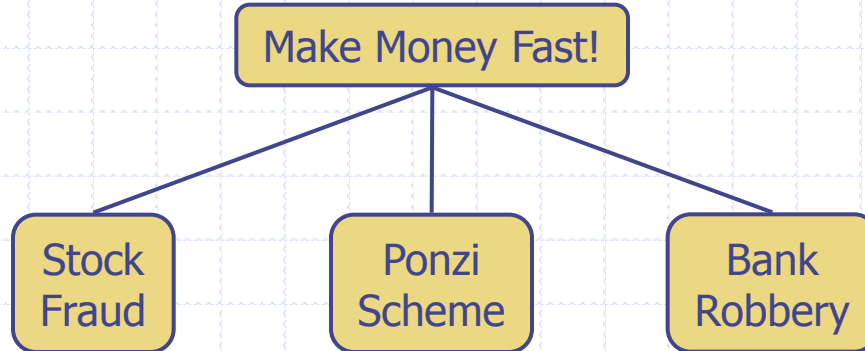
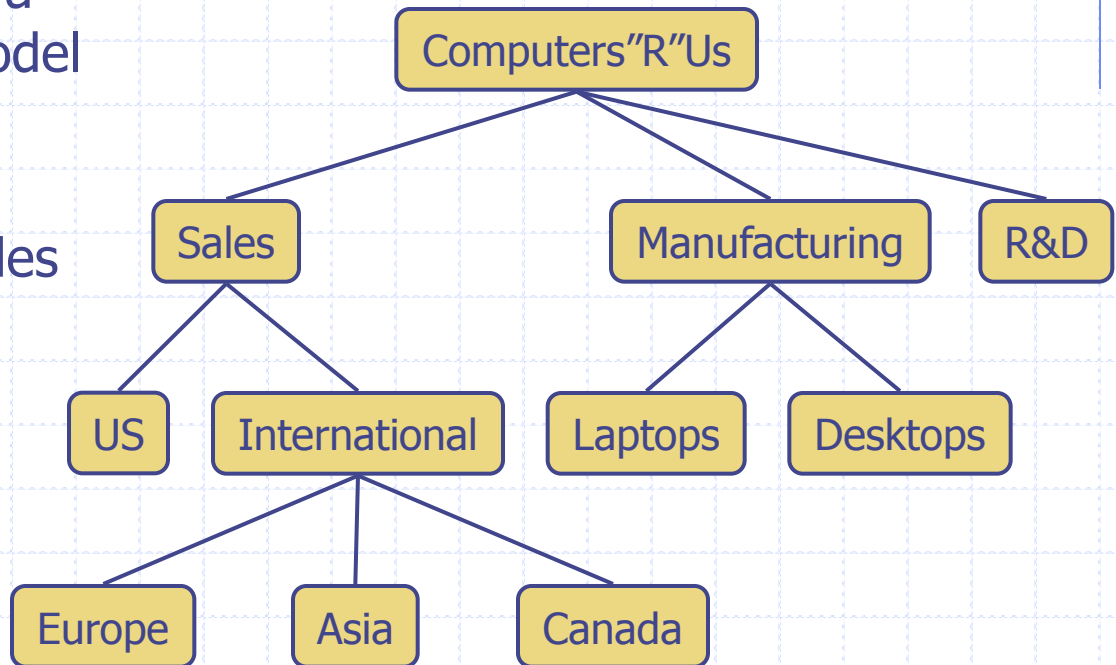


# Trees



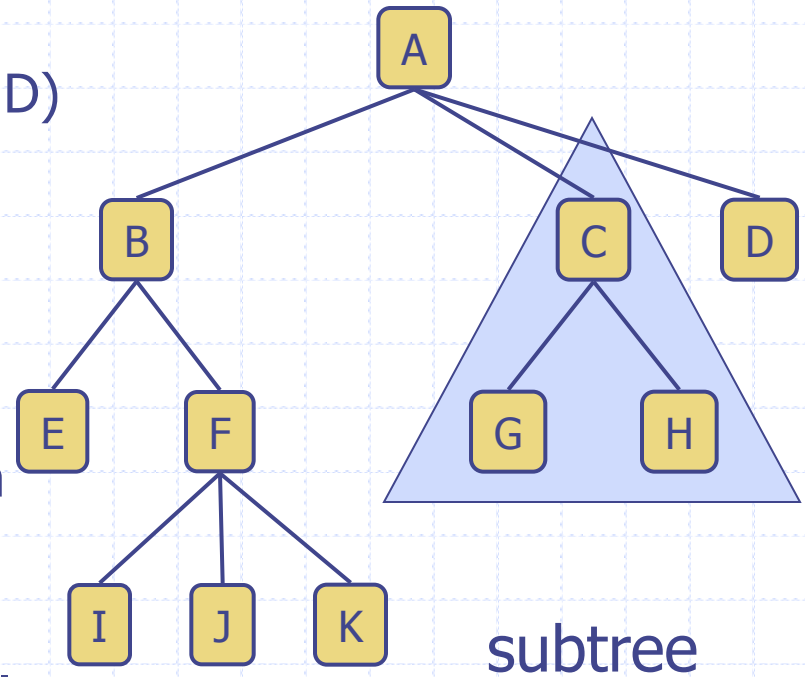
# What is a Tree

- ❑ In computer science, a tree is an abstract model of a hierarchical structure
- ❑ A tree consists of nodes with a parent-child relation
- ❑ Applications:
  - Organization charts
  - File systems
  - Programming environments



# Tree Terminology

- ❑ Root: node without parent (A)
- ❑ Internal node: node with at least one child (A, B, C, F)
- ❑ External node (a.k.a. leaf ): node without children (E, I, J, K, G, H, D)
- ❑ Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- ❑ Depth of a node: number of ancestors
- ❑ Height of a tree: maximum depth of any node (3)
- ❑ Descendant of a node: child, grandchild, grand-grandchild, etc.
- ❑ Subtree: tree consisting of a node and its descendants



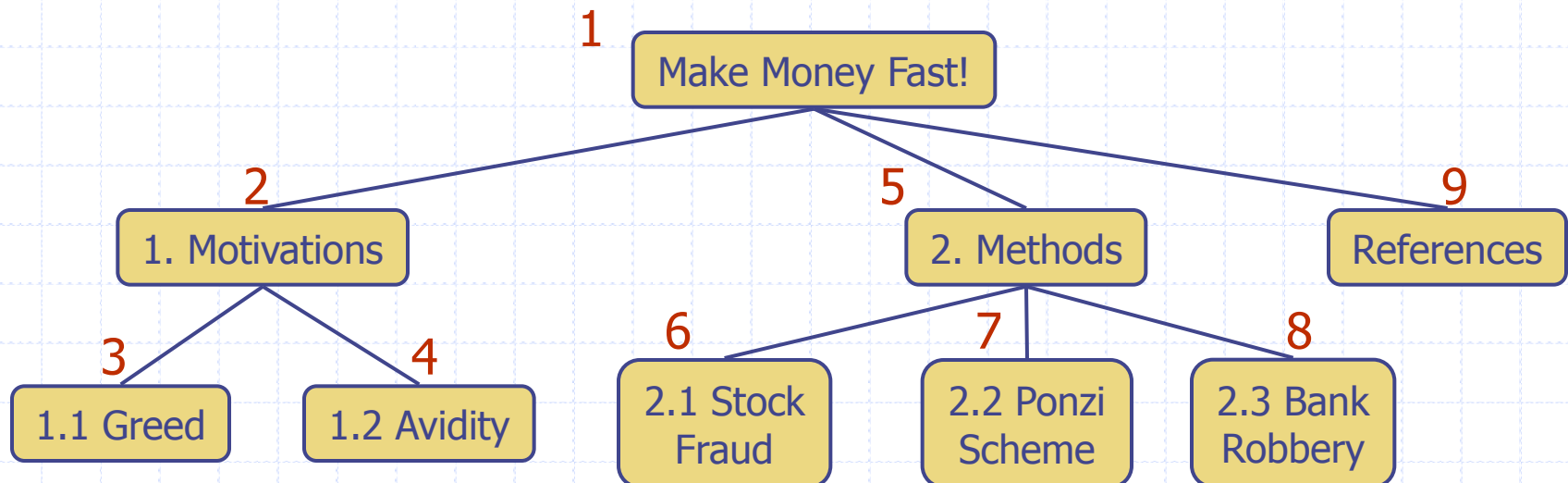
# Tree ADT

- We use positions to abstract nodes
- Generic methods:
  - integer `size()`
  - boolean `isEmpty()`
  - Iterator `iterator()`
  - Iterable `positions()`
- Accessor methods:
  - position `root()`
  - position `parent(p)`
  - Iterable `children(p)`
- ◆ Query methods:
  - boolean `isInternal(p)`
  - boolean `isExternal(p)`
  - boolean `isRoot(p)`
- ◆ Update method:
  - element `replace (p, o)`
- ◆ Additional update methods may be defined by data structures implementing the Tree ADT

# Preorder Traversal

- A traversal visits the nodes of a tree in a systematic manner
- In a preorder traversal, a node is visited before its descendants
- Application: print a structured document

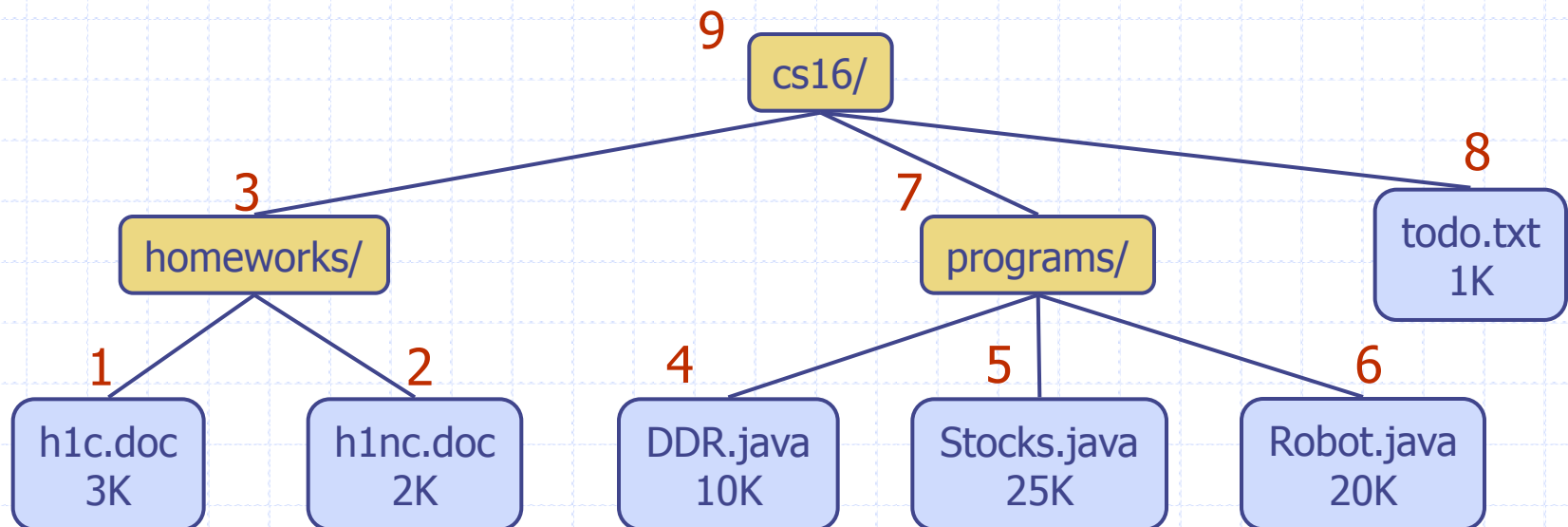
**Algorithm** *preOrder(v)*  
*visit(v)*  
**for each** child *w* of *v*  
*preorder(w)*



# Postorder Traversal

- In a postorder traversal, a node is visited after its descendants
- Application: compute space used by files in a directory and its subdirectories

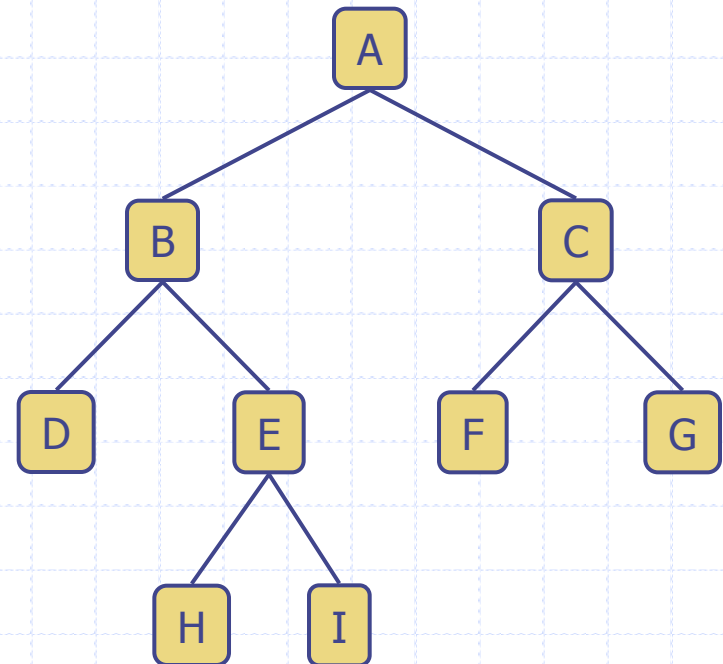
**Algorithm** *postOrder*( $v$ )  
for each child  $w$  of  $v$   
    *postOrder* ( $w$ )  
*visit*( $v$ )



# Binary Trees

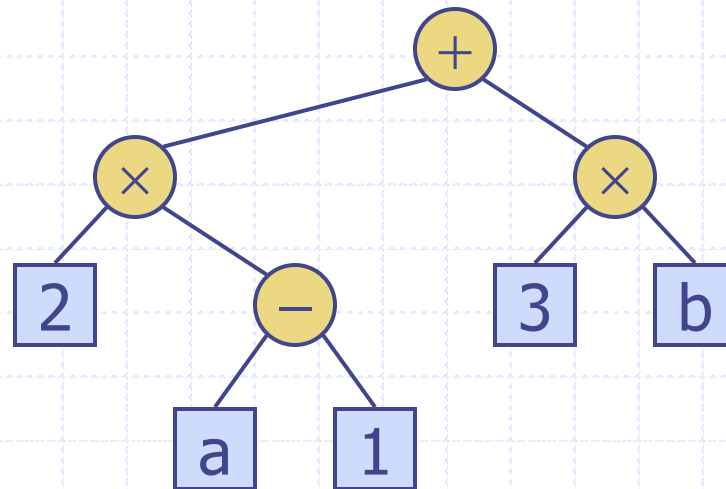
- A binary tree is a tree with the following properties:
  - Each internal node has at most two children (exactly two for **proper** binary trees)
  - The children of a node are an ordered pair
- We call the children of an internal node **left child** and **right child**
- Alternative recursive definition: a binary tree is either
  - a tree consisting of a single node, or
  - a tree whose root has an ordered pair of children, each of which is a binary tree

- Applications:
  - arithmetic expressions
  - decision processes
  - searching



# Arithmetic Expression Tree

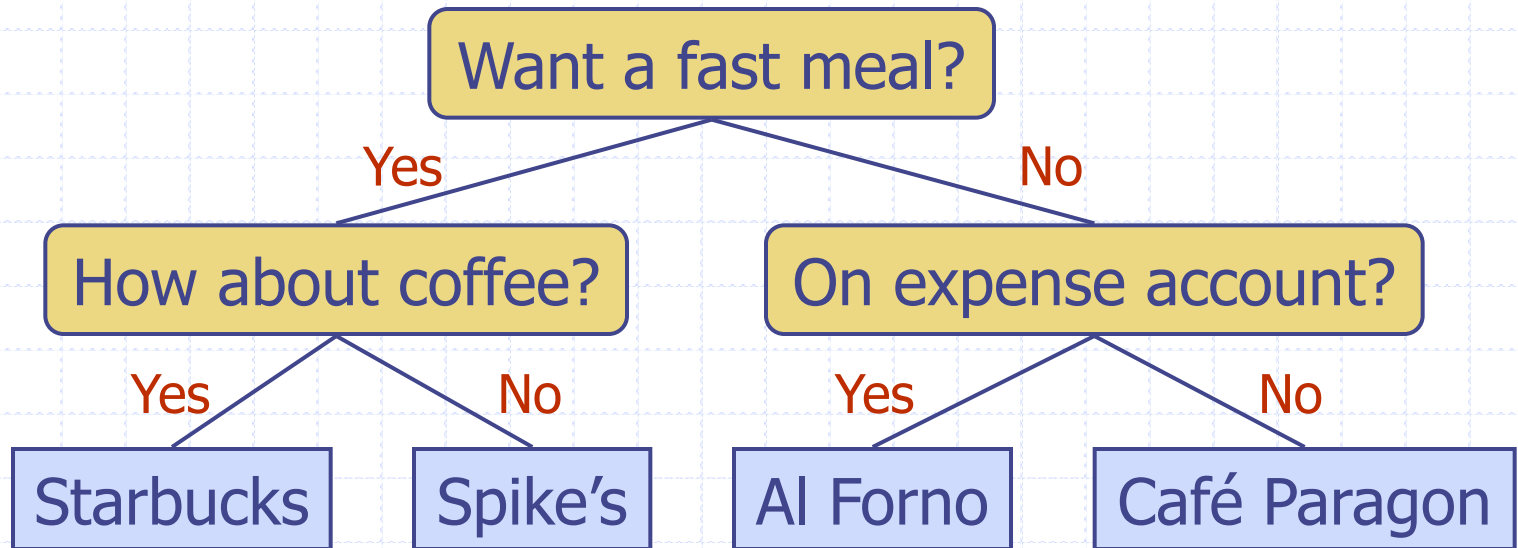
- Binary tree associated with an arithmetic expression
  - internal nodes: operators
  - external nodes: operands
- Example: arithmetic expression tree for the expression  $(2 \times (a - 1) + (3 \times b))$





# Decision Tree

- Binary tree associated with a decision process
  - internal nodes: questions with yes/no answer
  - external nodes: decisions
- Example: dining decision



# Properties of Proper Binary Trees

## □ Notation

$n$  number of nodes

$e$  number of external nodes

$i$  number of internal nodes

$h$  height

## ◆ Properties:

■  $e = i + 1$

■  $n = 2e - 1$

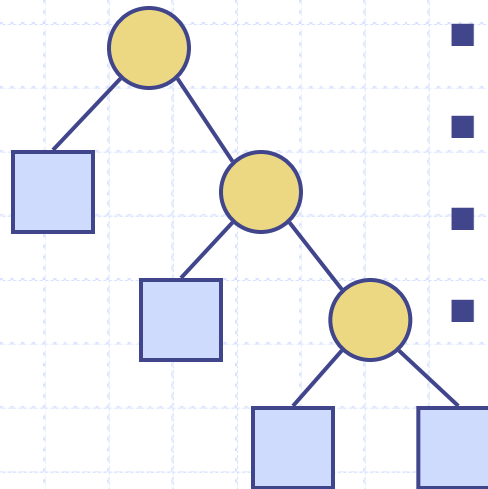
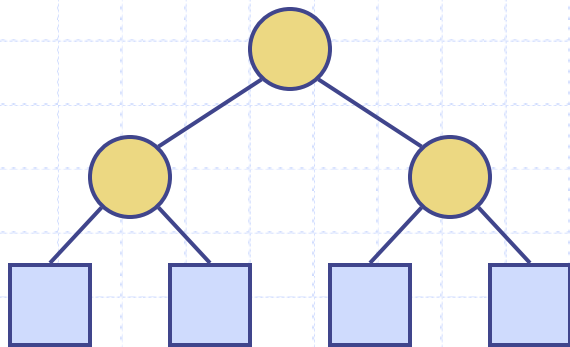
■  $h \leq i$

■  $h \leq (n - 1)/2$

■  $e \leq 2^h$

■  $h \geq \log_2 e$

■  $h \geq \log_2 (n + 1) - 1$



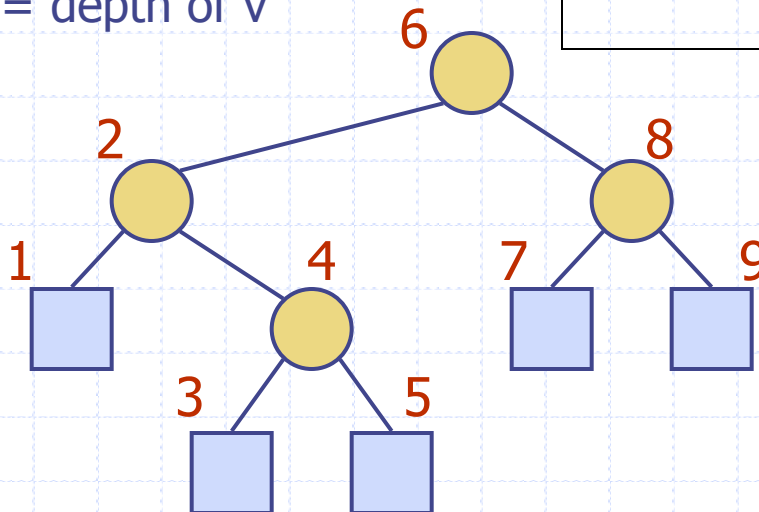
# BinaryTree ADT

- The BinaryTree ADT extends the Tree ADT, i.e., it inherits all the methods of the Tree ADT
- Additional methods:
  - position **left**(p)
  - position **right**(p)
  - boolean **hasLeft**(p)
  - boolean **hasRight**(p)
- Update methods may be defined by data structures implementing the BinaryTree ADT

# Inorder Traversal

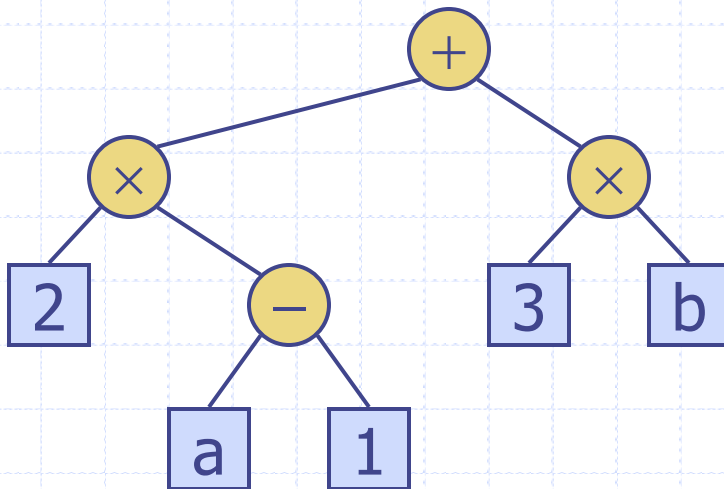
- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Application: draw a binary tree
  - $x(v)$  = inorder rank of  $v$
  - $y(v)$  = depth of  $v$

```
Algorithm inOrder( $v$ )  
  if hasLeft ( $v$ )  
    inOrder (left ( $v$ ))  
  visit( $v$ )  
  if hasRight ( $v$ )  
    inOrder (right ( $v$ ))
```



# Print Arithmetic Expressions

- Specialization of an inorder traversal
  - print operand or operator when visiting node
  - print "(" before traversing left subtree
  - print ")" after traversing right subtree



**Algorithm** *printExpression(v)*

**if** *hasLeft* (v)

*print*("(")

*inOrder* (*left*(v))

*print*(v.*element* ())

**if** *hasRight* (v)

*inOrder* (*right*(v))

*print* (")")

$((2 \times (a - 1)) + (3 \times b))$

# Evaluate Arithmetic Expressions

- Specialization of a postorder traversal
  - recursive method returning the value of a subtree
  - when visiting an internal node, combine the values of the subtrees

**Algorithm** *evalExpr(v)*

**if** *isExternal*(v)

**return** *v.element* ()

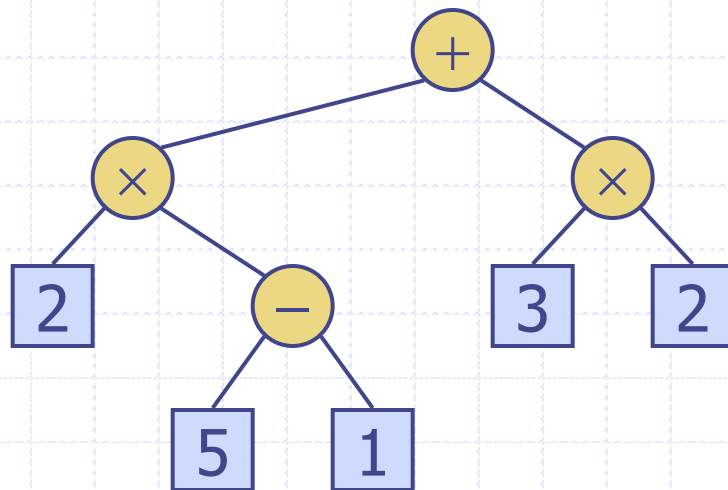
**else**

*x* ← *evalExpr*(*leftChild*(v))

*y* ← *evalExpr*(*rightChild*(v))

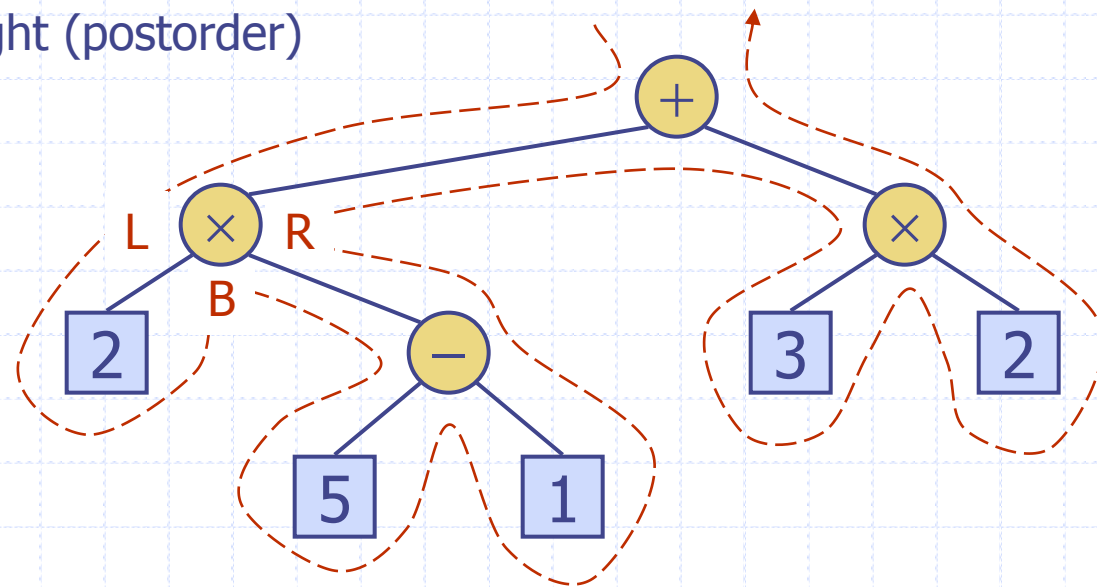
◇ ← operator stored at v

**return** *x* ◇ *y*



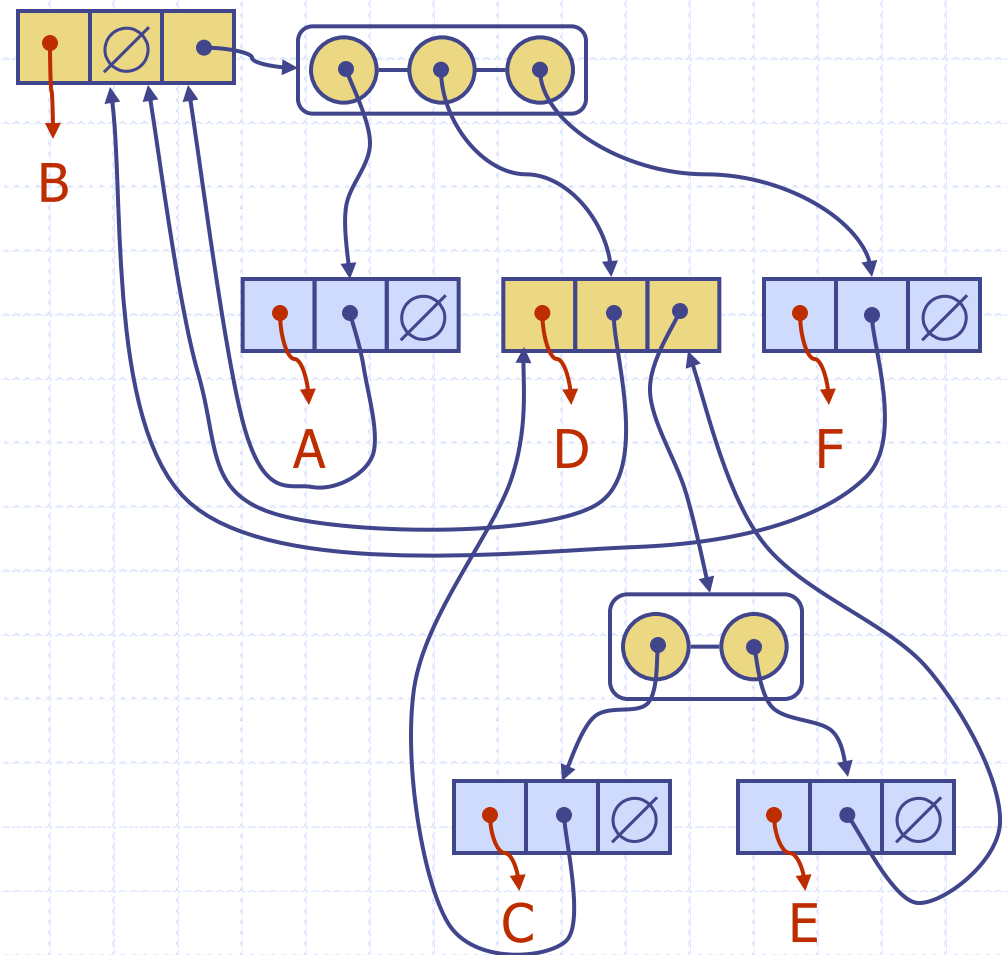
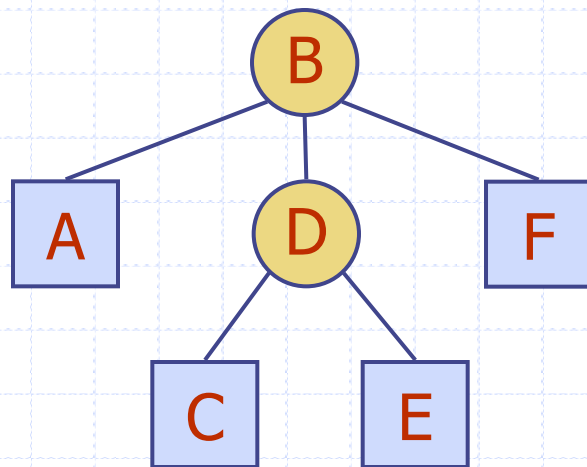
# Euler Tour Traversal

- Generic traversal of a binary tree
- Includes a special cases the preorder, postorder and inorder traversals
- Walk around the tree and visit each node three times:
  - on the left (preorder)
  - from below (inorder)
  - on the right (postorder)



# Linked Structure for Trees

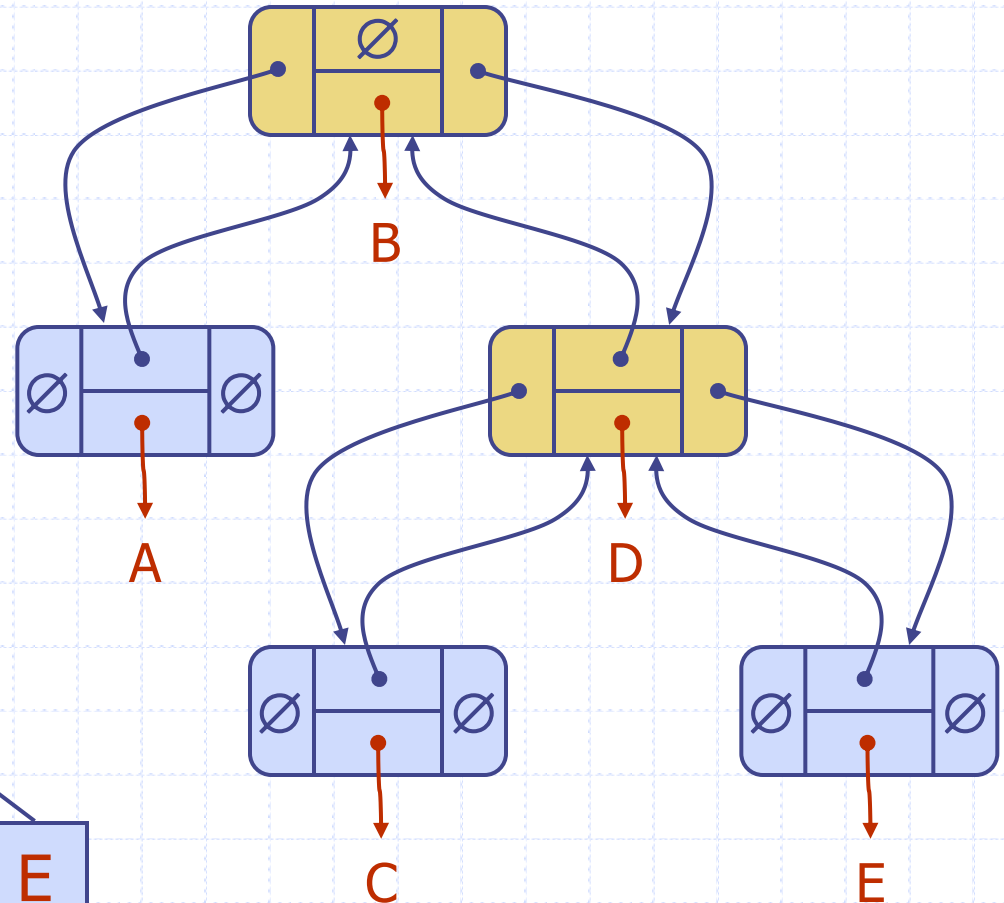
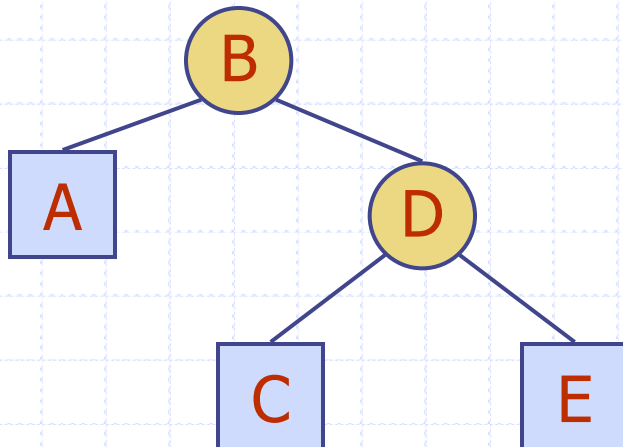
- A node is represented by an object storing
  - Element
  - Parent node
  - Sequence of children nodes
- Node objects implement the Position ADT





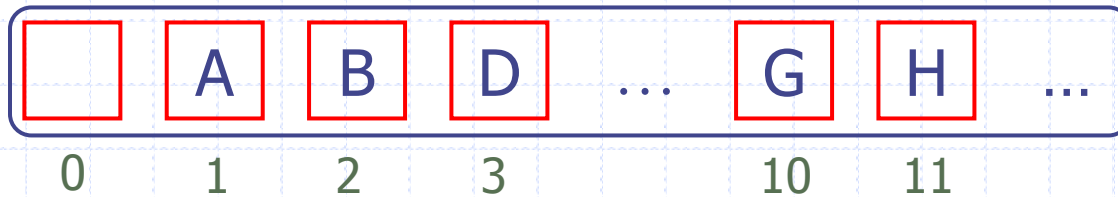
# Linked Structure for Binary Trees

- A node is represented by an object storing
  - Element
  - Parent node
  - Left child node
  - Right child node
- Node objects implement the Position ADT

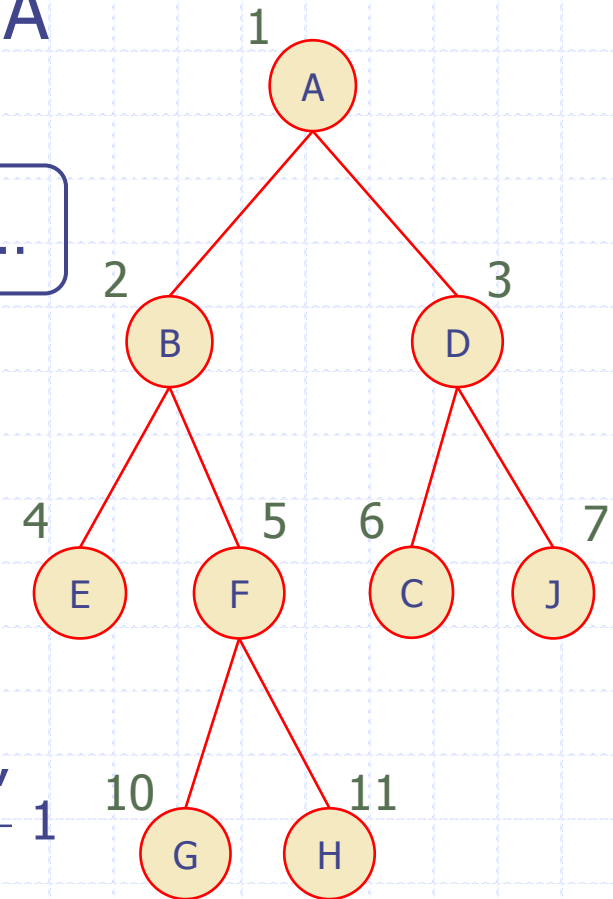


# Array-Based Representation of Binary Trees

- Nodes are stored in an array  $A$



- Node  $v$  is stored at  $A[\text{rank}(v)]$ 
  - $\text{rank}(\text{root}) = 1$
  - if node is the left child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node}))$
  - if node is the right child of  $\text{parent}(\text{node})$ ,  
 $\text{rank}(\text{node}) = 2 \cdot \text{rank}(\text{parent}(\text{node})) + 1$



# Template Method Pattern

- Generic algorithm
- Implemented by abstract Java class
- Visit methods redefined by subclasses
- Template method `eulerTour`
  - Recursively called on left and right children
  - A `TourResult` object with fields `left`, `right` and `out` keeps track of the output of the recursive calls to `eulerTour`

```
public abstract class EulerTour <E, R> {
    protected BinaryTree<E> tree;
    public abstract R execute(BinaryTree<E> T);
    protected void init(BinaryTree<E> T) { tree = T; }
    protected R eulerTour(Position<E> v) {
        TourResult<R> r = new TourResult<R>();
        visitLeft(v, r);
        if (tree.hasLeft(p))
            { r.left=eulerTour(tree.left(v)); }
        visitBelow(v, r);
        if (tree.hasRight(p))
            { r.right=eulerTour(tree.right(v)); }
        return r.out;
    }
    protected void visitLeft(Position<E> v, TourResult<R> r) {}
    protected void visitBelow(Position<E> v, TourResult<R> r) {}
    protected void visitRight(Position<E> v, TourResult<R> r) {}
}
```

# Specializations of EulerTour

- Specialization of class EulerTour to evaluate arithmetic expressions
- Assumptions
  - Nodes store **ExpressionTerm** objects with method **getValue**
  - **ExpressionVariable** objects at external nodes
  - **ExpressionOperator** objects at internal nodes with method **setOperands(Integer, Integer)**

```
public class EvaluateExpressionTour
    extends EulerTour<ExpressionTerm, Integer> {
    public Integer execute
        (BinaryTree<ExpressionTerm> T) {
        init(T);
        return eulerTour(tree.root());
    }
    protected void visitRight
        (Position<ExpressionTerm> v,
         TourResult<Integer> r) {
        ExpressionTerm term = v.element();
        if (tree.isInternal(v)) {
            ExpressionOperator op = (ExpressionOperator) term;
            op.setOperands(r.left, r.right);
            r.out = term.getValue();
        }
    }
}
```