1. Consider the edge-weighted graph shown below. We have partially run some MST algorithm and it has determined that edges $\{A, B\}$, $\{B, H\}$, $\{G, F\}$, and $\{D, E\}$ belong to an MST. Using the *cut property* of MSTs figure out all the edges that can safely be added to this partial MST. In other words, let $X = \{\{A, B\}, \{B, H\}, \{G, F\}, \{D, E\}\}$ and your task is to write down all edges $e \notin X$ such $X \cup \{e\}$ is part of some MST. For each edge e that you write down, also write down the corresponding cut (S, V - S) that helped you decide that e is safe.



Solution: There are three edges e, each of which can be added to X, such that $X \cup \{e\}$ is still a part of the MST. Below, for each edge e, I provide cuts (S, V - s) such that e is a lightest edge crossing this cut.

- $\begin{array}{l} e = \{B,C\}; \, S = \{A,B,H\}.\\ e = \{E,F\}; \, S = \{F,G\} \mbox{ or } S = \{A,B,F,G,H\}.\\ e = \{C,D\}; \, S = \{C\} \mbox{ or } S = \{A,B,C,F,G,H\}, \mbox{ or } S = \{A,B,C,H\}. \end{array}$
- 2. Consider an input to the SETCOVER problem with |B| = 10000 (i.e., we need to "cover" a set B with 10,000 elements). Suppose that somehow we know that B can be optimally covered by using 20 sets. Further suppose that after 5 iterations of the greedy algorithm for SETCOVER, it has picked sets that cover 7000 elements.
 - (a) Fill in the blank in the following statement:

Based on this information, it is possible to say that in the next iteration, the greedy algorithm will pick a set of size at least ______.

(b) Explain your answer for (a) in 1-2 sentences.

Solution: (a) 150, (b) After 5 iterations of greedy 3000 elements in *B* remain uncovered. The optimal solution covers these elements with 20 sets and therefore there is at least one set among those picked by the optimal solution that covers at least 3000/20 = 150 elements in *B*. Therefore in the next iteration, the greedy algorithm will pick a set that covers at least 150 elements in *B*.