

- HW 7 due tonight
- DS 8 available later today, due Sunday
- HW 8 available tomorrow, due next Thursday

## Last time

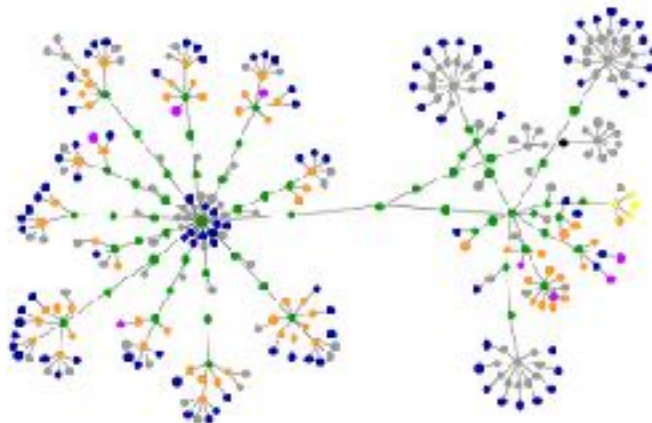
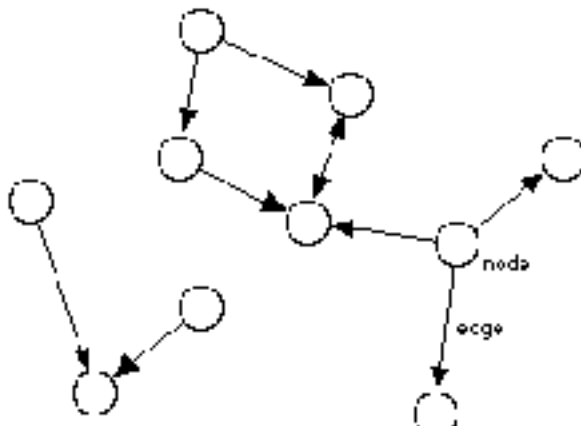
- Intro to optimization algorithms, greedy algorithms
- Introduced graphs (the computer science/mathematical kind), not the charts you've been plotting with pylab

## Today

- Graph representations
- Basic graph algorithms

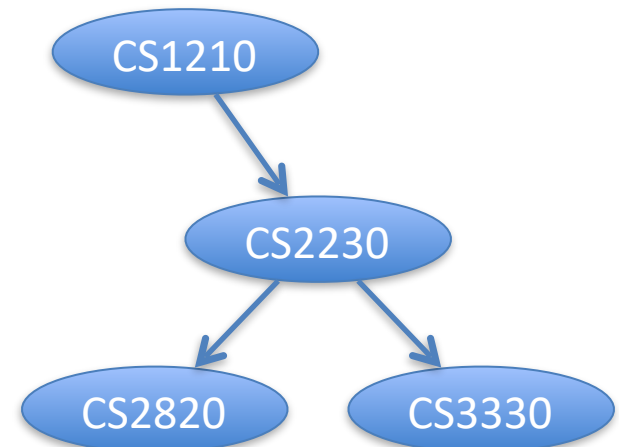
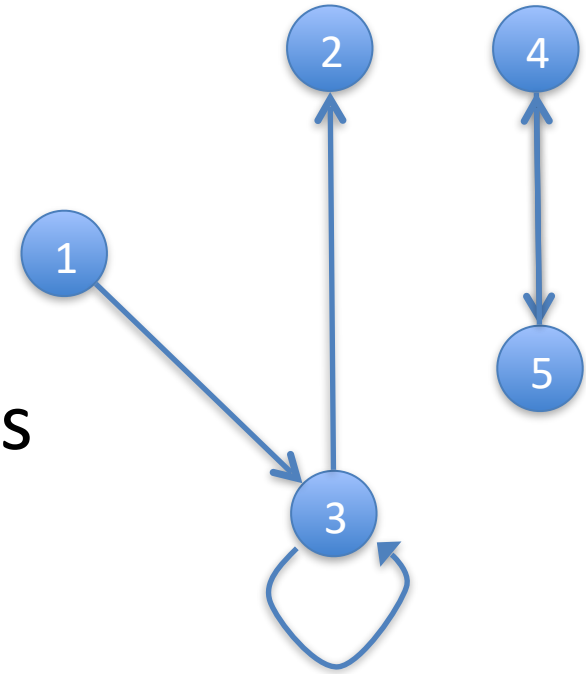
# Graphs and optimization problems based on graphs

- *Many* important real-world problems can be modeled as optimization problems on **graphs**
- A graph is:
  - A set of nodes (vertices)
  - A set of edges (arcs) representing connections between pairs of nodes
- There are several types of graphs:
  - **Directed.** Edges are “one way” from source to destination)
  - **Undirected.** Edges have no particular direction – can travel either way, “see” each node from other, etc.
  - **Weighted.** Edges have associated numbers called weights that can be used to represent cost, time, flow capacity, etc.
- See, e.g., Ch. 2 of free online book – Think Complexity - <http://greenteapress.com/complexity/html/thinkcomplexity003.html>



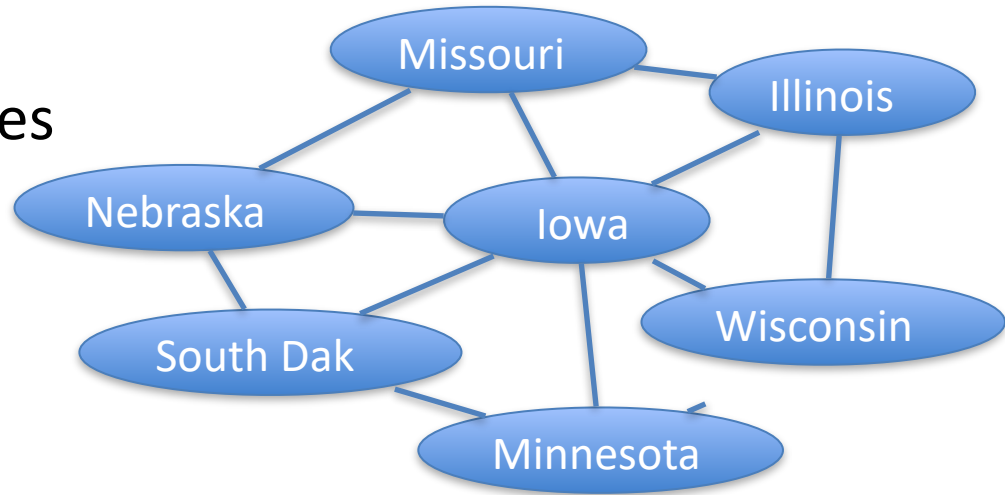
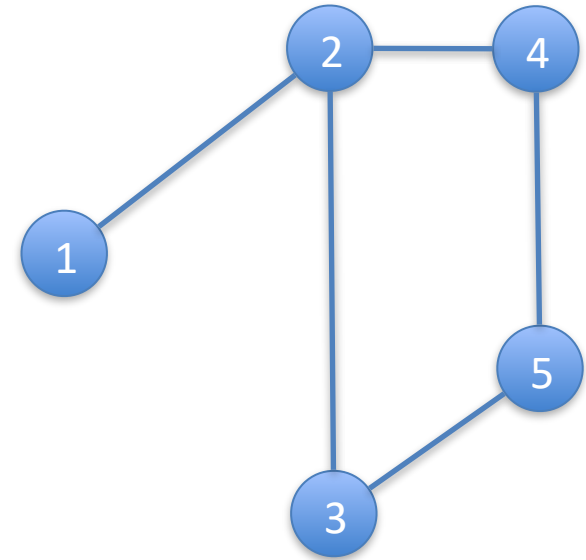
# Directed graph

- edges are “one way” from source to destination
- Can have two (one each way) between a pair of nodes
- Node can have edge to self
- Example relationships:
  - course prerequisite
  - hyperlink between web pages
  - street between intersections
  - Twitter follower
  - Infection spread from-to



## Undirected graph

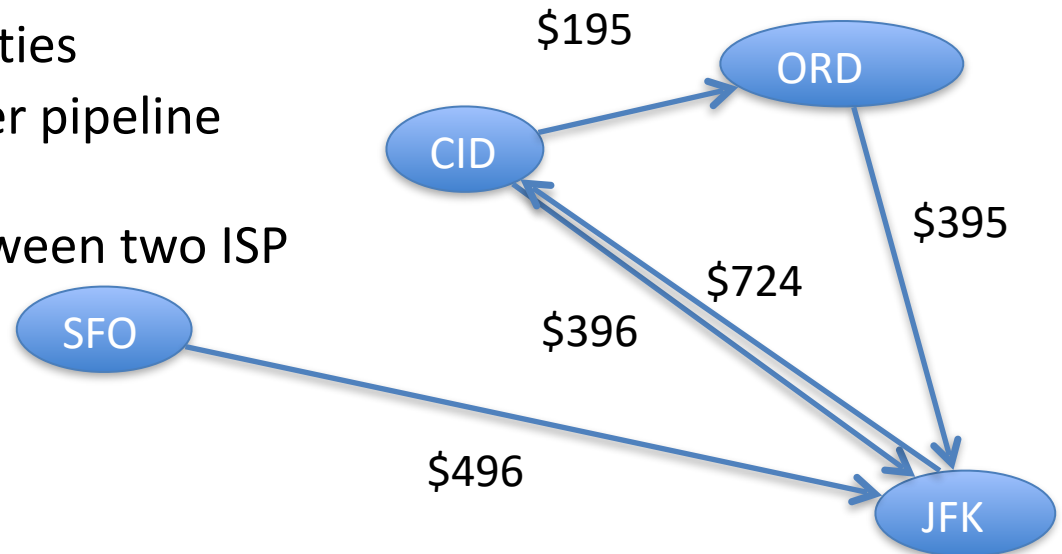
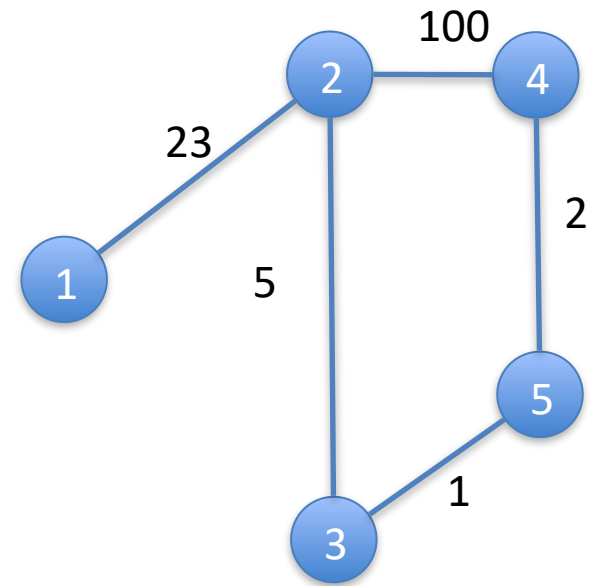
- Edges have no direction. Can “travel” either direction
- Can have only one edge between a pair of nodes\*
- Node cannot have edge to self
- Example relationships:
  - Facebook friend
  - Bordering countries/states



\*another kind of graph – multigraph – relaxes this rule

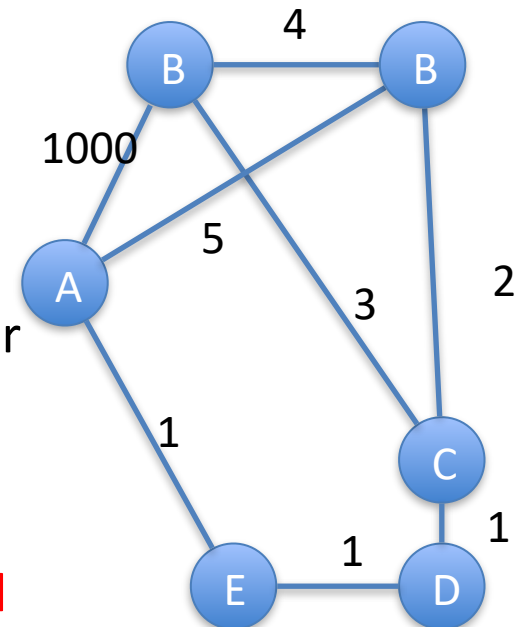
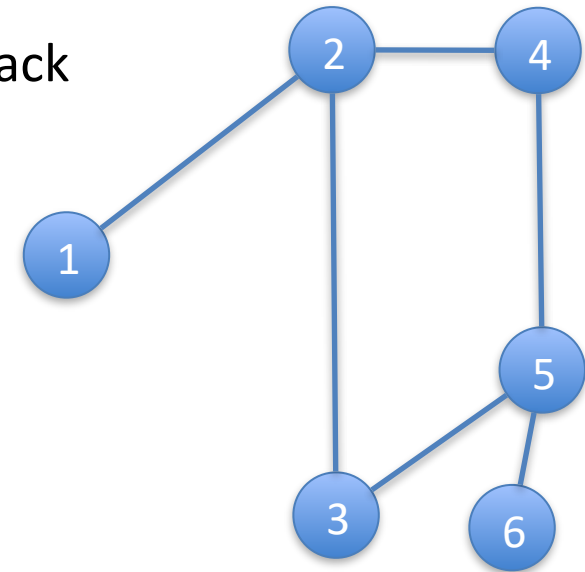
# Weighted graphs

- Variant of both directed and undirected graphs in which each edge has an associate number called a **weight** or cost
- Edge weight provides additional information about the relationship between the nodes.
- Example relationships:
  - Airfare between two cities
  - Distance between two cities
  - Flow capacity of oil/water pipeline between two points
  - Network bandwidth between two ISP nodes

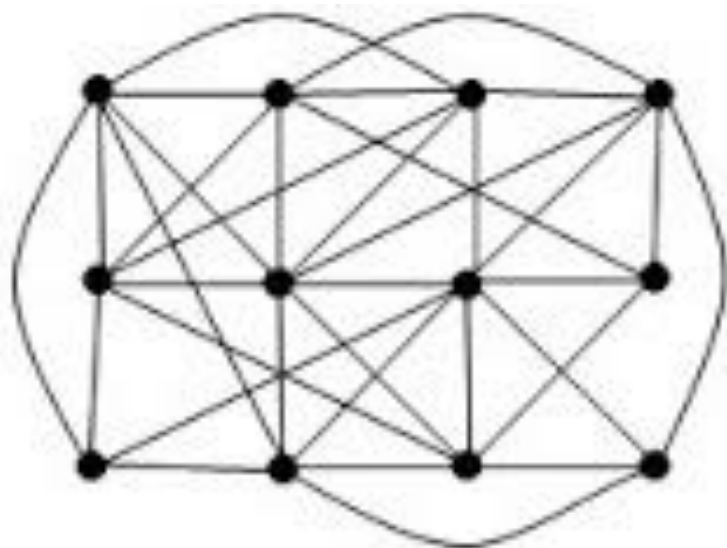
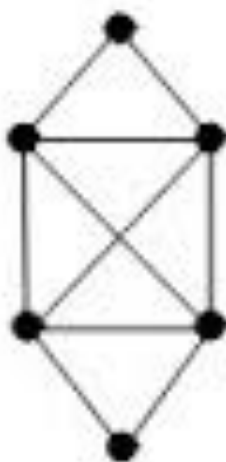
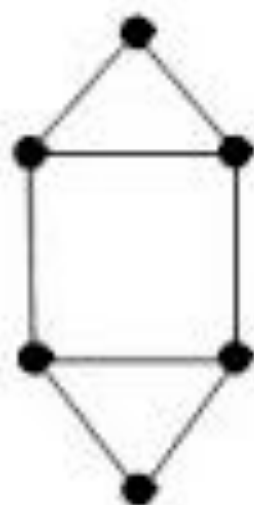
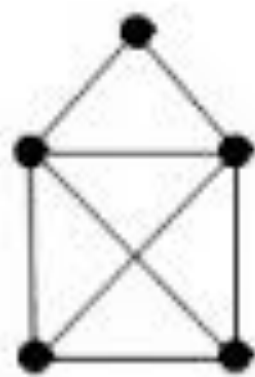
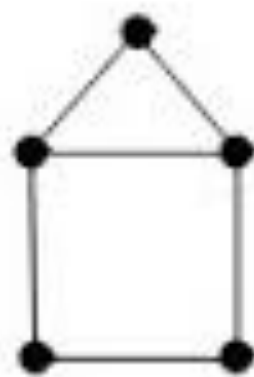


# Classic graph problems

- Determine if a graph has a cycle, a path that loops back to start points (e.g. 2-4-5-3-2)
- Find a path (non-branching) that traverses each (undirected) edge exactly once
  - [Leonhard Euler and the Bridges of Königsberg](#)
  - Not possible in graph on top right
- Find the shortest path between source  $s$  and destination
  - Different algorithms for weighted/unweighted graphs
- Find longest path between source and destination
- Find a path that visits each vertex exactly once
  - A, E, D, C, F, B, A in example on bottom right
- Path of minimum cost that visits each vertex once
  - A, E, D, C, B, F, A (cost 15) in example
- Assign no more than  $n$  different colors to vertices under constraint that no pair of connected vertices has the same color



Some of these are easy (have fast algorithms), others hard (no known efficient solution)

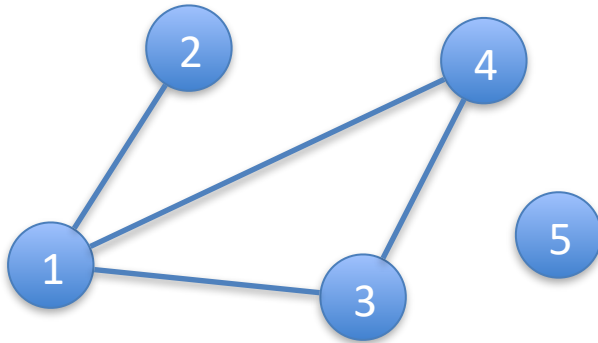


# Representing graphs

- How can we represent general graph in Python?
  - Need to keep track of nodes
  - Need to keep track of edges
- Several ways to represent graphs have been developed
  - List of nodes and list of edges
  - Adjacency matrix
  - Adjacency lists
  - Dictionary of dictionaries
  - Efficiency of algorithms that solve graph problems can vary greatly depending on how graph are representated
  - a strong influence on choice is the fact that one of the most common things needed in graph algorithms is access to immediate neighbors of a node (nodes that are destinations of edges for which “current” node is source)



# Adjacency matrix



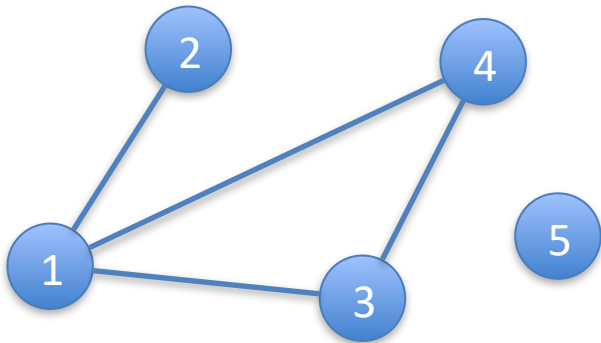
	1	2	3	4	5
1	False	True	True	True	False
2	True	False	False	False	False
3	True	False	False	True	False
4	True	False	True	False	False
5	False	False	False	False	False

- Appealingly simple to understand and implement
- Use, e.g. a list of lists containing True/False, 0/1, or similar
- NOT the most common graph representation for most problems. Can you think of a reason why?
  - Consider representing Facebook friends graph where each node is a FB user and an edge exists between two nodes whenever the two are FB friends.
  - One billion nodes. Adjacency matrix 1B x 1B in size! Your computer doesn't have that much storage. But FB graph **can** be represented in computer! How?
  - The 1B x 1B would be mostly False/0 – most people don't have huge number of friends. Should be representable in closer to 1B \* median number of friends. Other representations enable this huge memory savings.

# Adjacency list

Use a dictionary with

- Nodes as keys
- Values are lists of neighbor nodes



KEY	VALUE
1	[2, 4, 3]
2	[1]
3	[1, 4]
4	[3, 1]
5	[]

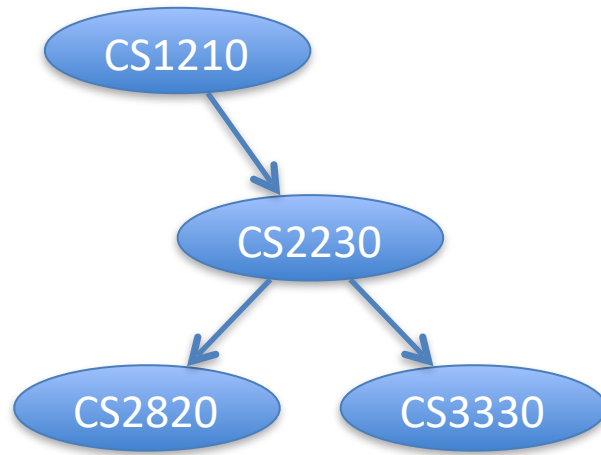
Compared to adjacency matrix:

+ Much less space (when, as is common, most nodes have only a small relatively small number of neighbors). Facebook graph. People have hundreds of friends, not many millions

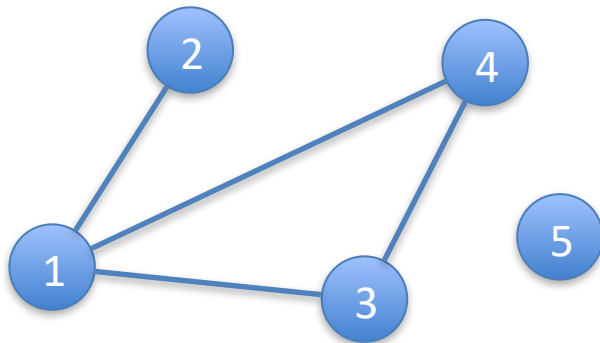
- Query of “does edge (i,j) exist?” not  $O(1)$ . Must search list associated with node i to see if j is there. Turns out this is not crucial in many graph algorithms. (could address this using dictionary of dictionaries but often not necessary)

# Adjacency list graph representation

Suitable for both undirected and directed graphs  
(and can be use for weighted graphs as well)



KEY	VALUE
CS2230	[CS2820, CS3330]
CS2820	[]
CS1210	[CS2230]
CS3330	[]



KEY	VALUE
1	[2, 4, 3]
2	[1]
3	[1, 4]
4	[3, 1]
5	[]

# An **adjacency list** representation for undirected graphs in Python

## Two classes: Node and Graph

[basicgraph.py](#)

### Node

- properties:
  - name : string
  - status: string (we'll use this to “mark” nodes during traversals)
- methods
  - getName
  - `__repr__` : we'll print nodes as *<name>*

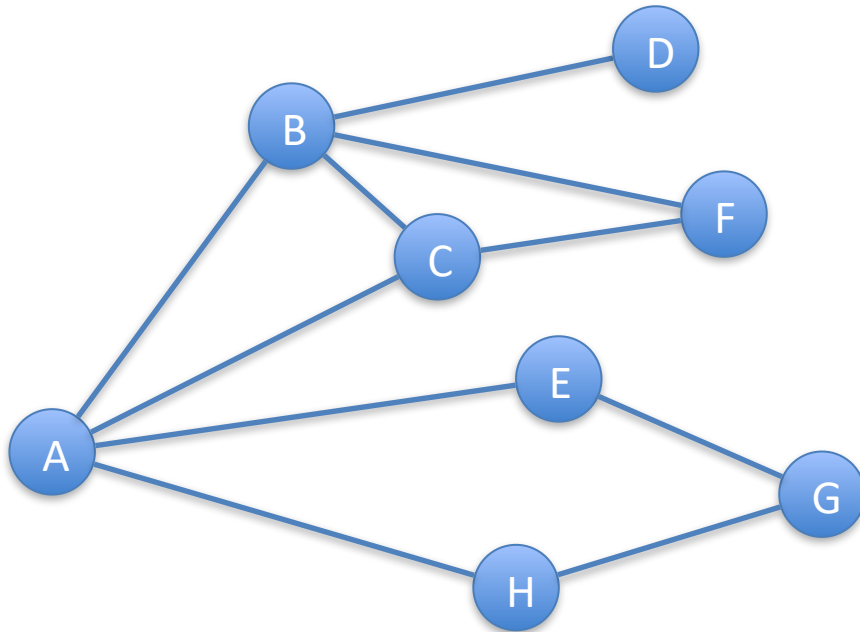
Note: in your HW8 you'll add one or more additional properties that help with traversing/walking through graphs to solve specific problems

# Adjacency list representation for undirected graphs

## Graph

- properties
  - nodes: a list of Node objects
  - adjacencyLists: a dictionary with all nodes as keys. The value associated with a key n1 (where n1 is a node) is a list of all the nodes, n2, for which (n1,n2) is an edge.
- methods
  - addNode(node) : nodes must be added to graphs before edges
  - addEdge(node1, node2) : presumes both nodes in graph already
  - neighborsOf(node) : returns list of neighboring nodes
  - getNode(name)
  - hasNode(node)
  - hasEdge(node1, node2) : return T if edge node1-node2 in graph
  - \_\_repr\_\_

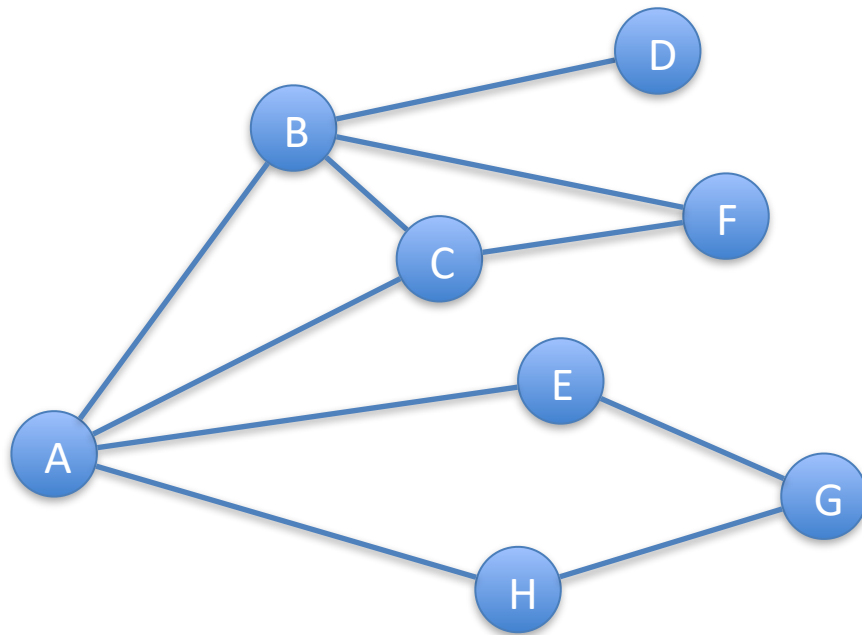
[basicgraph.py](#)



KEY	VALUE
A	[B,C,E,H]
B	[A,C,D,F]
C	[A,B,F]
D	[B]
E	[A,G]
F	[B,C]
G	[E,H]
H	[A,G]

This graph is generated by `genDemoGraph()` in `basicGraph.py`

*Note: for exams, you need to be able to 1) draw graph given adjacency list dictionary, and/or 2) show adjacency list dictionary given graph drawing*



KEY	VALUE
A	[B,C,E,H]
B	[A,C,D,F]
C	[A,B,F]
D	[B]
E	[A,G]
F	[B,C]
G	[E,H]
H	[A,G]

As I've said, many real-world problems can be represented as problems involving graphs. The algorithms to solve those problems often involve **graph traversals**, organized exploration or “walkthroughs” of the graph.

Two famous ones are: depth-first search and breadth-first search. I will present breadth-first search.

You will not be responsible for knowing the details of breadth-first search (for exam purposes) but you need to understand it well enough to *use and extend* it in HW8.

# Word ladder puzzles

CAT

???

???

DOG

CAT

COT

???

DOG

CAT

COT

DOT

DOG

Find 3-letter English words for ??? Positions. Each must differ from previous and next word in only one location

This problem is easily representable and solvable using graphs!



# Next time

- Graph traversals
  - Breadth first search
  - Depth first search
- HW 8