CS:5810 Formal Methods in Software Engineering

Case Study: Hotel Lock System

Acknowledgments

These notes are based on an Alloy example in the book:

[Jack06] Daniel Jackson. *Software abstractions – Logic, Language, and Analysis.* The MIT press, 2006

The Task

- Model in Alloy the disposable card key system used in most hotels for locking and unlocking guest rooms
- The system uses recordable locks, which prevent previous guests from entering a room once its has been re-assigned
- We will model both static and dynamic aspects of the system

Problem Description [Jack06]

"[…] the hotel issues a new key to the next occupant, which recodes the lock, so that previous keys will no longer work.

The lock is a simple, stand-alone unit […] with a memory holding the current key combination.

A hardware device […] [within the lock] generates a sequence of pseudorandom numbers."

Problem Description [Jack06]

"The lock is opened either by the current key combination, or by its successor;

If a key with the successor is inserted, the successor is made to be the current combination, so that the old combination will no longer be accepted.

This scheme requires no communication between the front desk and the door lock."

Problem Description [Jack06]

“By synchronizing the front desk and the door locks initially, and by using the same pseudorandom generator,

the front desk can keep its records of the current combinations in step with the doors themselves.”

Signatures and Fields

Signatures: `Key`, `Room`, `Guest`, `FrontDesk`

- `Key` refers to the key combination stored in the magnetic strip of the card
- `FrontDesk` stores at any time a mapping
  - between each room and its most recent key combination (if any), and
  - between each room and its current guest

Signatures and Fields

`Room` refers to the room lock

Each room (lock) has

- an associated set of possible keys, and
- exactly one current key at a time

Each key belongs to at most one room

Each guest has zero or more keys at any time

Signatures and Fields

```
module hotel
open util/ordering [Key] as KO
```

Signatures and Fields

```
module hotel
open util/ordering [Key] as KO

sig Key {}

sig Room {
  keys: set Key,
  var currentKey: Key
}

sig Guest {
  var keys: set Key
}

one sig FrontDesk {
  var lastKey: Room -> lone Key,
  var occupant: Room -> Guest
}
```

```
fun sig FDlastKey : Room -> Key {
  FrontDesk.lastKey
}

fun sig FDoccupant: Room -> Guest {
  FrontDesk.occupant
}
```

Room Constraint

Each key belongs to at most one room

```
fact {
  all k: Key | lone (Room <: keys).k
}
```

```
module hotel
open util/ordering [Key] as KO

sig Key {}

sig Room {
  keys: set Key,
  var currentKey: Key
}

sig Guest { var keys: set Key }

one sig FrontDesk {
  var lastKey: Room -> lone Key,
  var occupant: Room -> Guest
}

fun sig FDlastKey : Room -> Key {
  FrontDesk.lastKey
}

fun sig FDoccupant: Room -> Guest
   {
  FrontDesk.occupant
}
```

New Key Generation

Given a key k and a set ks of keys,

nextKey returns the smallest key (in the key ordering) in ks that follows k

```
fun nextKey [k: Key, ks: set Key]: set Key
{
  KO/min [KO/nexts[k] & ks]
}
```

Initial State

```
module examples/hotel
open util/ordering [Key] as KO

sig Key {}

sig Room {
  keys: set Key,
  var currentKey: Key
}

sig Guest {
  var keys: set Key
}

one sig FrontDesk {
  var lastKey: Room -> lone Key,
  var occupant: Room -> Guest
}
```

No constraints

No guests have keys

the record of each room's key at the front desk is synchronized with the current combination of the lock itself

No rooms are occupied

Hotel Operations: Initial State

```
pred init {
  -- no guests have keys
  no Guest.keys
  -- the roster at the front desk shows
  -- no room as occupied
  no FDoccupant
  -- the record of each room’s key at the
  -- front desk is synchronized with the
  -- current combination of the lock itself
  all r: Room |
    r.FDlastKey = r.currentKey
}
```

Hotel Operations: Guest Entry

```
pred entry [ g: Guest, r: Room, k: Key ]
```

- Preconditions:
  - The key used to open the lock is one of the keys the guest is holding
- Pre and Post Conditions:
  - The key on the card
    - either matches the lock's current key, and the lock remains unchanged (not a new guest), or
    - matches its successor, and the lock is advanced (new guest)
- Frame conditions:
  - no changes to the state of other rooms, to the set of keys held by guests, or to the records at the front desk

Hotel Operations: Guest Entry

```
pred entry [ g: Guest, r: Room, k: Key ] {

  -- preconditions
     -- the key used to open the lock is one of the keys
     -- held by the guest
     k in g.keys

  -- pre and post conditions
     (-- not a new guest
      k = r.currKey and r.currKey' = r.currKey
      or
      -- new guest
      k = nextKey[r.currKey, r.keys] and r.currKey' = k
     )

  -- frame conditions
     noFrontDeskChange
     noRoomChange[Room - r]
     noGuestChange[Guest]
}
```

Hotel Operations: Check-out

```
pred checkOut [ g: Guest ]
```

- Preconditions:
  - the guest occupies one or more rooms
- Postconditions:
  - the guest’s rooms become available
- Frame conditions:
  - Nothing changes but the `occupant` relation

Hotel Operations: Check-out

```
pred checkOut [g: Guest] {

  -- preconditions
     -- the guest occupies one or more rooms
     some FDoccupant.g

  -- postconditions
     -- the guest's rooms become available
     FDoccupant' = FDoccupant - (Room -> g)

  -- frame conditions
     FDlastKey' = FDlastKey
     noRoomChange[Room]
     noGuestChange[Guest]
}
```

Hotel Operations: Check-in

```
pred checkIn [ g: Guest, r: Room, k: Key ]
```

- Preconditions:
  - the room is available
  - the input key is the successor of the last key in the sequence associated to the room
- Postconditions:
  - the guest holds the input key and becomes the new occupant of the room
  - the input key becomes the room’s current key
- Frame conditions:
  - Nothing changes but the occupant relation and the guest’s relations

Hotel Operations: Check-in

```
pred checkIn [ g: Guest, r: Room, k: Key ] {
  -- preconditions
     -- the room has no current occupant
     no r.FDoccupant
     -- the input key is the successor of the last key in
     -- the sequence associated to the room
     k = nextKey[r.FDlastKey, r.keys]

  -- postconditions
     -- the guest becomes the new occupant of the room
     FDoccupant' = FDoccupant + (r -> g)
     -- the guest holds the input key
     g.keys' = g.keys + k
     -- the input key becomes the room's current key
     FDlastKey' = FDlastKey ++ (r -> k)

  -- frame conditions
     noRoomChange[Room]
     noGuestChange[Guest - g] }
```

Trace Generation

- The first time step satisfies the initialization conditions
- Any pair of consecutive time steps are related by
  - an entry operation, or
  - a check-in operation, or
  - a check-out operation

Trace Generation

```
fact Traces {
  init

  always
    some g: Guest, r: Room, k: Key |
      entry[g, r, k] or
      checkin[g, r, k] or
      checkout[g]
}
```

Analysis

Let’s check if unauthorized entries are possible:

- If a guest `g` enters room `r` at time `t`, and the front desk records show `r` as occupied at that time, then `g` must be a recorded occupant of `r`.

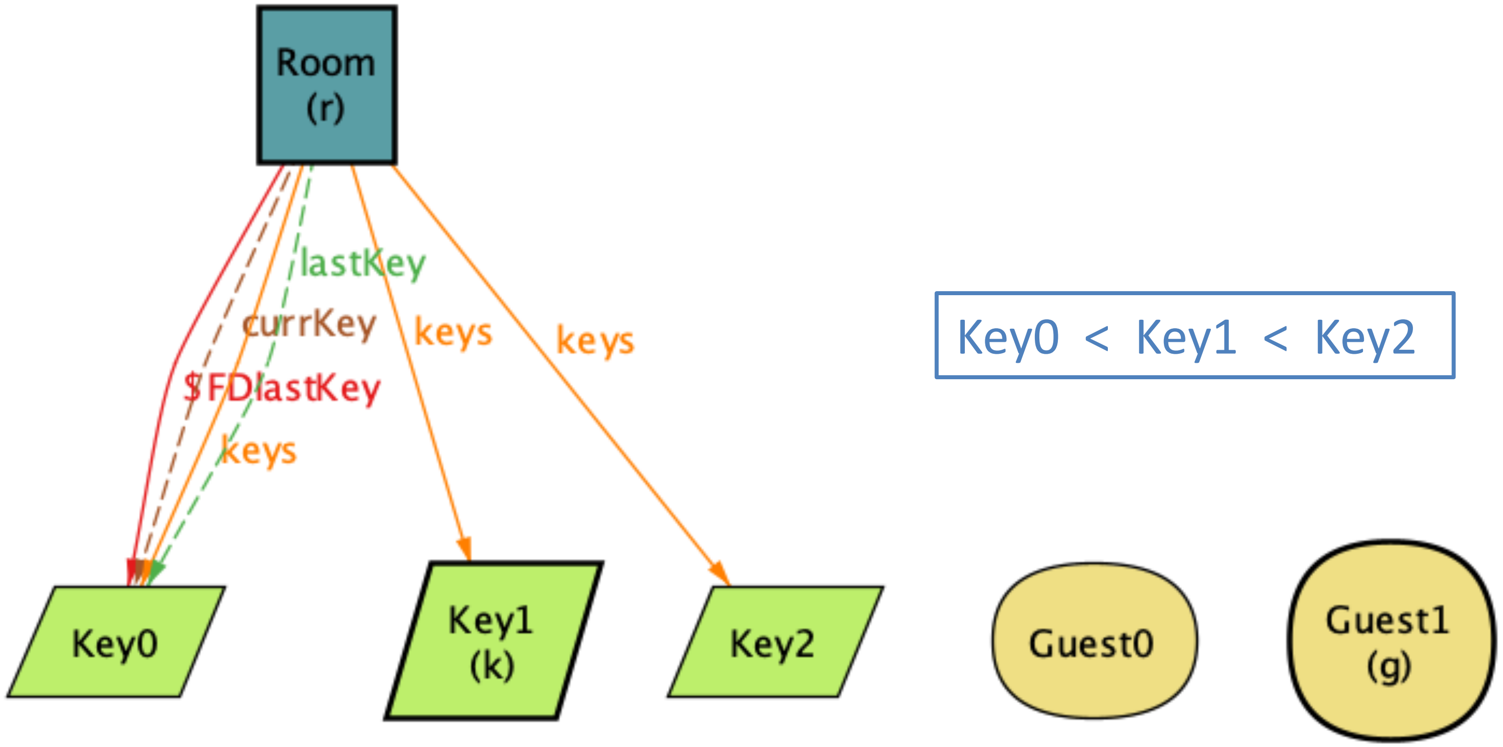
```
assert noBadEntry {
  always all r: Room, g: Guest, k: Key |
    let o = r.FDoccupant |
      (entry[g, r, k] and some o) implies g in o
}
```

Analysis

```
check noBadEntry for 3
but 1 Room, 2 Guest, 5 steps
```

- It is enough to check for problem already with just 2 guests and 2 rooms
- `steps`'s scope must be at least 5 because at least 4 steps are needed to execute each operation once
- There is a counterexample
  (see file `dynamic/hotel1-elec.als`)

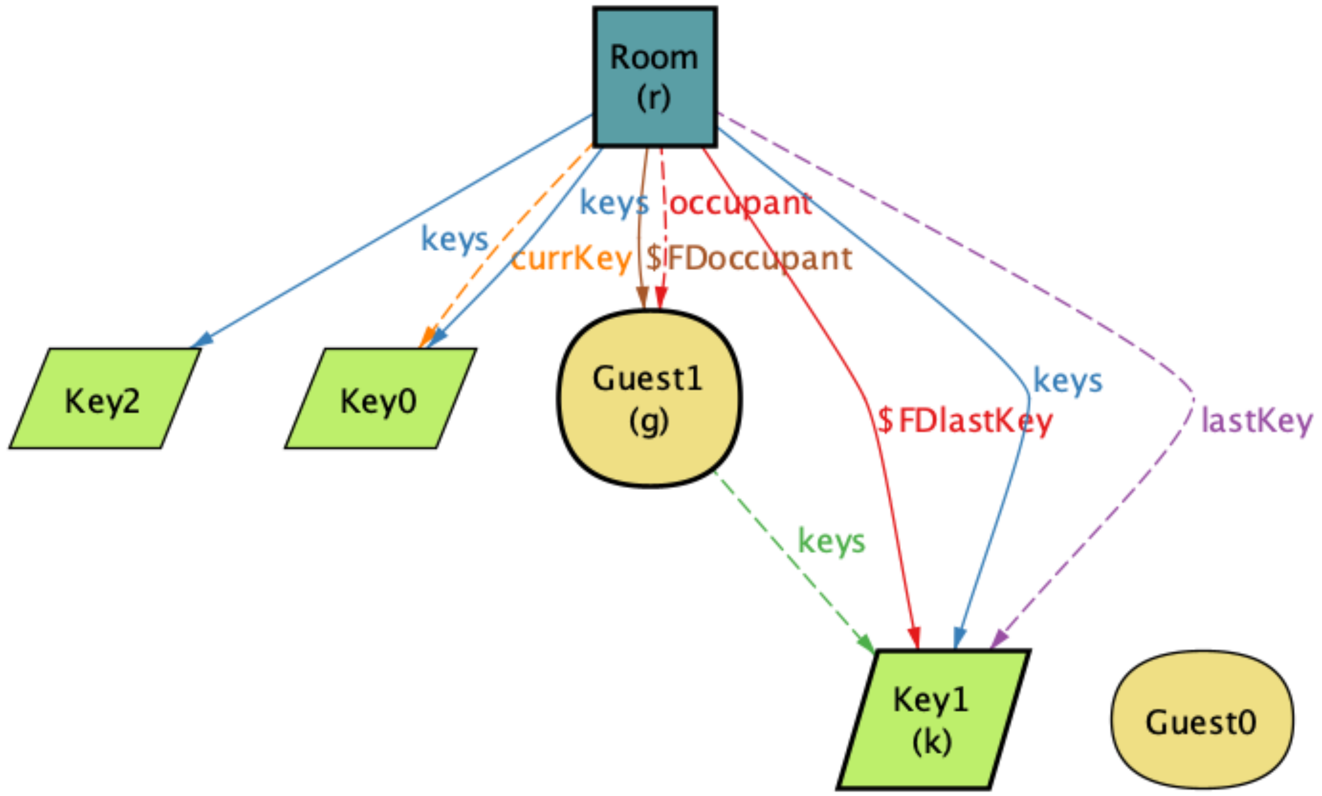
0: Initial State



Key0 < Key1 < Key2

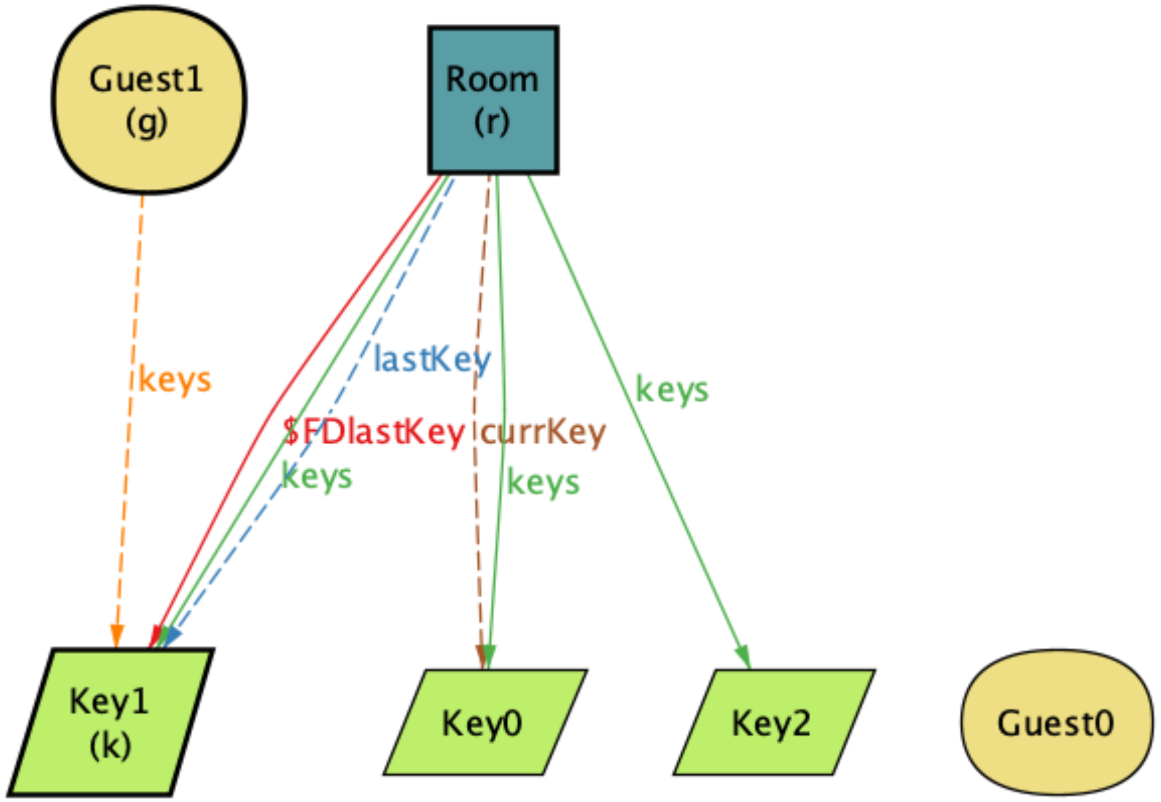
- initially, the current key of Room is Key0
- this is also reflected in the front desk's record

1: CheckIn Operation



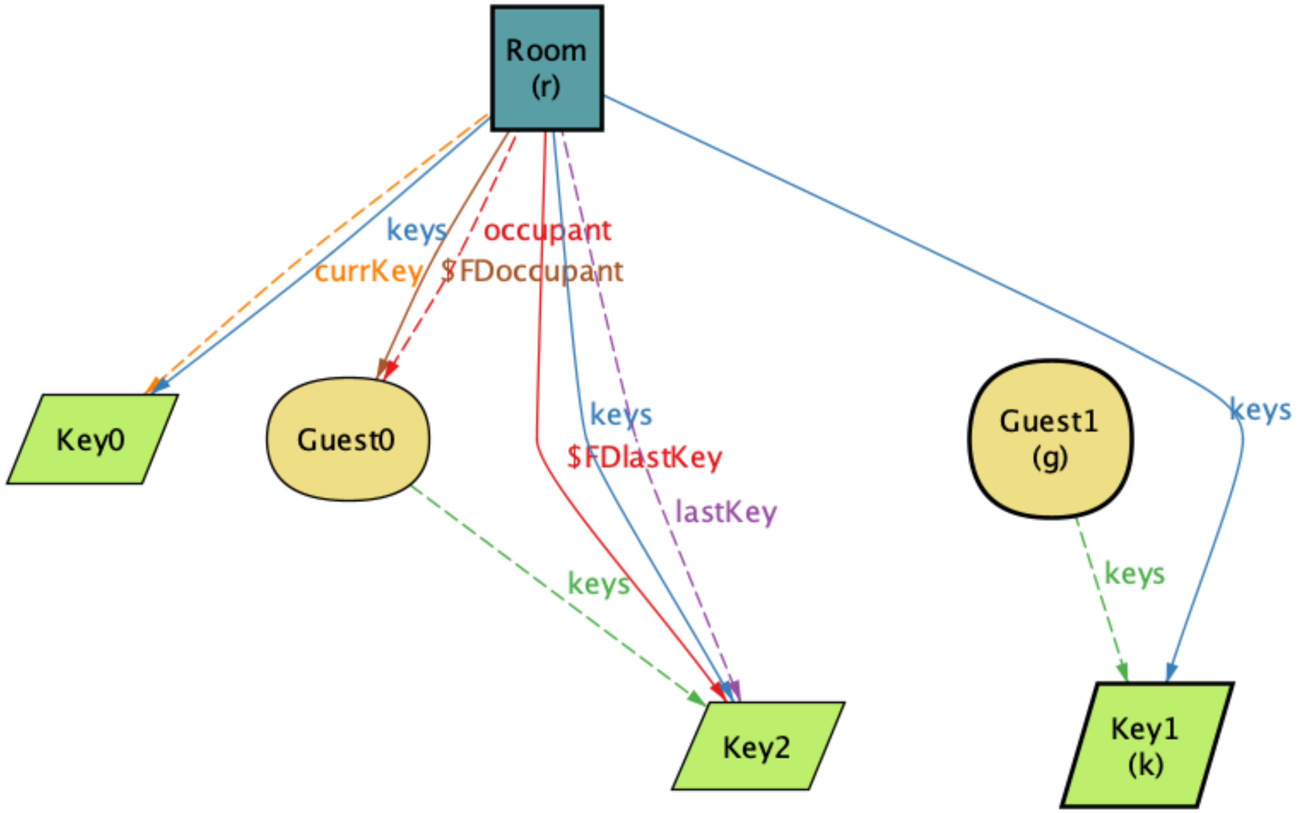
- Guest1 checks in to Room and receives key Key1
- the occupancy roster at the front desk is updated accordingly
- Key1 is recorded as the last key assigned to Room

2: CheckOut Operation



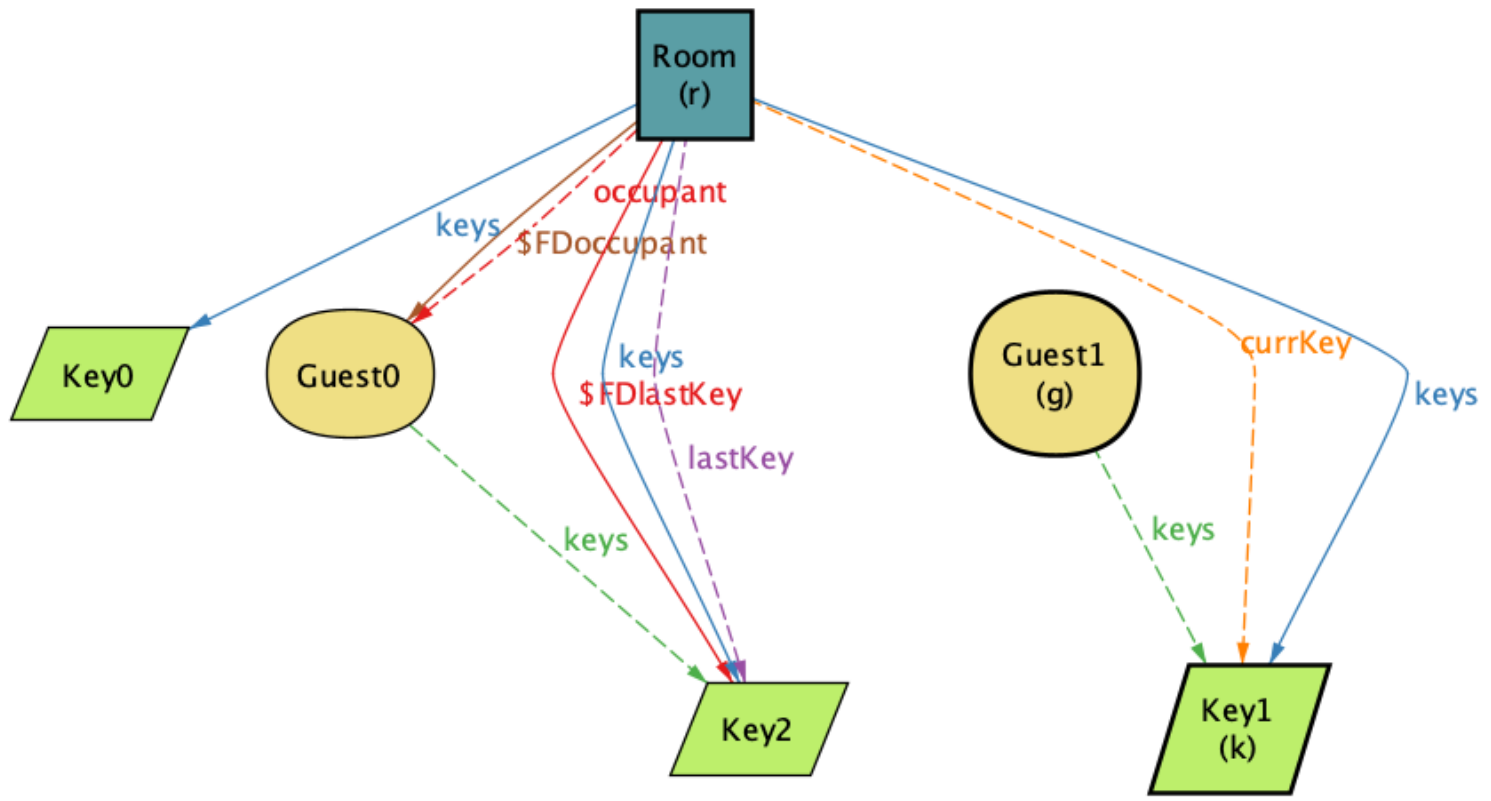
- Guest1 checks out, and the occupancy roster is cleared
- Since Guest1 never entered, Room's current key is still Key0
- Guest1 still holds Key1

3: CheckIn Operation



- Guest0 checks into Room and receives key Key2
- the occupancy roster at the front desk is updated accordingly
- Key2 is recorded as the last key assigned to Room

4: Entry Operation



- Guest1 presents Key1 to the lock of Room
- Since Room's current key is still Key0, Guest1 is admitted

Necessary Restriction

There must be no intervening operation
between a guest’s check-in and their room entry

```
pred noInterveningOps {
  always
    all g: Guest, r: Room, k: Key |
      checkIn[g, r, k] implies after entry[g, r, k]
}
```

Conditional Assertion

Make assertion under noInterveningOps assumption

```
  assert noBadEntry {

    noInterveningOps implies

      always all r: Room, g: Guest, k: Key |
        let o = r.FrontDesk.occupant |
          (entry[g, r, k] and some o) implies g in o
}
```

Analysis

- We check once again:

```
check noBadEntry for 3
   but 2 Room, 2 Guest, 5 steps
```

  – No counter-example (see file `dynamic/hotel2-elec.als`)

- For greater confidence, we increase the scope:

```
check noBadEntry for 5
   but 3 Room, 3 Guest, 20 steps
```

  – No counterexamples