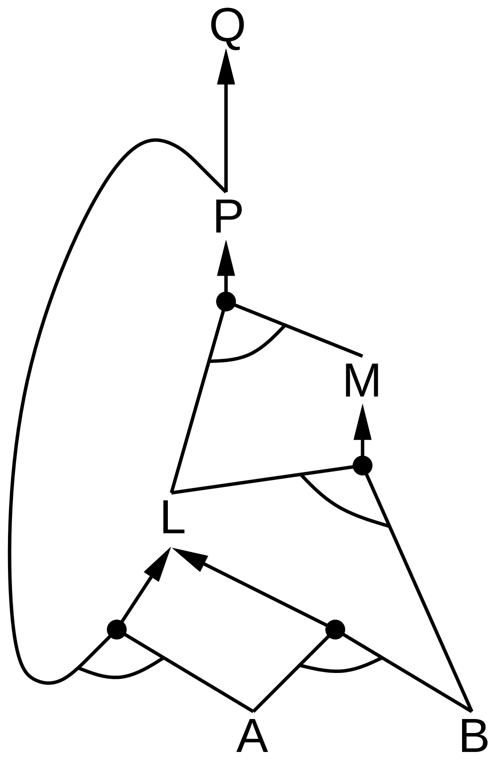
$B \wedge L \implies M$

$A \wedge P \implies L$

$A \wedge B \implies L$

A

B



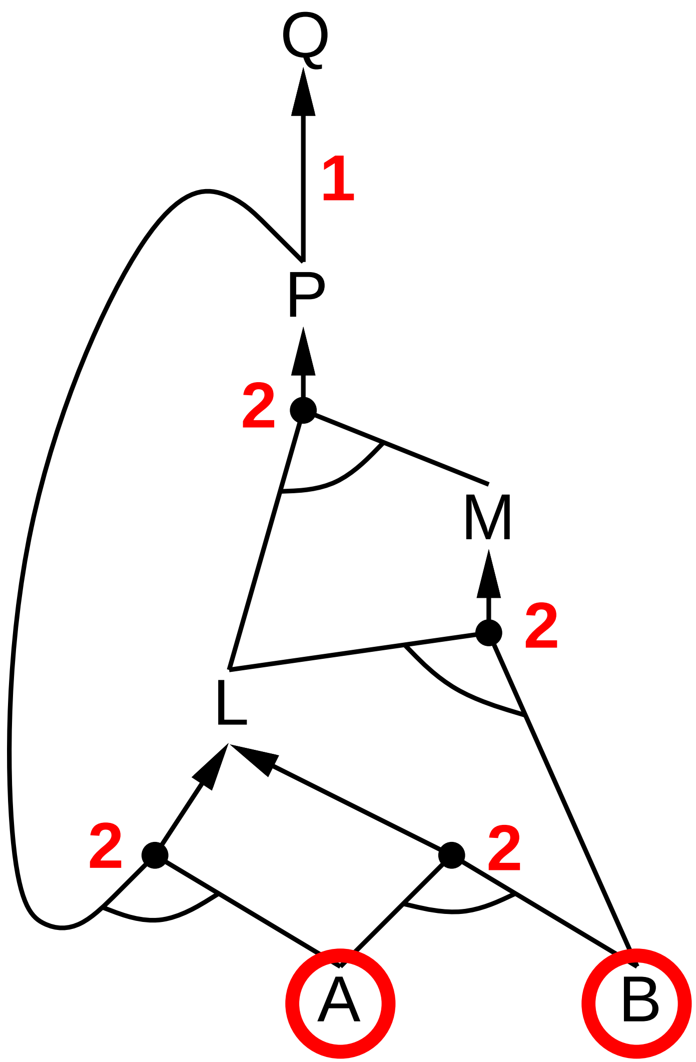
Forward chaining algorithm

function PL-FC-ENTAILS?(KB, q) **returns** *true* or *false*

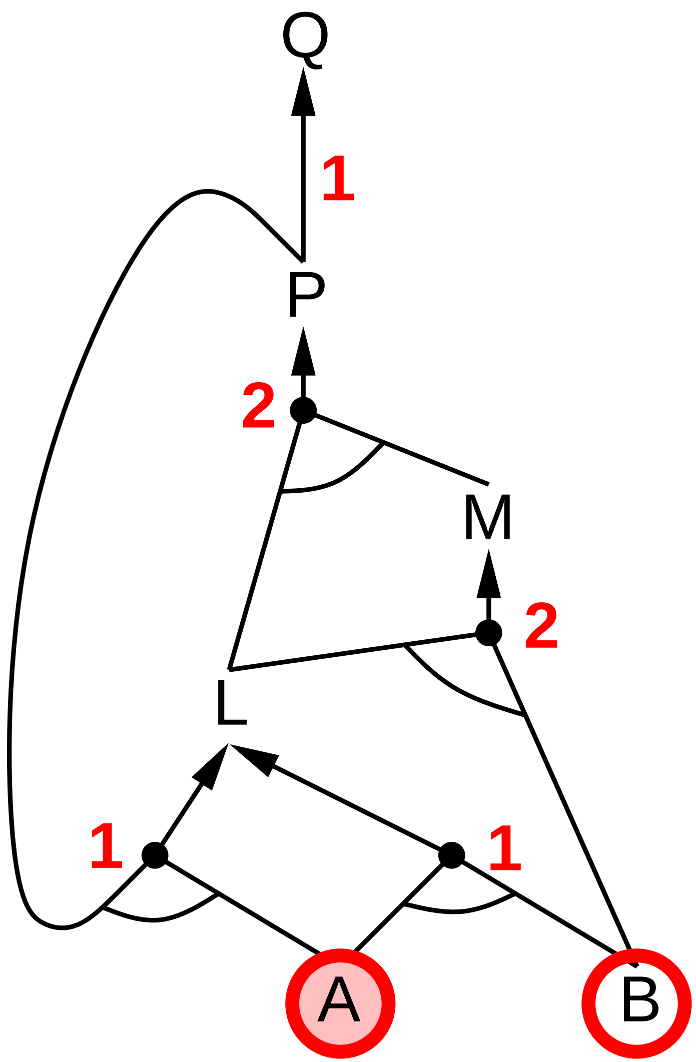
count, a table, indexed by clause, initially the number of premises
inferred, a table, indexed by symbol, each entry initially *false*
agenda, a list of symbols, initially the symbols known to be true

```
while agenda is not empty do
    p ← Pop(agenda)
    unless inferred[p] do
        inferred[p] ← true
        for each Horn clause c in whose premise p appears do
            decrement count[c]
            if count[c] = 0 then do
                if Head[c] = q then return true
                Push(Head[c], agenda)
return false
```

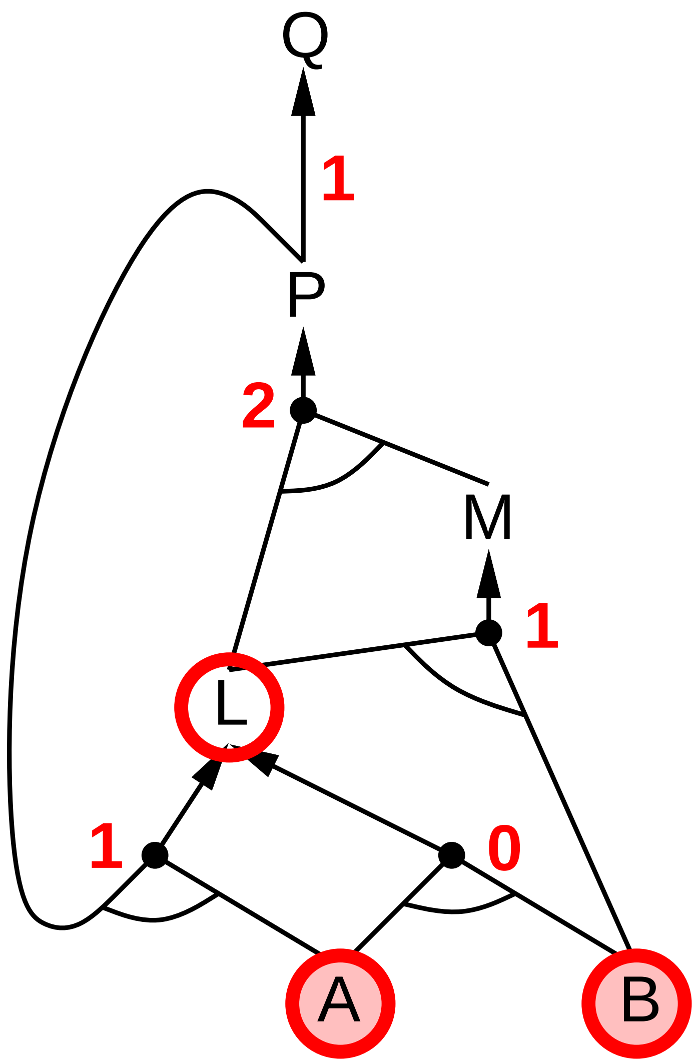
Forward chaining example



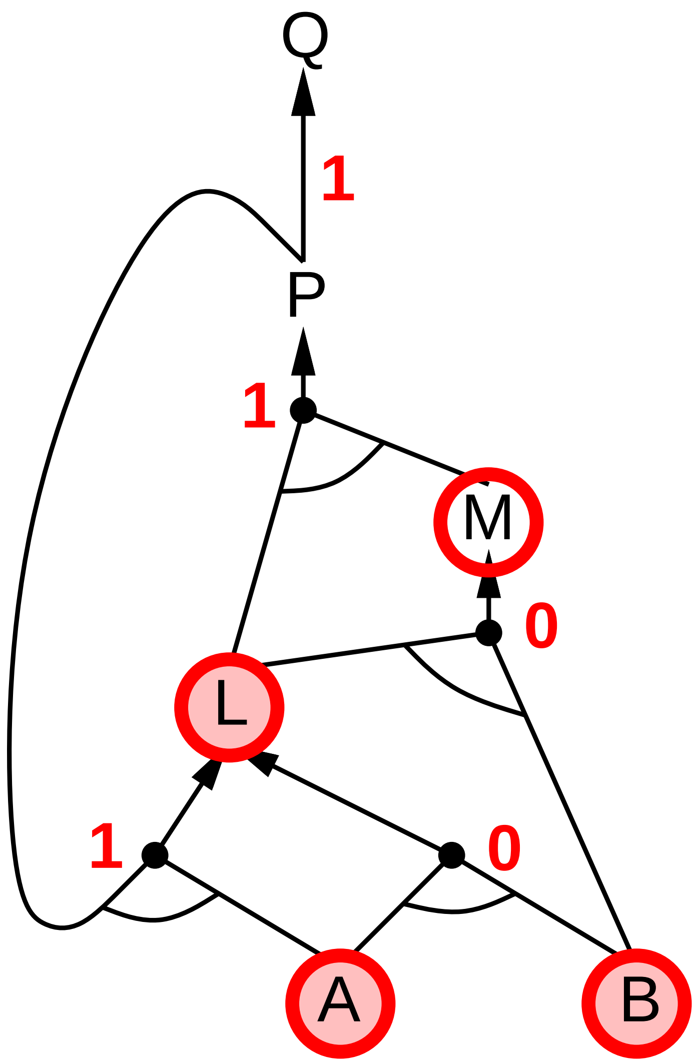
Forward chaining example



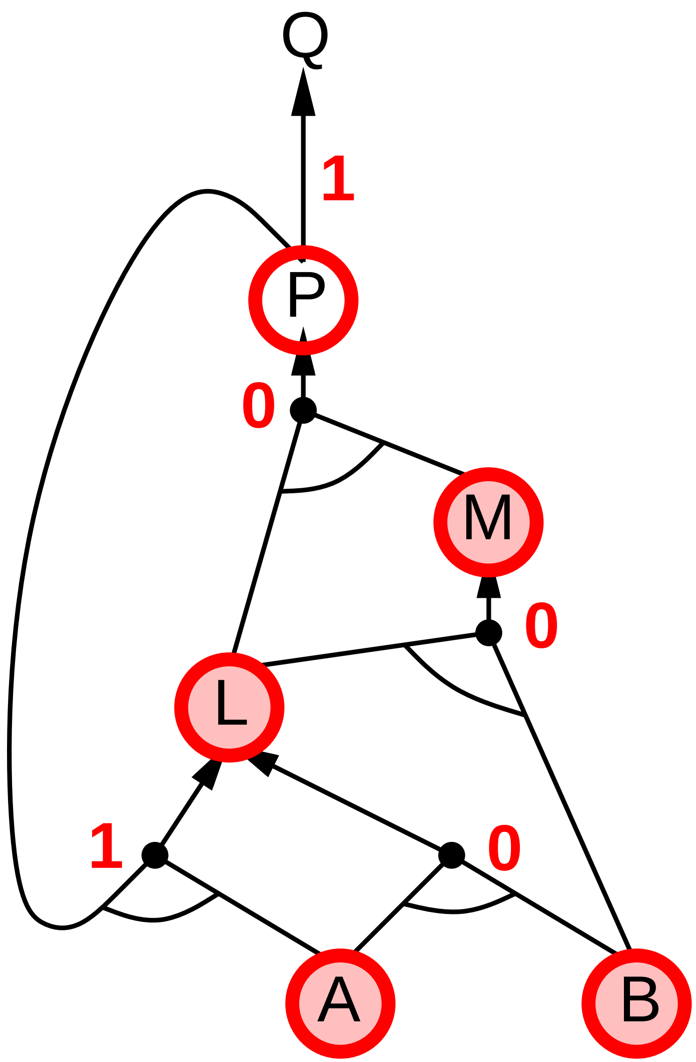
Forward chaining example



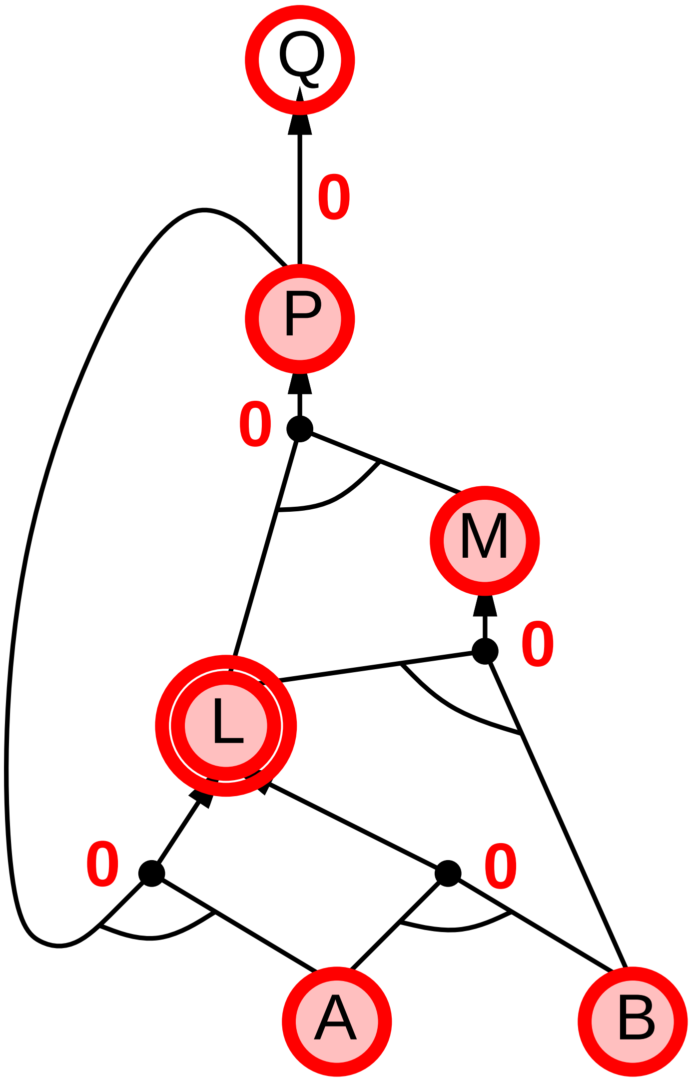
Forward chaining example



Forward chaining example



Forward chaining example



Proof of completeness (sketch)

FC derives every atomic sentence that is entailed by KB

1. FC reaches a fixed point where no new atomic sentences (prop. symbols) are inferred

2. Consider the final state as an interpretation m, assigning true to the inferred symbols and false to the other symbols

3. Claim: Every clause in the original KB is satisfied by m

   Proof: Suppose a clause $a_1 \wedge \ldots \wedge a_k \Rightarrow b$ is falsified by m
   Then $a_1 \wedge \ldots \wedge a_k$ is satisfied by m while b is not
   But then b was not inferred, contradicting the assumption that the algorithm had reached a fixed point!

4. Hence m is a model of KB

5. 5. If $KB \models q$, q is true in every model of KB, including m

Backward chaining

Idea: work backwards from the query q

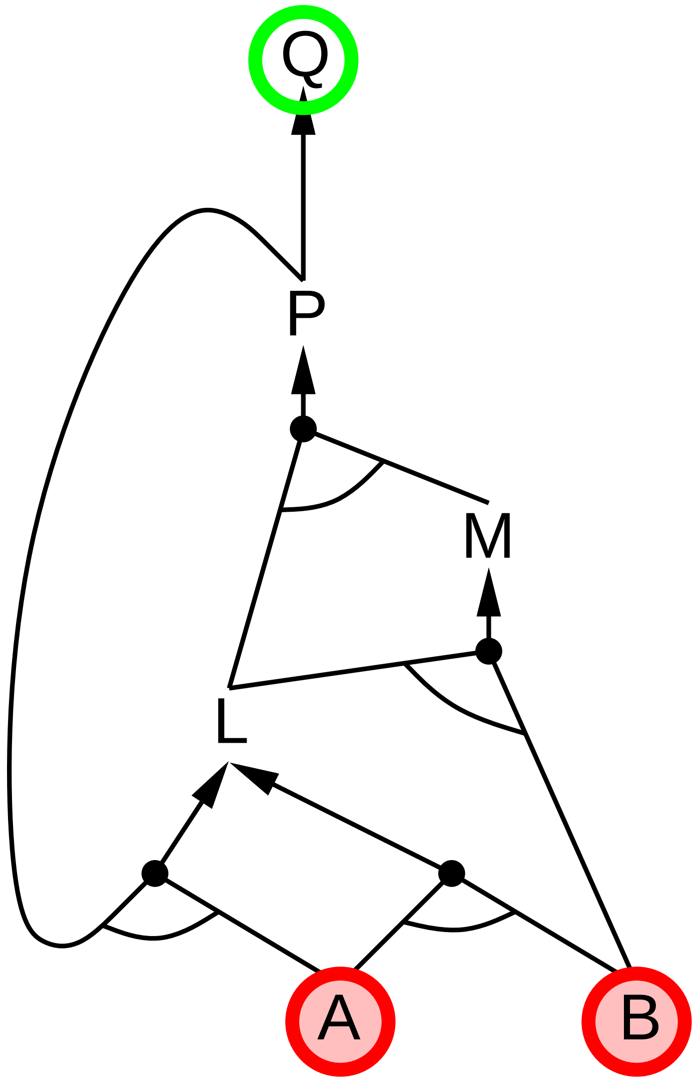
to infer q by BC,
check if q is known already, or
infer by BC all premises of some rule concluding q

Avoid loops: check if new subgoal is already on the goal stack

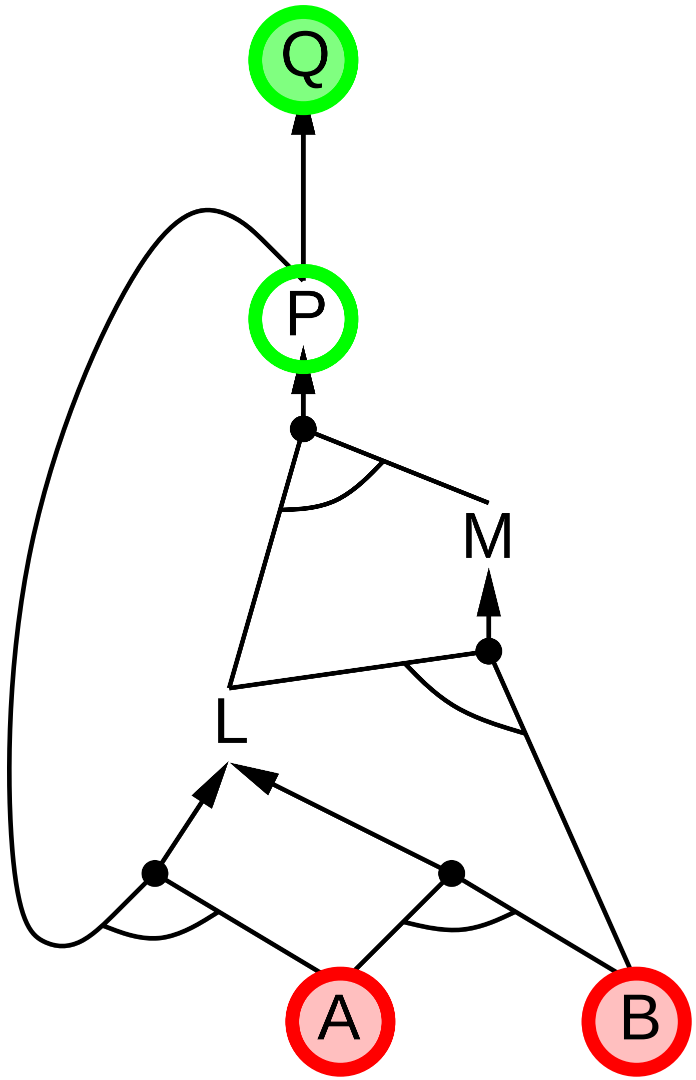
Avoid repeated work: check if new subgoal

1. has already been inferred, or
2. has already failed

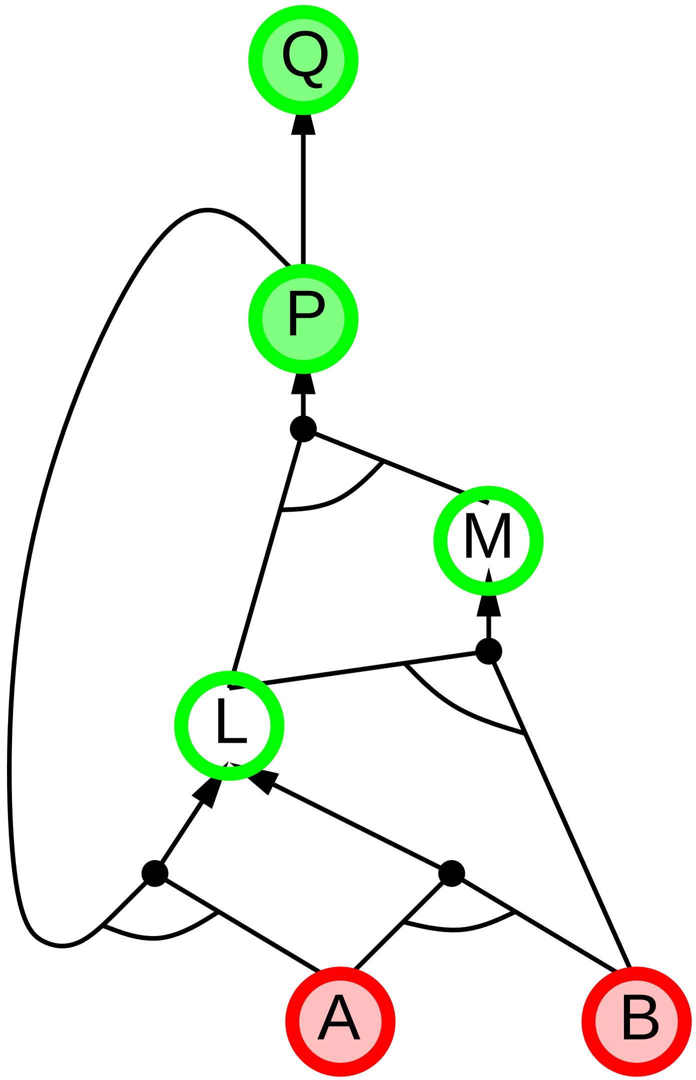
Backward chaining example



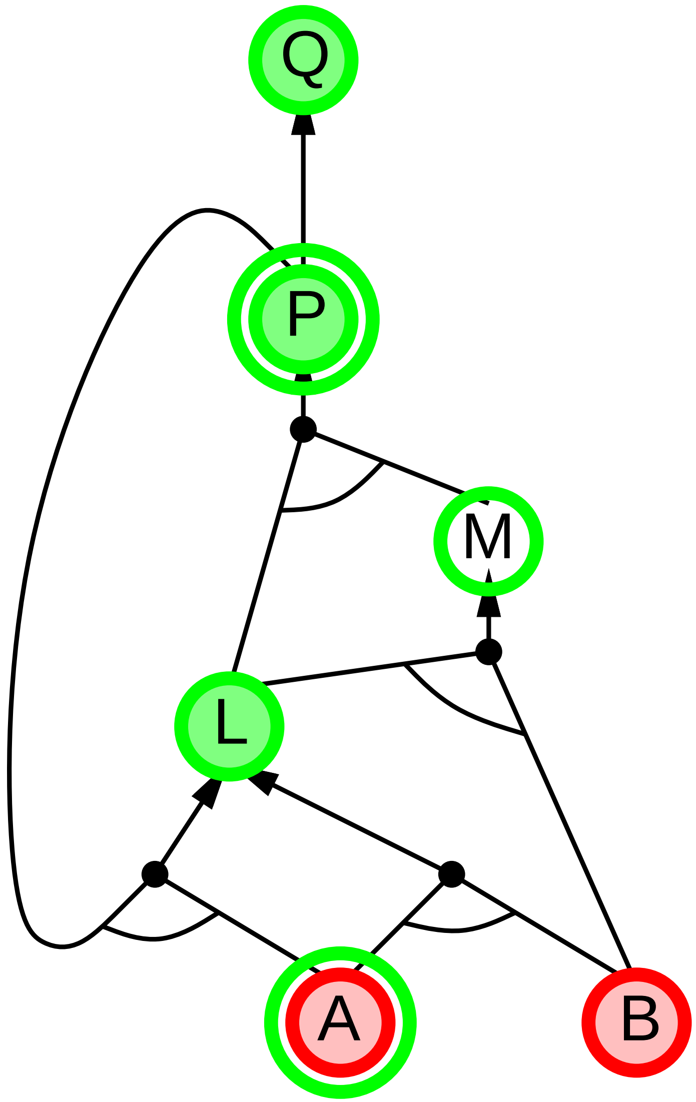
Backward chaining example



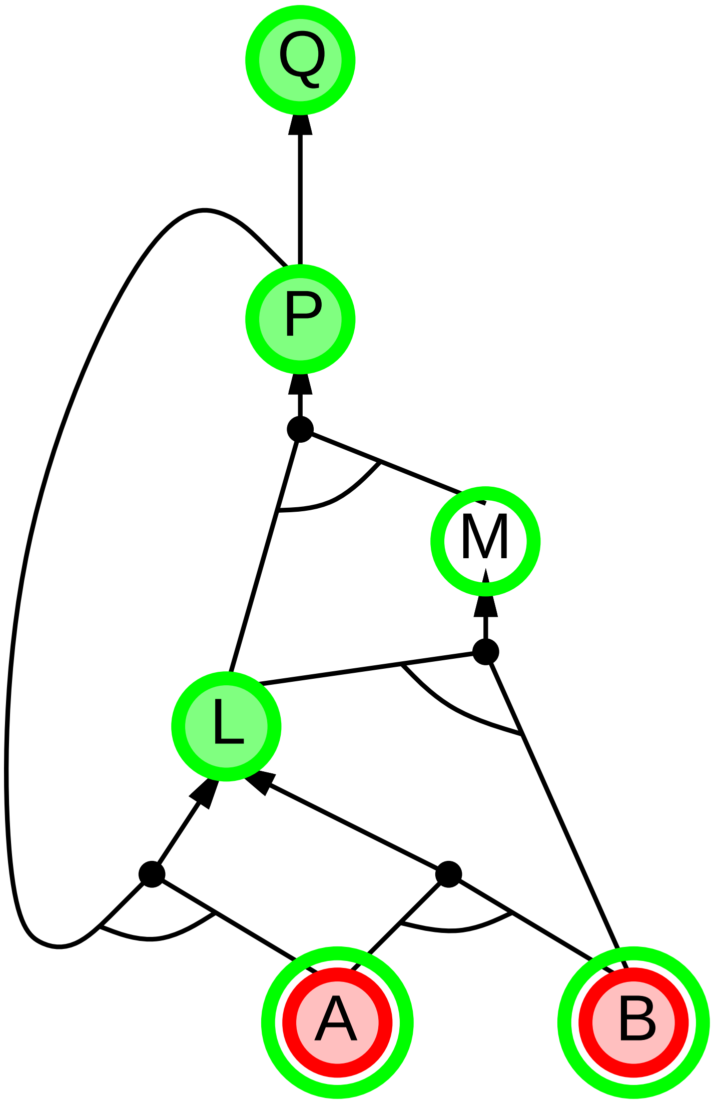
Backward chaining example



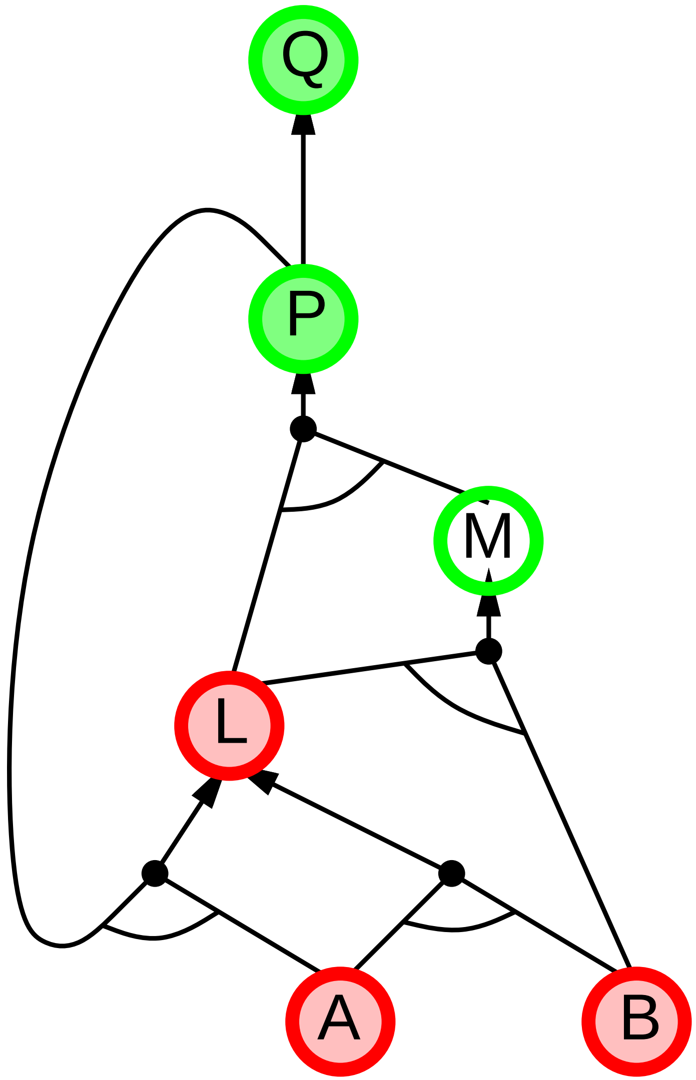
Backward chaining example



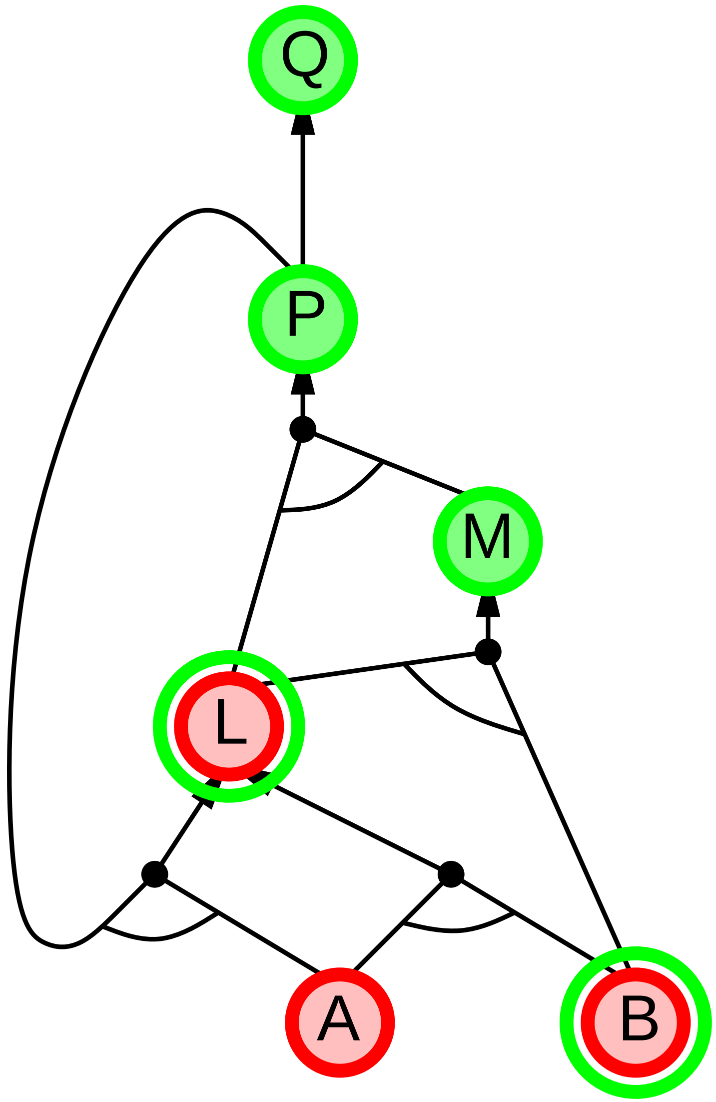
Backward chaining example



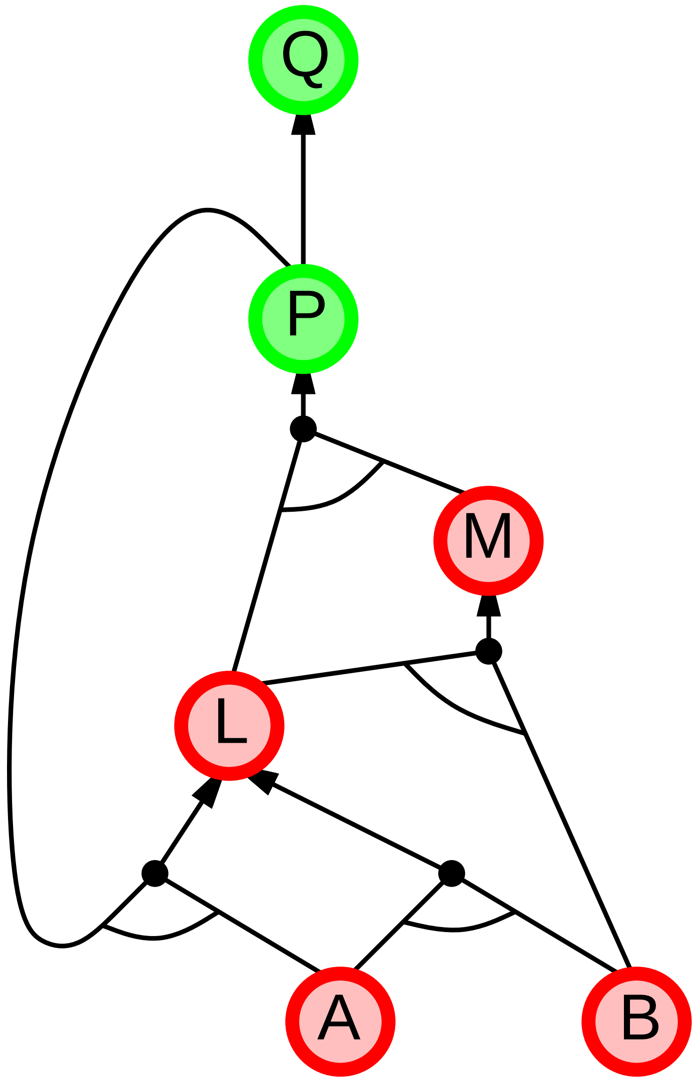
Backward chaining example



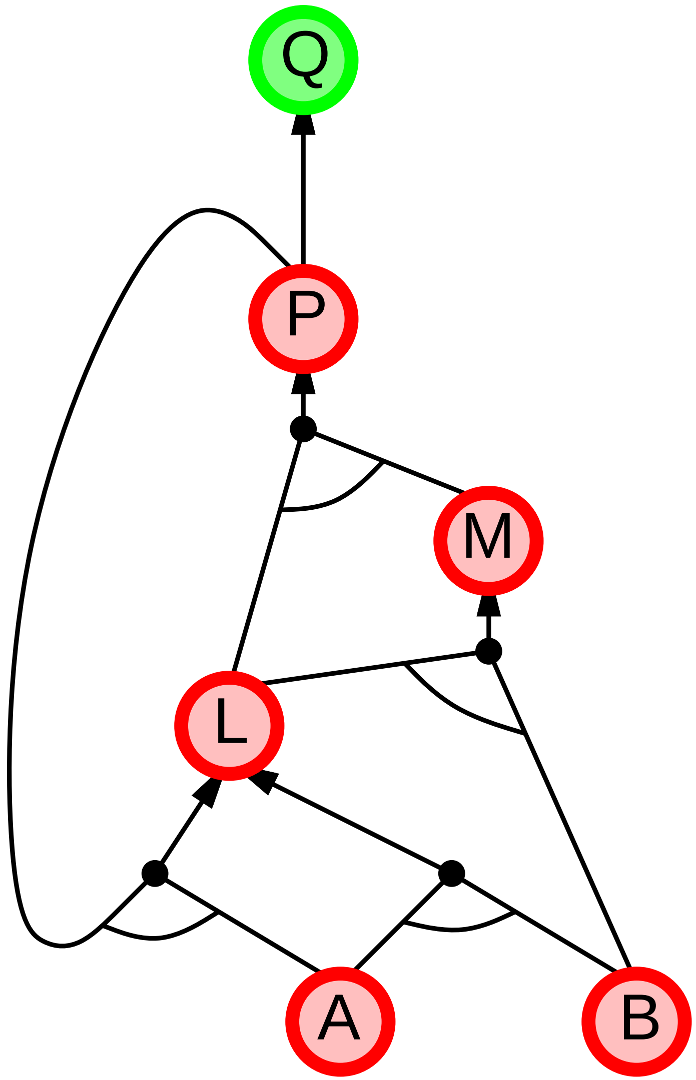
Backward chaining example



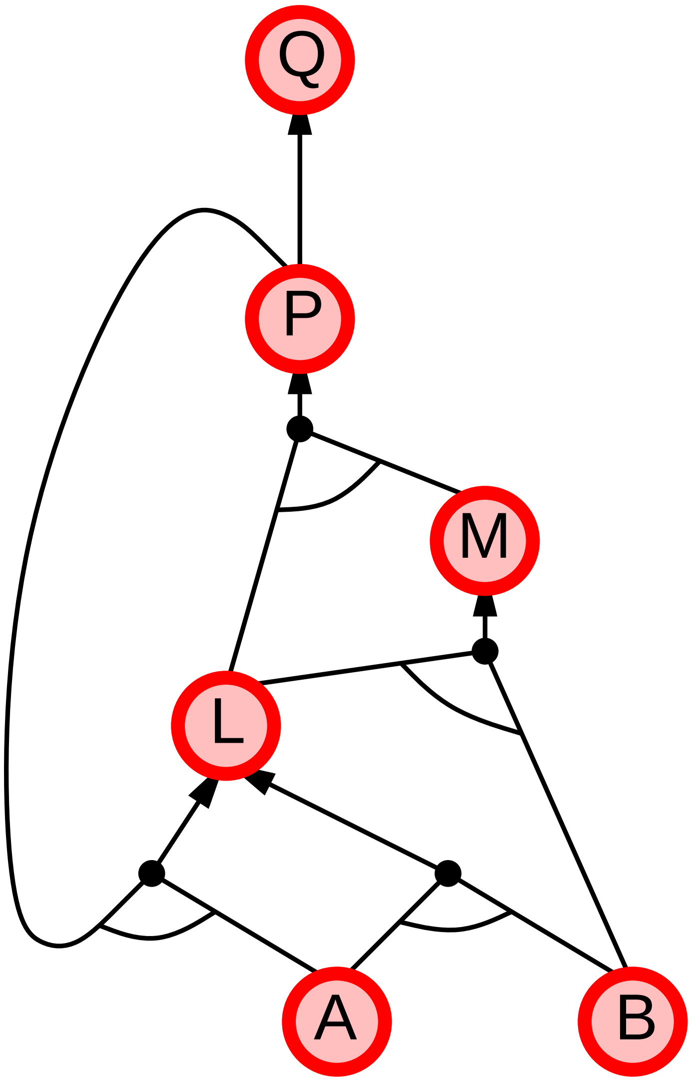
Backward chaining example



Backward chaining example



Backward chaining example



Forward vs. Backward Chaining

FC is data-driven, cf. automatic, unconscious processing
e.g., object recognition, routine decisions

May do lots of work that is irrelevant to the goal

BC is goal-driven, appropriate for problem-solving,
e.g., Where are my keys? How do I get into a PhD program?

Complexity of BC can be much smaller than linear in size of KB

Model Checking methods

The most effective procedures for propositional satisfiability are based on CSP techniques

Variable domain: $\{true, false\}$

Constraints: sets of clauses

Heuristic Improvements:

- unit propagation
- variable and value ordering
- intelligent backtracking
- clause learning
- random restarts
- clever indexing
- subproblem decomposition

DPLL Procedure

function DPLL-SATISFIABLE?(s) **returns** *true* or *false*
inputs: s, a sentence in propositional logic

$clauses \leftarrow$ the set of clauses in the CNF representation of s
$symbols \leftarrow$ a list of the proposition symbols in s
return DPLL($clauses$, $symbols$, [])

DPLL Procedure (cont.)

function DPLL(*clauses*, *symbols*, *model*) **returns** *true* or *false*

if every clause in *clauses* is satisfied by *model* **then return** *true*
if some clause in *clauses* is falsified by *model* **then return** *false*
P, *value* $\leftarrow$ FIND-PURE-SYMBOL(*symbols*, *clauses*, *model*)
if P is non-null **then**
    return DPLL(*clauses*, *symbols* - P, $(P \mapsto \mathit{value}) :: \mathit{model}$)
P, *value* $\leftarrow$ FIND-UNIT-CLAUSE(*clauses*, *model*)
if P is non-null **then**
    return DPLL(*clauses*, *symbols* - P, $(P \mapsto \mathit{value}) :: \mathit{model}$)
$P \leftarrow$ FIRST(*symbols*)
rest $\leftarrow$ REST(*symbols*)
return DPLL(*clauses*, *rest*, $(P \mapsto \mathit{true}) :: \mathit{model}$) **or**
        DPLL(*clauses*, *rest*, $(P \mapsto \mathit{false}) :: \mathit{model}$)

DPLL Exercise

Use DPLL to check the satisfiability of the following set of clauses

(1) $\neg p_1 \lor p_2$	(2) $\neg p_3 \lor p_4$	(3) $\neg p_6 \lor \neg p_5 \lor \neg p_2$
(4) $\neg p_5 \lor p_6$	(5) $p_5 \lor p_7$	(6) $\neg p_1 \lor p_5 \lor \neg p_7$