CS:5810
Formal Methods in Software Engineering

Introduction to Alloy
Part 1

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Outline

• Introduction to basic Alloy constructs using a simple example of a static model
  – How to define sets, subsets, relations with multiplicity constraints
  – How to use Alloy’s quantifiers and predicate forms

• Basic use of the Alloy Analyzer 4 (AA)
  – Loading, compiling, and analyzing a simple Alloy specification
  – Adjusting basic tool parameters
  – Using the visualization tool to view instances of models
Roadmap

Alloy: Rationale and Use Strategies

– What types of systems have been modeled with Alloy
– What types of questions can AA answer
– What is the purpose of each of the sections in an Alloy specification

Alloy Specifications

– Parameterized conditionals
– Indexed relations
– Graphical representations of Alloy models
– More complex examples
Alloy --- Why was it created?

Lightweight

small and easy to use, and capable of expressing common properties tersely and naturally

Precise

having a simple and uniform mathematical semantics

Tractable

amenable to efficient and fully automated semantic analysis (within scope limits)
Alloy --- Comparison

**UML**

- Has similarities (graphical notation, OCL constraints) but it is neither lightweight, nor precise
- UML includes many modeling notions omitted from Alloy (use-cases, state-charts, code architecture specs)
- Alloy’s diagrams and relational navigation are inspired by UML

**Z**

- Precise, but intractable. Stylized typography makes it harder to work with.
- Z is more expressive than Alloy, but more complicated
- Alloy’s set-based semantics is inspired by Z
Alloy --- What is it used for?

Alloy is a textual modeling language aimed at expressing structural and behavioral properties of software systems.

It is not meant for modeling code architecture (à la class diagrams in UML).

But there might be a close relationship between the Alloy specification and an implementation in an OO language.
Alloy --- Example Applications

The Alloy 4 distribution comes with several example models that together illustrate the use of Alloy’s constructs.

Examples

- Specification of a distributed spanning tree
- Model of a generic file system
- Model of a generic file synchronizer
- Tower of Hanoi model
- ...
Alloy in General

Alloy is general enough that it can model

– any domain of individuals and
– relations between them

We will then start with a few simple examples that are not necessarily about software
Example: Family Structure

We want to...

• Model parent/child relationships as primitive relations
• Model spousal relationships as primitive relations
• Model relationships such as “siblings” as derived relations
• Enforce certain biological constraints via 1st-order predicates (e.g., people have only one mother)
• Enforce certain social constraints via 1st-order predicates (e.g., a wife isn’t a sibling)
• Confirm or refute the existence of certain derived relationships (e.g., no one has a wife with whom he shares a parent)
Example: addressBook

An **address book** for an email client that maintains a mapping from **names** to **addresses**

<table>
<thead>
<tr>
<th>FriendBook</th>
<th>WorkBook</th>
</tr>
</thead>
<tbody>
<tr>
<td>Ted -&gt; <a href="mailto:ted@gmail.com">ted@gmail.com</a></td>
<td>Pilard -&gt; <a href="mailto:lpilard@uiowa.edu">lpilard@uiowa.edu</a></td>
</tr>
<tr>
<td>Ryan -&gt; <a href="mailto:ryan@hotmail.com">ryan@hotmail.com</a></td>
<td>Ryan -&gt; <a href="mailto:ryan@uiowa.edu">ryan@uiowa.edu</a></td>
</tr>
</tbody>
</table>
Atoms and Relations

In Alloy, everything is built from atoms and relations.

An atom is a primitive entity that is

- *indivisible*: it cannot be broken down into smaller parts
- *immutable*: its properties do not change over time
- *uninterpreted*: it does not have any built in property
  (the way numbers do for example)

A relation is a structure that relates atoms. It is a set of tuples, each tuple being a sequence of atoms
Atoms and Relations: Examples

• **Unary relations**: a set of names, a set of addresses and a set of books
  
  \[
  \begin{align*}
  \text{Name} &= \{(N0),(N1),(N2)\} \\
  \text{Addr} &= \{(D0),(D1)\} \\
  \text{Book} &= \{(B0),(B1)\}
  \end{align*}
  \]

• A **binary relation** from names to addresses
  
  \[
  \text{address} = \{(N0,D0),(N1,D1)\}
  \]

• A **ternary relation** from books to name to addresses
  
  \[
  \text{addr} = \{(B0,N0,D0),(B0,N1,D1),(B1,N1,D2)\}
  \]
Relations

Size of a relation: the number of tuples in the relation

Arity of a relation: the number of atoms in each tuple of the relation

relations with arity 1, 2, and 3 are said to be unary, binary, and ternary relations

Examples.

– relation of arity 1 and size 1: myName = \{(N0)\}
– relation of arity 2 and size 3: address = \{(N0,D0),(N1,D1),(N2,D1)\}
Main Components of Alloy Models

• Signatures and Fields
• Predicates and Functions
• Facts
• Assertions
• Commands and scopes
Signatures and Fields

Signatures
- Describe classes of entities we want to reason about
- Sets defined in signatures are fixed (dynamic aspects can be modeled by time-dependent relations)

Fields
- Define relations between signatures

Simple constraints
- Multiplicities on signatures
- Multiplicities on relations
Signatures

• A signature introduces a set of atoms

• The declaration

\[
\text{sig } A \ {} \\
\text{introduces a set named } A
\]

• A set can be introduced as an extension of another; thus

\[
\text{sig } A_1 \text{ extends } A \ {} \\
\text{introduces a set } A_1 \text{ that is a subset of } A
\]
Signatures

sig A {}
sig B {}
sig A1 extends A {}
sig A2 extends A {}

• **A1** and **A2** are extensions of **A**
• A signature declared independently of any other one is a **top-level signature**, e.g., **A** and **B**
• Extensions of the same signature are **mutually disjoint**, as are top-level signatures
Signatures

abstract sig A {}
sig B {}
sig A1 extends A {}
sig A2 extends A {}

• A signature can be introduced as a subset of another
  sig A3 in A {}

• An abstract signature has no elements except those belonging to its extensions or subsets

• All extensions of an abstract signature A form a partition of A
Fields

- **Relations** are declared as *fields* of signatures
  - Writing
    ```
    sig A {f: e}
    ```
    introduces a relation $f$ of type $A \times e$, where $e$ is an expression denoting a product of signatures.

- **Examples**: (with signatures $A$, $B$, $C$)
  - Binary Relation:
    ```
    sig A { f1: B }
    ```
    // $f1$ is a subset of $A \times B$
  - Ternary Relation:
    ```
    sig A { f2: B -> C }
    ```
    // $f2$ is a subset of $A \times B \times C$
Example Signatures and Fields

Family Structure:

abstract sig Person {
    children: Person,
    siblings: Person
}

sig Man, Woman extends Person {}

sig Married in Person {
    spouse: Married
}
Example: Family Structure

**Alloy Model**

```
abstract sig Person {}
sig Man extends Person {}
sig Woman extends Person {}
sig Married in Person {}
```

**Graphical Representation**

- Abstract class **Person**
- **Man** extends Person
- **Woman** extends Person
- **Married** in Person

The diagram shows the class hierarchy with `Person` as the root, `Man` and `Woman` as subclasses, and `Married` as a specialization within the `Person` class.
ModelInstances

The Alloy Analyzer will generate instances of models so that we can see if they match our intentions. Which of the following are instances of our current model?

abstract sig Person {}
sig Man extends Person {}
sig Woman extends Person {}
sig Married in Person {}

A. Person = {(P0),(P1),(P2)}
   Man = {(P1),(P2)}
   Married = {}
   Woman = {(P0)}

B. Person = {(P0),(P1),(P2)}
   Man = {(P1),(P2)}
   Married = {}
   Woman = {(P0),(P1)}

C. Person = {(P0),(P1),(P2),(P3)}
   Man = {(P0),(P1),(P2),(P3)}
   Married = {{(P2),(P3)}}
   Woman = {}

D. Person = {(P0),(P1)}
   Man = {(P0)}
   Married = {(P1)}
   Woman = {}

E. Person = {(P0),(P1)}
   Man = {(P0)}
   Married = {(P1)}
   Woman = {(P1)}
Example: Family Structure

Alloy Model with siblings

```alloy
abstract sig Person {
  siblings: Person
}
sig Man extends Person {}
sig Woman extends Person {}
sig Married in Person {}
```

siblings is a binary relation
it is a subset of Person x Person

Example instance

- Person = {(P0), (P1)}
- Man = {(P0), (P1)}
- Married = {}
- Woman = {}

siblings = {(P0,P1), (P1,P0)}

Intuition: P0 and P1 are siblings
Multiplicities

Allow us to constrain the sizes of sets

- A multiplicity keyword placed before a signature declaration constraints the number of element in the signature’s set
  \[
  \text{m sig A \{\}}
  \]

- We can also make multiplicities constraints on fields:
  \[
  \text{sig A \{f: m e\}}
  \]
  \[
  \text{sig A \{f: e1 m -> n e2\}}
  \]

There are four multiplicities

- \text{set} : any number
- \text{some} : one or more
- \text{lone} : zero or one
- \text{one} : exactly one
Multiplicities: Examples

Without multiplicity:

A set of colors, each of which is red, yellow or green

abstract sig Color {}
sig Red, Yellow, Green extends Color {}
Multiplicities: Examples

• A file system in which each directory contains any number of objects, and each alias points to exactly one object.

```plaintext
abstract sig Object {}
sig Directory extends Object {contents: set Object}
sig File extends Object {}
sig Alias in File {to: one Object}
```

• The default multiplicity is **one**, so:

```plaintext
sig A {f: e} and sig A {f: one e}
```

are equivalent.
Multiplicities: Examples

• A book maps names to addresses
  – There is at most one address per Name
  – An address is associated to at least one name

```
sig Name, Addr {}
sig Book {
    addr: Name some -> lone Addr
}
```
Multiplicities: Examples

- A collection of weather forecasts, each of which has a field `weather` associating every city with exactly one weather condition

  ```
  sig Forecast {weather: City -> one Weather}
  sig City {}
  abstract sig Weather {}
  one sig Rainy, Sunny, Cloudy extends Weather {}
  ```

- Instance:

  ```
  City = {(Iowa City), (Chicago)}
  Rainy = {(rainy)}
  Sunny = {(sunny)}
  Cloudy = {(cloudy)}
  Forecast = {((f1), (f2))}
  weather = { (f1, Iowa City, rainy), (f1, Chicago, rainy),
               (f2, Iowa City, rainy), (f2, Chicago, sunny) }
  ```
Multiplicities and Binary Relations

• \textbf{sig } S \{ f: \text{ lone } T \}

  – says that, for each element \( s \) of \( S \), \( f \) maps \( s \) to \textbf{at most} a single value in \( T \)

Conventional name: \textbf{partial function}

• Potential instances:
Multiplicities and Binary Relations

- **sig $S \{f: \text{one } T\}$**
  - says that, for each element $s$ of $S$, $f$ maps $s$ to exactly one value in $T$

  *Conventional name: total function*

- Potential instances:

  - $s_1 \rightarrow t_1$
  - $s_2 \rightarrow t_2$
  - $s_3 \rightarrow t_3$
  - $s_4 \rightarrow t_4$

  - $s_1 \rightarrow t_1$  [Crossed out]
  - $s_2 \rightarrow t_2$
  - $s_3 \rightarrow t_3$
  - $s_4 \rightarrow t_4$

  - $s_1 \rightarrow t_1$
  - $s_2 \rightarrow t_2$
  - $s_3 \rightarrow t_3$
  - $s_4 \rightarrow t_4$

  - $s_1 \rightarrow t_1$  [Crossed out]
  - $s_2 \rightarrow t_2$
  - $s_3 \rightarrow t_3$
  - $s_4 \rightarrow t_4$

  - $s_1 \rightarrow t_1$
  - $s_2 \rightarrow t_2$
  - $s_3 \rightarrow t_3$
  - $s_4 \rightarrow t_4$

  - $s_1 \rightarrow t_1$  [Crossed out]
  - $s_2 \rightarrow t_2$
  - $s_3 \rightarrow t_3$
  - $s_4 \rightarrow t_4$

  - $s_1 \rightarrow t_1$
  - $s_2 \rightarrow t_2$
  - $s_3 \rightarrow t_3$
  - $s_4 \rightarrow t_4$
Multiplicities and Ternary Relations

- \textbf{sig} \, S \, \{ f: T \rightarrow \text{one} \, V \} \\
  \quad \text{For each element} \, s \, \text{of} \, S, \, \text{over the triples that start with} \, s: \, f \, \text{maps each} \, T\text{-element to exactly one} \, V\text{-element}

- Potential instances:
Multiplicities and Ternary Relations

• \texttt{sig }S \{ f : \texttt{T lone } \rightarrow \texttt{V}\} \\
  – For each element \texttt{s} of \texttt{S}, over the triples that start with \texttt{s}: \texttt{f} maps at most one \texttt{T}-element to the same \texttt{V}-element

• Potential instances:
Multiplicities and Relations

• Other kinds of relational structures can be specified using multiplicities

• Examples:
  – \texttt{sig S \{f: some T\}} ... total relation
  – \texttt{sig S \{f: set T\}} ... partial relation
  – \texttt{sig S \{f: T set -> set V\}}
  – \texttt{sig S \{f: T one -> V\}}
  – ...

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Cardinality Constraints

Multiplicities can also be applied to whole expressions denoting relations

– **some** e  \( e \) is non-empty
– **no** e  \( e \) is empty
– **lone** e  \( e \) has at most one tuple
– **one** e  \( e \) has exactly one tuple
Example: Family Structure

• How would you use multiplicities to define the children relation?

  \[\text{sig}\ \text{Person}\ \{\text{children}: \text{set}\ \text{Person}\}\]

  – Intuition: each person has zero or more children

• How would you use multiplicities to define the spouse relation?

  \[\text{sig}\ \text{Married}\ \{\text{spouse}: \text{one}\ \text{Married}\}\]

  – Intuition: each married person has exactly one spouse
Summarizing

Alloy Model

abstract sig Person {
    children: set Person,
    siblings: set Person
}

sig Man, Woman extends Person {}

sig Married in Person {
    spouse: one Married
}
Exercises

• Start the Alloy Analyzer:
• Load file `family-1.als` from the Resources section of the course website
• Execute it
• Analyze the model instance
• Look at the generated instance
• Does it look correct?
• What, if anything, would you change about it?
Model Instance

Instance found:

Person = {Man0, Man1, Man2}
Man = {Man0, Man1, Man2}
Woman = {}
Married = {Man0, Man1, Man2}

children = { (Man0, Man0), (Man0, Man1),
             (Man1, Man0),
             (Man2, Man1), (Man2, Man2) }

siblings = { (Man0, Man0), (Man0, Man1),
              (Man1, Man0), (Man1, Man2),
              (Man2, Man2) }

spouse = { (Man1, Man0), (Man0, Man2), (Man2, Man0) }
Man can be his own child?

Instance found:

Person = \{\text{Man0, Man1, Man2}\}
Man = \{\text{Man0, Man1, Man2}\}
Woman = \{\}
Married = \{\text{Man0, Man1, Man2}\}

children = \{ (\text{Man0, Man0}), (\text{Man0, Man1}), (\text{Man1, Man0}), (\text{Man2, Man1}), (\text{Man2, Man2}) \}
siblings = \{ (\text{Man0, Man0}), (\text{Man0, Man1}), (\text{Man1, Man0}), (\text{Man1, Man2}), (\text{Man2, Man2}) \}

spouse = \{ ((\text{Man1, Man0}), (\text{Man0, Man2}), (\text{Man2, Man0})) \}
Multiple Fathers?

Instance found:

Person = {Man0, Man1, Man2}
Man = {Man0, Man1, Man2}
Woman = {}
Married = {Man0, Man1, Man2}

children = { (Man0, Man0), (Man0, Man1),
             (Man1, Man0),
             (Man2, Man1), (Man2, Man2)
}
siblings = { (Man0, Man0), (Man0, Man1),
             (Man1, Man0), (Man1, Man2),
             (Man2, Man2)
}
spouse = {(Man1, Man0), (Man0, Man2), (Man2, Man0)}
Own-Siblings?

Instance found:

Person = {Man0, Man1, Man2}
Man = {Man0, Man1, Man2}
Woman = {}
Married = {Man0, Man1, Man2}

children = { (Man0, Man0), (Man0, Man1),
              (Man1, Man0),
              (Man2, Man1), (Man2, Man2) }

siblings = { (Man0, Man0), (Man0, Man1),
               (Man1, Man0), (Man1, Man2),
               (Man2, Man2) }

spouse = { (Man1, Man0), (Man0, Man2), (Man2, Man0) }
Asymmetric Siblings?

Instance found:

Person = {Man0, Man1, Man2}
Man = {Man0, Man1, Man2}
Woman = {}
Married = {Man0, Man1, Man2}

children = { (Man0,Man0), (Man0,Man1),
             (Man1,Man0),
             (Man2,Man1), (Man2,Man2) }
siblings = { (Man0,Man0), (Man0,Man1),
             (Man1,Man0), (Man1,Man2), (Man2,Man2),
             (Man2,Man2) }  No (Man2,Man1)?

spouse = {(Man1,Man0), (Man0,Man2), (Man2,Man0)}
Children-Siblings?

Instance found:

Person = \{\text{Man}0, \text{Man}1, \text{Man}2\}
Man = \{\text{Man}0, \text{Man}1, \text{Man}2\}
Woman = \{
\}
Married = \{\text{Man}0, \text{Man}1, \text{Man}2\}

children = \{(\text{Man}0, \text{Man}0), (\text{Man}0, \text{Man}1), (\text{Man}1, \text{Man}0), (\text{Man}2, \text{Man}1), (\text{Man}2, \text{Man}2)\}

siblings = \{(\text{Man}0, \text{Man}0), (\text{Man}0, \text{Man}1), (\text{Man}1, \text{Man}0), (\text{Man}1, \text{Man}2), (\text{Man}2, \text{Man}2)\}

spouse = \{((\text{Man}1, \text{Man}0), (\text{Man}0, \text{Man}2), (\text{Man}2, \text{Man}0))\}
Asymmetric Marriage?

Instance found:

Person = \{\text{Man0, Man1, Man2}\}
Man = \{\text{Man0, Man1, Man2}\}
Woman = \{
Married = \{\text{Man0, Man1, Man2}\}

children = \{(\text{Man0, Man0}), (\text{Man0, Man1}),
            (\text{Man1, Man0}),
            (\text{Man2, Man1}), (\text{Man2, Man2})\}

siblings = \{(\text{Man0, Man0}), (\text{Man0, Man1}),
               (\text{Man1, Man0}), (\text{Man1, Man2}),
               (\text{Man2, Man2})\}

spouse = \{(\text{Man1, Man0}), (\text{Man0, Man2}), (\text{Man2, Man0})\}

where is (\text{Man0, Man1})?
Model Weaknesses

• The model is underconstrained
  – It doesn’t match our domain knowledge
  – We can add constraints to enrich the model

• Under-constrained models are common early on in the development process
  – AA gives us quick feedback on weaknesses in our model
  – We can incrementally add constraints until we are satisfied with it
Adding Constraints

We’d like to enforce the following constraints which are simply matters of biology:

– No person can be their own parent (or more generally, their own ancestor)

– No person can have more than one father or mother

– A person’s siblings are those with the same parents
Adding Constraints

• We’d like to enforce the following social constraints

  – *The spouse relation is symmetric*

  – *A man’s wife cannot be one of his siblings*