

CS:5810 Formal Methods in Software Engineering

Dynamic Models in Alloy

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Overview

- Basics of dynamic models
 - Modeling a system's **states** and **state transitions**
 - Modeling **operations** causing transitions
- Simple example of operations

Static Models

- So far, we've used Alloy to define the allowable values of **state** components
 - values of **sets**
 - values of **relations**
- A model instance is a **set of state component values** that
 - Satisfies the **constraints** defined by multiplicities, fact, “realism” conditions, ...

Static Model Instances

```
Person = {Matt, Sue}
Man = {Matt}
Woman = {Sue}
Married = {}
spouse = {}
children = {}
siblings = {}
```

```
Person = {Matt, Sue}
Man = {Matt}
Woman = {Sue}
Married = {Matt, Sue}
spouse = {(Matt,Sue), (Sue,Matt)}
children = {}
siblings = {}
```

```
Person = {Matt, Sue, Sean}
Man = {Matt, Sean}
Woman = {Sue}
Married = {Matt, Sue}
spouse = {(Matt,Sue), (Sue,Matt)}
children = {(Matt,Sean), (Sue,Sean)}
siblings = {}
```

Dynamic Models

- Static models let us describe the legal **states** of a **dynamic** system
- However, we'd like also to be able to describe possible **transitions** between states

E.g.

- Two unmarried people become each other's spouses once they get married
- People go from being alive to not being alive when they die
- A person becomes someone's child after being born

Example

Family Model

```
abstract sig Person {
    children: set Person,
    siblings: set Person
}

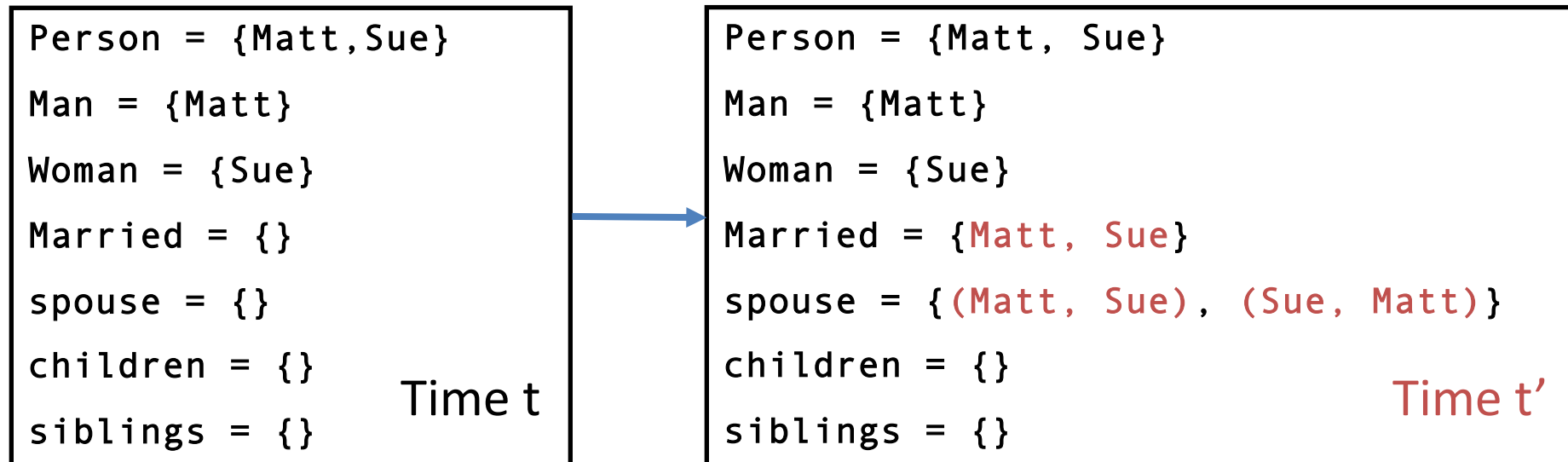
sig Man, Woman extends Person {}

sig Married in Person {
    spouse: one Married
}
```

State Transitions

Two people get married

- At time t , `spouse` = `{}`
 - At time t' , `spouse` = `{(Matt, Sue), (Sue, Matt)}`
- ⇒ We can add the notion of `time` in the `spouse` relation



Modeling State Transitions

- Until version 6, Alloy had **no predefined notion** of time and of state transition
- This is not really a problem since there are **several ways to model dynamic aspects** of a system in Alloy
- A **general** and relatively simple **way** is to:
 1. **introduce** a **Time signature** expressing time
 2. **add a time component** to each relation that changes over time

Family Model Signatures

```
abstract sig Person {  
    children: set Person,  
    siblings: set Person  
}
```

```
sig Man, Woman extends Person {}
```

```
sig Married in Person {  
    spouse: one Married  
}
```

Family Model Signatures with Time

```
sig Time {}
```

```
abstract sig Person {  
    children: Person set -> Time,  
    siblings: Person set -> Time  
}
```

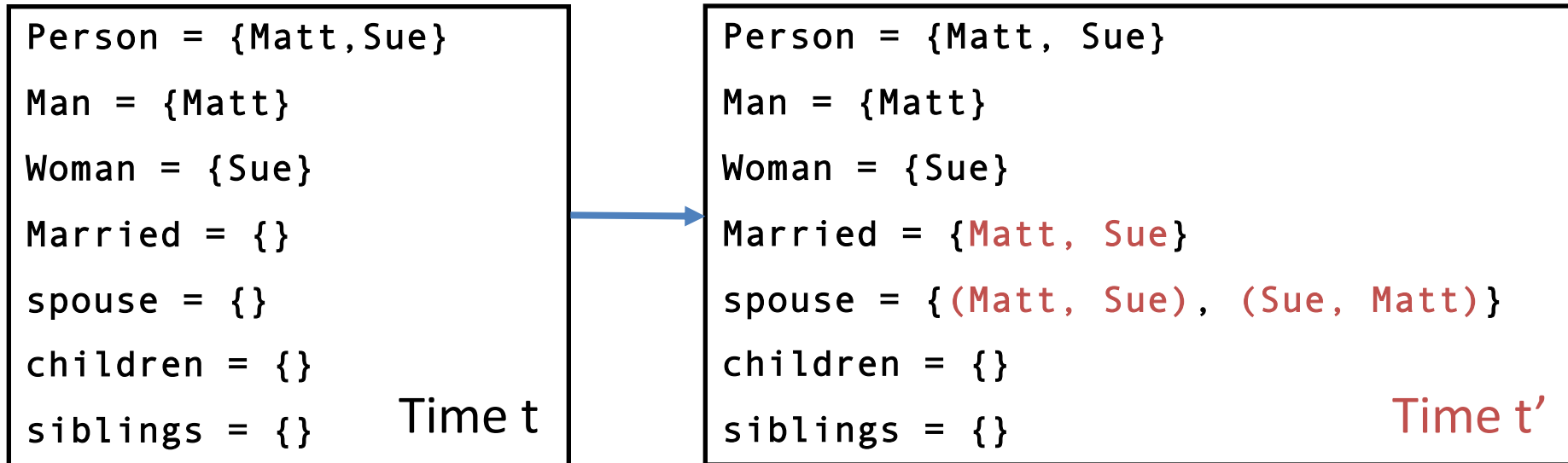
```
sig Man, Woman extends Person {}
```

```
sig Married in Person {  
    spouse: Married one -> Time  
}
```

Transitions

Two people get married

- At time t , `Married` = `{}`
- At time t' , `Married` = `{Matt, Sue}`

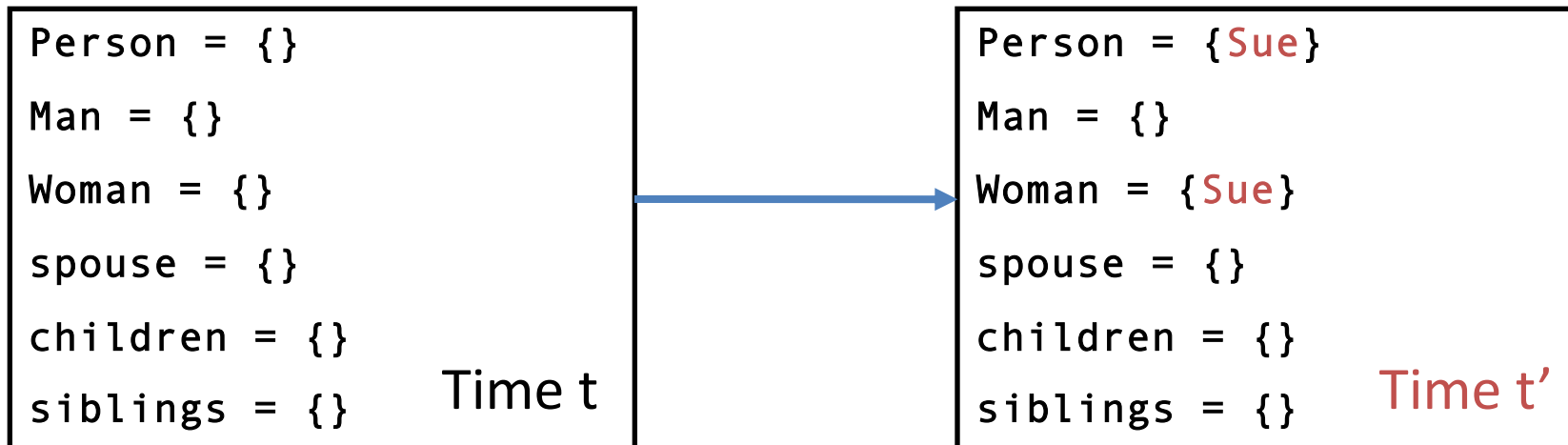


Transitions

A person is born

- At time t , `Person` = `{}`
- At time t' , `Person` = `{Sue}`

For simplicity, we will not use time-dependent signatures



Keeping Signatures Static

```
abstract sig Person {  
  children: Person set -> Time,  
  siblings: Person set -> Time,  
  spouse: Person lone -> Time  
}  
sig Man, Woman extends Person {}
```

```
sig Married in Person {  
  spouse: Married one -> Time  
}
```

Keeping Signatures Static

```
abstract sig Person {  
  children: Person set -> Time,  
  siblings: Person set -> Time,  
  spouse: Person lone -> Time,  
  alive: set Time  
}  
  
sig Man, Woman extends Person {}
```

Revising Constraints

```
abstract sig Person {
  children: Person set -> Time,
  siblings: Person set -> Time,
  spouse: Person lone -> Time,
  alive: set Time,
  parents: Person set -> Time
}

sig Man, Woman extends Person {}

fun parents[] : Person >Person { ~children }
fact parentsDef {
  all t: Time | parents.t = ~(children.t)
}
```

Revising Constraints

```
-- Time-dependent parents relation
fact parentsDef {
  all t: Time | parents.t = ~(children.t)
}

-- Two persons are blood relatives (at time t) iff
-- they have a common ancestor (at time t)
pred BloodRelatives [p, q: Person, t: Time]
{
  some p.*(parents.t) & q.*(parents.t)
}
```


Revising *Static* Constraints

-- People cannot be their own ancestors (at any time)

```
all t: Time | no p: Person |  
  p in p.^(parents.t)
```

-- No one can have more than one father or mother (at any time)

```
all t: Time | all p: Person |  
  lone (p.parents.t & Man)  
  and  
  lone (p.parents.t & Woman)
```

...

Revising *Static* Constraints

-- (At all times) your siblings are those people other than you

-- who have the same parents you have

```
all t: Time | all p: Person |  
  p.siblings.t = { q: Person - p | some q.parents.t and  
                                p.parents.t = q.parents.t }
```

-- (At all times) the spouse relation is symmetric

```
all t: Time |  
  spouse.t = ~(spouse.t)
```

Revising *Static* Constraints

-- (At all times) a spouse can't be a sibling

```
all t: Time | no p: Person |  
  some p.spouse.t and p.spouse.t in p.siblings.t
```

-- (At all times) people can't be married to a blood relative

```
all t: Time | no p: Person |  
  let s = p.spouse.t |  
    some s and BloodRelatives[p, s, t]
```

Revising *Static* Constraints

-- (At all times) a person can't have children with a blood relative

```
all t: Time | all p, q: Person |  
  (some (p.children.t & q.children.t) and p != q)  
  implies  
  not BloodRelatives[p, q, t]
```

A Better Approach: Mutable Relations

Alloy 6 incorporates an **implicit**, built-in notion of (discrete) **time**

- The meaning of an Alloy model is actually an infinite **sequence** of instances, or a **trace**
- Each instance in a trace corresponds to a state of a dynamic system
- Signatures/relations can **change** from state to state
- A set of **temporal operators** allows us to express properties over time as properties over traces

Mutable Fields: Relations that Change Over Time

```
abstract sig Color {}
```

```
one sig      Green, Yellow, Red extends Color {}
```

```
one sig TrafficLight
```

```
{
```

```
  var col: Color  -- value here can change over time
```

```
}
```

```
fun c : Color { TrafficLight.col }
```

Mutable Fields: Relations that Change Over Time

-- enum abbreviates a partition of a signature into singletons

```
enum Color { Green, Yellow, Red }
```

```
one sig TrafficLight
```

```
{
```

```
  var col: Color -- value here can change over time
```

```
}
```

```
fun c : Color { TrafficLight.col }
```

From Instances to Traces

- Models with **mutable** signatures and/or fields represent *dynamic systems*, systems that change over time
- Instead of standing for a set of instances, a dynamic model stands for a set of *traces*
- A trace is an infinite sequence of instances
 - An **instance** now describes just one possible *state* of a system
 - A trace describes a particular sequence of *state transitions* for the system

From Instances to Traces

An Alloy model captures the behavior of a system over time by means of constraints containing *temporal operators*

Temporal operators implicitly talk about (properties of) traces

Temporal Operators in Alloy 6

Formula	Intuitive meaning
<code>always p</code>	<code>p</code> holds from current state/instance forward in a trace
<code>historically p</code>	<code>p</code> holds from current state backward
<code>after p</code>	<code>p</code> holds in next state (after the current one)
<code>before p</code>	<code>p</code> holds in previous state (before the current one)
<code>eventually p</code>	<code>p</code> holds in current state or a later one
<code>once p</code>	<code>p</code> holds in current state or an earlier one
<code>p until q</code>	<code>q</code> holds eventually and <code>p</code> holds continuously until then
<code>p since q</code>	<code>p</code> has held continuously since last time <code>q</code> held
<code>e'</code>	denotes the value of <code>e</code> in next state

Example Traces

Time steps	1	2	3	4	5	6	7	8	9	...													
p (state prop.)	•	•	•	•	•		•	•	•	•	•				•	•	•	•	•	•	•	•	...
q (state prop.)						•								•	•								...
always p															•	•	•	•	•	•	•	•	...
historically p	•	•	•	•	•																		...
after p	•	•	•	•		•	•	•	•	•				•	•	•	•	•	•	•	•	•	...
before p		•	•	•	•	•		•	•	•	•	•				•	•	•	•	•	•	•	...
eventually q	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•								...
once q						•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	•	...
p until q	•	•	•	•	•	•								•	•								...
p since q						•	•	•	•	•	•			•	•	•	•	•	•	•	•	•	...

• = true

Temporal Operator Precedence



The Family Model with Mutable Fields

```
enum Liveness { Alive, Dead, Unborn }
```

```
abstract sig Person {  
  var children: set Person,  
  var parents: set Person,  
  var siblings: set Person,  
  var spouse: lone Person,  
  var liveness: Liveness  
}  
sig Man, Woman extends Person {}
```

Revising the Model

```
enum Liveness { Alive, Dead, Unborn }
```

```
abstract sig Person {  
  var children: set Person,  
  var spouse: lone Person,  
  var liveness: Liveness  
}
```

```
sig Man, Woman extends Person {}
```

```
fun parents : Person -> Person { ~children }
```

```
fun siblings [p: Person]: Person { {q: Person | ... } }
```

Revising Constraints

```
pred BloodRelatives [p, q: Person] {
    some p.*parents & q.*parents
}
pred isAlive [p: Person] { p.liveness = Alive }
pred isDead [p: Person] { p.liveness = Dead }
pred isUnborn [p: Person] { p.liveness = Unborn }

-- a newborn is someone who has just been born
pred newBorn[p: Person] {
    isAlive[p] and before isUnborn[p]
}

pred isMarried [p: Person] { some p.spouse }
```

Revising *Static* Constraints

-- People cannot be their own ancestors

always no p: Person | p in p.^parents

-- No one can have more than one father or mother

always all p: Person |
lone (p.parents & Man) and lone (p.parents & Woman)

-- The spouse relation is symmetric

always spouse = ~spouse

Revising *Static* Constraints

-- A spouse can't be a sibling

```
always no p: Person |  
    some p.spouse and p.spouse in p.siblings
```

-- People can't be married to a blood relative

```
always no p: Person | let s = p.spouse |  
    some s and BloodRelatives[p, s]
```

-- A person can't have children with a blood relative

```
always all disj p, q: Person |  
    some (p.children & q.children) implies  
    not BloodRelatives[p, q]
```

Adding *Temporal* Constraints

-- Dead people stay dead

```
always all p: Person |  
  isDead[p] implies after isDead[p]
```

-- Dead people were once alive

```
always all p: Person |  
  isDead[p] implies once isAlive[p]
```

-- No one lives forever

```
always all p: Person |  
  isAlive[p] implies eventually isDead[p]
```

Adding *Temporal* Constraints

-- Living people don't become unborn

```
always all p: Person |  
  isAlive[p] implies always not isUnborn[p]
```

-- Living people stay alive until they die

```
always all p: Person |  
  isAlive[p] implies (isAlive[p] until isDead[p])
```

-- Newborns have a father and a mother

```
always all p: Person | newBorn[p] implies  
  some m: Man | some w: Woman | p.parents = m + w
```

Adding *Temporal* Constraints

-- Children were born from previously alive parents

```
always all p, q: Person |  
  p in q.children implies  
    once (newBorn[p] and once isAlive[q])
```

-- People with parents have had those parents since birth

```
always all p, q: Person |  
  p in q.children implies  
    (p in q.children since newBorn[p])
```

Exercises

- Load `family-6-elec.a1s` in Alloy
- Execute it
- Analyze the model
- Look at the generated instance
- Does it look correct?
- What, if anything, would you change about it?

Dynamics as State Transitions

Recall

- The evolution of a dynamic system can be modeled as a set of **traces**
- Each trace is a **sequence** of **transitions** from one state to another

A transition can be thought of as **caused by** the application of a **state transformer**

A state transformer is an **operator** that modifies the current state

Possible Trace

```
Person = {Matt, Sue, Sean}
Man = {Matt, Sean}
Woman = {Sue}
spouse = {}
children = {}
liveness = {(Matt, U),
(Sue, A), (Sean, U)}
```

```
Person = {Matt, Sue, Sean}
Man = {Matt, Sean}
Woman = {Sue}
spouse = {(Matt, Sue), (Sue, Matt)}
children = {}
liveness = {(Matt, A), (Sue, A), (Sean, U)}
```

```
Person = {Matt, Sue, Sean}
Man = {Matt, Sean}
Woman = {Sue}
spouse = {}
children = {}
liveness = {(Matt, U),
(Sue, U), (Sean, U)}
```

```
Person = {Matt, Sue, Sean}
Man = {Matt, Sean}
Woman = {Sue}
spouse = {(Matt, Sue), (Sue, Matt)}
children = {(Matt, Sean), (Sue, Sean)}
liveness = {(Matt, A), (Sue, A), (Sean, A)}
```

Transitions

A person is born from parents

State transformer that
modifies the **children** and
liveness relations

```
Person = {Matt, Sue, Sean}
Man = {Matt, Sean}
Woman = {Sue}
spouse = {(Matt,Sue), (Sue,Matt)}
children = {}
liveness = {(Matt,Alive), (Sue,Alive),
            (Sean,Unborn)}
```



```
Person = {Matt, Sue, Sean}
Man = {Matt, Sean}
Woman = {Sue}
spouse = {(Matt,Sue), (Sue,Matt)}
children = {(Matt,Sean), (Sue,Sean)}
liveness = {(Matt,Alive), (Sue,Alive),
            (Sean,Alive)}
```


Expressing State Transitions in Alloy

A state transformer is modeled as a **predicate** over two states:

1. the state **right before** the transition (**current state**) and
2. the state **right after** it (**next state**)

We use the **temporal** operators of Alloy to express **constraints** on the current and the next state

(Single) **primed** field names refer to values in the **next** state

Expressing State Transformers

Pre-condition constraints

- Describe the states to which the transformer applies

Post-condition constraints

- Describes the effects of the transformer in generating the next state

Frame-condition constraints

- Describes what does not change between current state and next state of a transition

*Distinguishing the pre-, post- and frame-conditions
in comments provides useful documentation*

Example: Marriage

```
pred getMarried [p, q: Person] {  
  -- preconditions  
  -- p and q are both alive  
  isAlive[p] and isAlive[q]  
  -- neither is married  
  no (p + q).spouse  
  -- they are not blood relatives  
  not BloodRelatives[p, q]  
  -- post-conditions  
  -- p and q are each other's spouses  
  p.spouse' = q  
  q.spouse' = p  
  -- frame conditions  
  ??  
}
```

```
enum Liveness { Alive, Dead, Unborn }  
abstract sig Person {  
  var children: set Person,  
  var spouse: lone Person,  
  var liveness: Liveness }  
sig Man, Woman extends Person {}  
pred isAlive [p: Person] {  
  p.liveness = Alive  
}  
fun parents : Person -> Person { ~children }  
fun siblings [p: Person]: Person {  
  {q: Person | ... }  
}
```

spouse' is the next version of spouse

Frame Condition

How is each relation impacted by marriage?

- 5 relations :
 - children, ~~parents~~, ~~siblings~~
 - spouse
 - liveness
- The `parents` and `siblings` relations are **defined** in terms of the `children` relation
- Thus, the frame condition has only to consider `children`, `spouse` and `liveness`

Frame Condition Predicates

```
pred noChildrenChange [Ps: set Person] {  
  all p: Ps |  
    p.children' = p.children  
}
```

```
pred noSpouseChange [Ps: set Person] {  
  all p: Ps |  
    p.spouse' = p.spouse  
}
```

```
pred noLivenessChange [Ps: set Person] {  
  all p: Ps |  
    p.liveness' = p.liveness  
}
```

Marriage Operator

```
pred getMarried [p, q: Person]
{
  -- preconditions
  isAlive[p] and isAlive[q]
  no (p + q).spouse
  not BloodRelatives[p, q]
  -- post-conditions
  p.spouse' = q and q.spouse' = p
  -- frame conditions
  noSpouseChange[Person - (p + q)]
  noChildrenChange[Person]
  noLivenessChange[Person]
}
```

Instance of Marriage

...

```
pred someMarriage {  
    some m: Man | some w: Woman | getMarried[m, w]  
}  
  
-- there is a marriage initially  
run { someMarriage }  
  
-- there is a marriage initially or later  
run { eventually someMarriage }  
  
-- there is a marriage eventually but not initially  
run { not someMarriage and eventually someMarriage }
```

Birth from Parents Operator

```
pred isBornFromParents [p: Person, m: Man, w: Woman] {
-- Pre-conditions
  isUnborn[p]
  isAlive[w]
  once isAlive[m]

-- Post-condition
  liveness' = liveness + p
-- Post-condition and frame condition
  children' = children + (m -> p) + (w -> p)

-- Frame conditions
  noLivenessChange[Person - p]
  noSpouseChange[Person]
  noChildrenChange[Person - (m + w)] // redundant
}
```


Instance of Birth

```
pred someBirth {  
    some b: Person, m: Man, w: Woman |  
        isBornFromParents[b, m, w]  
}  
  
run { eventually someBirth }  
  
run { some b: Person, m: Man, w: Woman |  
    eventually (getMarried[m, w] and  
        eventually isBornFromParents[b, m, w])
```

Death Operator

```
pred dies [p: Person] {  
  -- Pre-condition  
  isAlive[p]  
  
  -- Post-condition  
  after isDead[p]  
  
  -- Post-condition and frame condition  
  let q = p.spouse |  
    spouse' = spouse - ((p -> q) + (q -> p))  
  
  -- Frame conditions  
  noChildrenChange[Person]  
  noLivenessChange[Person - p]  
}
```

Instance of Death

```
pred someDeath {  
    some p: Person | dies[p]  
}
```

```
run {  
    eventually someDeath  
}
```

```
run {  
    some p: Person |  
        isAlive[p] and after (isAlive[p] and eventually dies[p])  
}
```

Specifying Transition Systems

- A transition system can be defined as a set of **traces** (aka **executions**):
 - sequences of states generated by the operators
- In our family example, for every execution:
 - The initial state satisfies some initialization condition
 - All pairs of consecutive states are related by
 - a birth operation, or
 - a death operation, or
 - a marriage operation

Initial State Specification

`init` specifies constraints on the initial state

```
pred init {  
  no children  
  no spouse  
  #LivingPeople > 2  
  #Person > #LivingPeople  
}
```

```
fun LivingPeople [] : Person {  
  liveness.Alive  
}
```

Transition Relation Specification

`trans` specifies that each transition is a consequence of the application of one of the operators to some individuals

```
pred trans [] {  
  (some m: Man, w: Woman | getMarried [m, w])  
  or  
  (some p: Person, m: Man, w: Woman |  
    isBornFromParents [p, m, w])  
  or  
  (some p: Person | dies [p])  
  or  
  other ???  
}
```

The Need for a No-op

- For convenience, Alloy considers only **infinite** traces
- So, we need a do-nothing operator for systems that can have **finite** executions

```
pred other [] {  
    -- the relevant relations stay the same  
    children' = children  
    spouse' = spouse  
    liveness' = liveness  
}
```

System Specification

A **System** predicate specifies that each execution

- starts in a state satisfying the initial state condition and
- moves from one state to the next as specified by one of the operator predicates

```
pred System {  
    init and always trans  
}
```

```
run { System }
```


System Invariants

- Many of the facts that we stated in our static model now become *expected system invariants*
- These are properties that
 - should *hold in initial states*
 - should *be preserved by* system *transitions*
- We can check that a property is invariant for a given system *System* (within a given scope) by
 - encoding the property as a formula *F* and
 - checking the assertion *System implies always F*

Expected Invariants: Examples

-- People cannot be their own ancestors

```
assert a1 { System implies
  always no p: Person | p in p.^parents
}
check a1 for 6
```

-- No one can have more than one father or mother

```
assert a2 { System implies
  always all p: Person |
    lone (p.parents & Man) and
    lone (p.parents & Woman)
}
check a2 for 8
```

Exercises

- Load `family-7-elec.als` in Alloy
- Execute it
- Look at the generated instance
- Does it look correct?
- What if anything would you change about it?
- Check each of the given assertions
- Are they all valid?
- If not, how would you change the model to fix that?

Exercises

- Load `dynamic/trash-1-elec.als` in Alloy 5
- Complete the model as instructed there
- Execute it
- Check each of the assertions you have written
- Are they all valid?
- If not, how would you change the model to fix that?