

Quick Sort



MAY 1ST, 2013

Quick Sort: Review



```
def generalQuickSort(L, first, last):  
    # Base case: if first == last, then there is only one element in the  
    # slice that needs sorting. So there is nothing to do.  
  
    # Recursive case: if there are 2 or more elements in the slice L[first:last+1]  
    if first < last:  
        # Divide step: partition returns an index p such that  
        # first <= p <= last and everything in L[first:p] is <= L[p]  
        # and everything in L[p+1:last+1] is >= L[p]  
        p = partition(L, first, last)  
  
        # Conquer step  
        generalQuickSort(L, first, p-1)  
        generalQuickSort(L, p+1, last)  
  
    # Combine step: there is nothing left to do!
```

Quick Sort: Review

The partition function



```
def partition(L, first, last):
    # We pick the element L[first] as the "pivot" around which we partition the list
    p = first

    # We process the rest of the elements, one-by-one, in left-to-right order
    for current in range(p+1, last+1):

        # If L[current] is smaller than the pivot, it needs to move into the first block,
        # to the left of the pivot.
        if L[current] < L[p]:
            swap(L, current, p)
            swap(L, p, p+1)
            p = p + 1

    return p
```

Quick Sort: Review

The wrapper function and the *swap* function



```
def quickSort(L):  
    generalQuickSort(L, 0, len(L)-1)
```

```
def swap(L, i, j):  
    temp = L[i]  
    L[i] = L[j]  
    L[j] = temp
```

Partition Function in Action



- **Initial list:** [6, 2, 4, 1, 6, 10, 2, 11, 8, 7]
- 6 is selected as the pivot.

[] 6 []

[2] 6 []

[2, 4] 6 []

[2, 4, 1] 6 []

[2, 4, 1] 6 [6]

[2, 4, 1] 6 [6, 10]

[2, 4, 1, 2] 6 [10, 6]

[2, 4, 1, 2] 6 [10, 6, 11]

[2, 4, 1, 2] 6 [10, 6, 11, 8]

[2, 4, 1, 2] 6 [10, 6, 11, 8, 7]

- **Final list:** [2, 4, 1, 2, 6, 10, 6, 11, 8, 7]
- Function returns index 4

The partition may be skewed!



- **Initial list:** [16, 2, 4, 1, 6, 10, 2, 11, 8, 7]

```
[ ] 16 [ ]  
[2] 16 [ ]  
[2, 4] 16 [ ]  
[2, 4, 1] 16 [ ]  
[2, 4, 1, 6] 16 [ ]  
[2, 4, 1, 6, 10] 16 [ ]  
[2, 4, 1, 6, 10, 2] 16 [ ]  
[2, 4, 1, 6, 10, 2, 11] 16 [ ]  
[2, 4, 1, 6, 10, 2, 11, 8] 16 [ ]  
[2, 4, 1, 6, 10, 2, 11, 8, 7] 16 [ ]
```

- **Final list:** [2, 4, 1, 6, 10, 2, 11, 8, 7, 16]
- Function returns index 9

Quick Sort in Action



The initial list is: [6, 2, 4, 11, 6, 10, 2]

[2, 4, 2] [6] [6, 10, 11] (partition on [6, 2, 4, 11, 6, 10, 2])

[] [2] [4, 2] (partition on [2, 4, 2])

[2] [4] [] (partition on [4, 2])

[] [6] [10, 11] (partition on [6, 10, 11])

[] [10] [11] (partition on [10, 11])

The sorted list is: [2, 2, 4, 6, 6, 10, 11]

Running Time Comparison



- On lists with 100,000 elements constructed at random. Selection sort took 5 minutes on lists of this size.

Finished constructing the lists...

Time for Merge Sort is: 0.678801059723

Time for Quick Sort is: 0.980933904648

Finished constructing the lists...

Time for Merge Sort is: 0.682029008865

Time for Quick Sort is: 0.987423181534

Finished constructing the lists...

Time for Merge Sort is: 0.67242193222

Time for Quick Sort is: 0.985061883926

Puzzle: A different Experiment



- The input list is [0, 1, 2, ...] of size 10,000.

Finished constructing the lists...

Time for Merge Sort is: 0.0404422283173

Time for Quick Sort is: 4.38273501396

Finished constructing the lists...

Time for Merge Sort is: 0.0395169258118

Time for Quick Sort is: 4.36549711227

Finished constructing the lists...

Time for Merge Sort is: 0.0384669303894

Time for Quick Sort is: 4.36951899529

- Why does quick sort take 100 times more time??

Solution



- For $[0, 1, 2, 3, \dots, n-1]$, n units of work yields
 $[]_0 [1, 2, 3, \dots, n-1]$
- An additional $n-1$ units of work yields
 $[]_1 [2, 3, \dots, n-1]$
- An additional $n-2$ units of work yields
 $[]_2 [3, 4, \dots, n-1]$

So total work is $n + (n-1) + (n-2) + \dots + 1$, which is roughly $n^2/2$.

Ideal behavior



- n units of work yield

[.....] pivot [.....]

- n units of work yield

[.....] pivot [.....] [.....] pivot [.....]

- n units of work yield

[...] pivot [...] [...] pivot [...] [...] pivot [...] [...] pivot [...]

-

We go down $\log(n)$ levels, for a total of $n \log(n)$ units of work.

How to pick a good pivot?



- Randomize! (Just pick a random element as the pivot, instead of the first element).
- Add this line of code at the beginning of partition:
`swap(L, first, random.randint(first, last))`
- Now the running times, even on a sorted input list are comparable:

Finished constructing the lists...

Time for Merge Sort is: 0.040990114212

Time for Quick Sort is: 0.0971350669861