

## Warm-up Analysis of Luby's Distributed Graph Coloring Algorithm

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Assume that  $P(v)$  is initially set to  $\{1, 2, \dots, 2\Delta\}$  for all nodes  $v$ . This assumption is just for this “warm-up” analysis; in the actual algorithm,  $P(v) = \{1, 2, \dots, \text{degree}(v) + 1\}$  initially for all nodes  $v$ . Note that this assumption implies that  $|P(v)| \geq \Delta$  for all nodes  $v$ , *throughout* the algorithm. Fix a round  $i$  and let  $u$  be an arbitrary node that has not been colored after the first  $i - 1$  rounds. Let  $X_u$  denote the event that  $u$  is colored in round  $i$ . We will show that  $\Pr(X_u) \geq 1/4$ .

Let  $W_{c,u}$  denote the event that node  $u$  has selected color  $c$ , as a “tentative” color. Note that  $\Pr(W_{c,u}) = 1/|P(u)|$  if  $c \in P(u)$  and otherwise  $\Pr(W_{c,u}) = 0$ .

$$\begin{aligned}
 \Pr(X_u) &= \Pr(\exists c \in P(u) : W_{c,u} \wedge \forall v \in N(u) : \overline{W_{c,v}}) \\
 &= \sum_{c \in P(u)} \Pr(W_{c,u} \wedge \forall v \in N(u) : \overline{W_{c,v}}) \\
 &= \sum_{c \in P(u)} \Pr(W_{c,u}) \cdot \Pr(\forall v \in N(u) : \overline{W_{c,v}}) \\
 &= \frac{1}{|P(u)|} \sum_{c \in P(u)} \Pr(\forall v \in N(u) : \overline{W_{c,v}}) \\
 &= \frac{1}{|P(u)|} \sum_{c \in P(u)} \prod_{v \in N(u)} \Pr(\overline{W_{c,v}}) \tag{1}
 \end{aligned}$$

Now we focus on finding a lower bound for  $\Pr(\overline{W_{c,v}})$ . If  $c \notin P(v)$  then  $\Pr(\overline{W_{c,v}}) = 1$ . If  $c \in P(v)$  then  $\Pr(\overline{W_{c,v}}) = 1 - 1/|P(v)|$ . Therefore, independent of whether  $c \in P(v)$ , we see that  $\Pr(\overline{W_{c,v}}) \geq (1 - 1/|P(v)|)$ . Since  $|P(v)| \geq \Delta$ , this implies that  $\Pr(\overline{W_{c,v}}) \geq (1 - 1/\Delta)$ . Plugging this lower bound into the right hand side of (1) we get

$$\Pr(X_u) \geq \frac{1}{|P(u)|} \sum_{c \in P(u)} \prod_{v \in N(u)} \left(1 - \frac{1}{\Delta}\right). \tag{2}$$

Since  $|N(u)| \leq \Delta$ ,

$$\prod_{v \in N(u)} \left(1 - \frac{1}{\Delta}\right) \geq \left(1 - \frac{1}{\Delta}\right)^\Delta.$$

Plugging this into (2), we get

$$\Pr(X_u) \geq \frac{1}{|P(u)|} \sum_{c \in P(u)} \left(1 - \frac{1}{\Delta}\right)^\Delta = \left(1 - \frac{1}{\Delta}\right)^\Delta.$$

Now we use the fact that  $(1 - 1/x)^x \geq 1/4$  for all  $x \geq 2$  to obtain the result  $\Pr(X_u) \geq 1/4$  for  $\Delta \geq 2$ .