

Distributed Algorithms for Graph coloring

A few known results

- Any **tree** can be colored using **two colors** only
- Any **graph** whose maximum node degree is Δ can be colored using **$(\Delta+1)$** colors
- Any **planar graph** can be colored using four colors, but no distributed algorithm is known and the centralized algorithm is also extremely cumbersome.
- Any **tree** of size **n** can be colored using three colors in **$\log^*(n)$** rounds (Cole and Vishkin)

$O(\log^* n)$ coloring of a tree

Cole & Vishkin's algorithm

$\log^*(n)$ is the *smallest number of log-operations* needed to bring n down to ≤ 2 . For example, let $n = \text{one trillion}$. Now,

$$\log(\text{one trillion}) = 40,$$

$$\log(\log(\text{one trillion})) = 5.322, \text{ and}$$

$$\log(\log(\log(\log(\text{one trillion})))) < 2.$$

This means that $\log^*(\text{one trillion}) = 4$. This also illustrates that the function grows very slowly with the value of the argument.

$O(\log^* n)$ coloring of a tree

Consider a **rooted tree**. The algorithm assumes that initially the color of each node is its id.

Each non-root node **v** is aware of its *parent* **p(v)**.

Interpret each color **c** as a *little-endian* bit string $c_{k-1} c_{k-2} c_{k-3} \dots c_0$, and let **|c|** denote the size of the bit string.

$O(\log^* n)$ coloring of a tree

program *reduce* for a rooted tree

The root node labels itself with color 0 followed by bit 0 of its old color;

{Program for each non-root node v in a round}

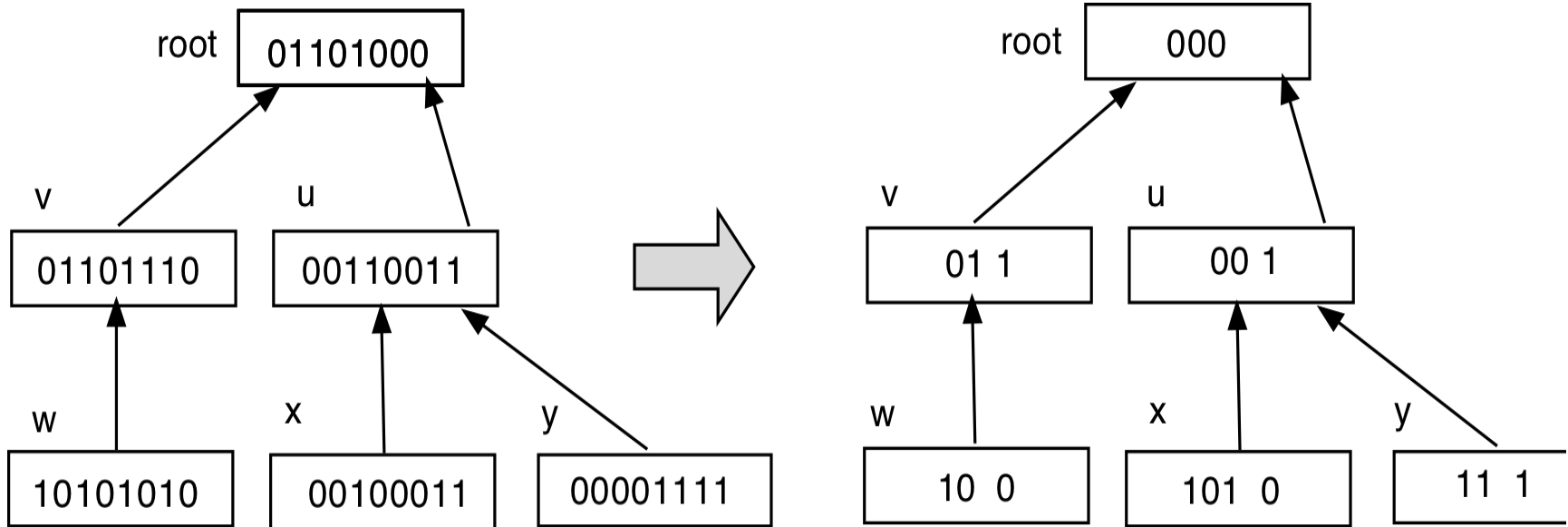
do $c(v) \geq 6 \rightarrow$

{Let j = smallest index where the bit strings of $c(v)$ and $c(p(v))$ differ}

new color $c(v) :=$ the bit string for j followed by bit j of $c(v)$

od

$O(\log^* n)$ coloring of a tree



(a) initial colors

(b) after one step

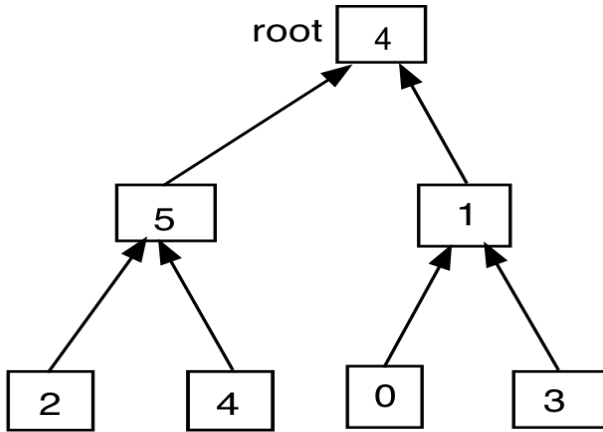
$O(\log^* n)$ coloring of a tree

The shift-down operation

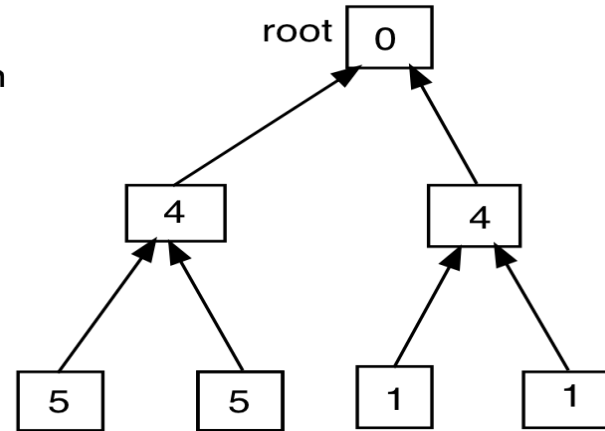
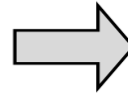
The root picks a new color < 6 and different from its current color, and

Each non-root node v concurrently executes $c(v) := c(p(v))$

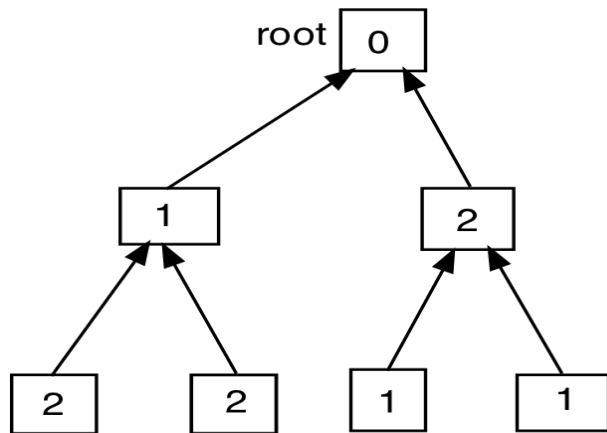
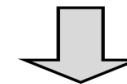
$O(\log^* n)$ coloring of a tree



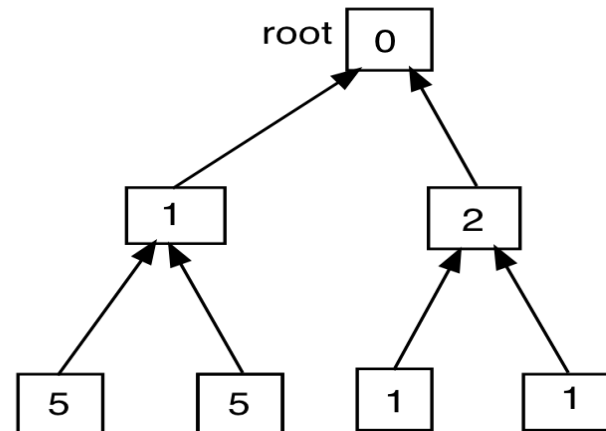
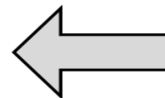
shift-down



eliminate 4



eliminate 5



$O(\log^* n)$ coloring of a tree

Reduction of the color palette size from six to three

$z := 3;$

do $z \leq 5 \rightarrow$

$c(v) = z \rightarrow$ pick a color from $\{0, 1, 2\}$ not used by the neighbors of v ;

$z := z + 1$

od