

Context-sensitive Examples

This document presents three examples of “context-sensitive” grammars. They are drawn from the book *Models of Computation and Formal Languages* by R. G. Taylor.

Example 11.1.2 The language $L = \{a^i b^j c^i d^j \mid i, j \geq 1\}$ is not context-free. There do exist grammars that generate L , however. Here is one:

$$\begin{aligned} S &\rightarrow AB \\ A &\rightarrow aAX \mid aX \\ B &\rightarrow bBd \mid bYd \\ Xb &\rightarrow bX \\ XY &\rightarrow Yc \\ Y &\rightarrow \varepsilon \end{aligned}$$

The logic of the grammar is basically as follows. The productions for nonterminal A will be used to fix the value of i while the productions for nonterminal B will fix the value of j . Each occurrence of nonterminal X is replaced with terminal c but not before the production $Xb \rightarrow bX$ has been used to position it between the b s and d s. Nonterminal Y serves as a placemaker situated between the b s and the d s and can hence be used in positioning the X s. The problem, of course, as the reader will likely have noticed, is that this grammar is not quite context-sensitive given the production for nonterminal Y (editor’s note: Taylor defines a context-sensitive production $\alpha \rightarrow \beta$ to be one where $\text{len}(\alpha) \leq \text{len}(\beta)$). But it is fairly easy to obtain an equivalent grammar that is context-sensitive by adjusting the productions for S .

$$\begin{aligned} \text{(i)-(ii)} \quad S &\rightarrow aAB \mid aB \\ \text{(iii)-(iv)} \quad A &\rightarrow aAX \mid aX \\ \text{(v)-(vi)} \quad B &\rightarrow bBd \mid bYd \\ \text{(vii)} \quad Xb &\rightarrow bX \\ \text{(viii)} \quad XY &\rightarrow Yc \\ \text{(ix)} \quad Y &\rightarrow c \end{aligned}$$

Placemaker Y is now itself being replaced with terminal c , so we do not need quite so many occurrences of nonterminal X . A derivation of word $aabbccddd$ appears below.

$$\begin{aligned} S &\Rightarrow aAB && \text{by (i)} \\ &\Rightarrow aaXB && \text{by (iv)} \\ &\Rightarrow aaXXbBd && \text{by (v)} \\ &\Rightarrow aaXbbBdd && \text{by (v)} \\ &\Rightarrow aaXbbbYddd && \text{by (vi)} \\ &\Rightarrow aabXbbYddd && \text{by (vii)} \\ &\Rightarrow aabbXbYddd && \text{by (vii)} \end{aligned}$$

\Rightarrow aabbbXYddd by (vii)
 \Rightarrow aabbbYcddd by (viii)
 \Rightarrow aabbbccddd by (ix)

It should be apparent that every derivation of a word over $\Sigma = \{a,b,c,d\}$ will terminate with an application of production (ix). Moreover, by the time (ix) is used to eliminate nonterminal Y, all occurrences of nonterminal X must already have been converted to cs; once Y has disappeared, production (viii) will no longer be applicable.

Example 11.1.3 The language $L = \{a^i b^j c^k \mid 1 \leq i \leq j \leq k\}$ is not context-free. Again, there do exist grammars that generate L, however. One such grammar has the 10 productions.

$S \rightarrow aS'bX \mid abX$
 $S' \rightarrow aS'bC \mid S'bC \mid S'C \mid bC \mid C$
 $Cb \rightarrow bC$
 $CX \rightarrow Xc$
 $X \rightarrow c$

The productions for S and S' are first used to fix the values i, j, and k, in effect. Nonterminal X is introduced as a placemaker situated immediately to the right of all the bs. Precisely as in the preceding Example 11.1.2, occurrences of nonterminal C are moved to the right past the bs before being converted to cs. Finally, placemaker X is itself converted to c.

Example 11.1.4 We previously considered the theoretical linguist's use of context-free grammars in giving an account of the syntax of English. We close the present section with a brief discussion of the application within linguistics of context-sensitive rules. As an easy illustration, we describe the potential of such rules to interpret the phenomenon of subject-verb agreement with respect to number—that is, singular or plural—as reflected in sentences:

- (a) The child runs.
- (b) The men run.

To this end, we introduce the following rules (editor's note: syntax categories are S for sentence, NP for noun-phrase, VP for verb-phrase, N for noun, V for verb, and Det for determiner):

$S \rightarrow NP VP$
 $NP \rightarrow Det N_{sing} \mid Det N_{plur}$
 $N_{sing} VP \rightarrow N_{sing} V_{sing}$
 $N_{plur} VP \rightarrow N_{plur} V_{plur}$
 $Det \rightarrow \mathbf{the}$

$N_{\text{sing}} \rightarrow \mathbf{child}$

$N_{\text{plur}} \rightarrow \mathbf{men}$

$V_{\text{sing}} \rightarrow \mathbf{runs}$

$V_{\text{plur}} \rightarrow \mathbf{run}$

Note that the fourth and fifth rules here are context-sensitive but not context-free. Although it is possible to account for the subject-verb agreement using context-free rules only, the two context-sensitive rules capture neatly our intuition that the number of the subject determines that of the verb.